Decomposition of Monomial Algebras: Applications and Algorithms

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ABSTRACT. Considering finite extensions $K[A] \subseteq K[B]$ of positive affine semigroup rings over a field K we have developed in [BEN] an algorithm to decompose K[B] as a direct sum of monomial ideals in K[A]. By computing the regularity of homogeneous semigroup rings from the decomposition, we have confirmed the Eisenbud-Goto conjecture in a range of new cases not tractable by standard methods. Here we first illustrate this technique and its implementation in our *Macaulay2* package *MonomialAlgebras* by computing the decomposition and the regularity step by step for an explicit example. We then focus on ring-theoretic properties of simplicial semigroup rings. From the characterizations given in [BEN], we develop and prove explicit algorithms testing various properties, including being Buchsbaum, Cohen-Macaulay, Gorenstein, normal, and seminormal. All algorithms are implemented in our *Macaulay2* package.

INTRODUCTION. Let *B* be a positive affine semigroup, that is, *B* is a finitely generated subsemigroup of \mathbb{N}^m for some *m*. Let *K* be a field and *K*[*B*] the affine semigroup ring associated to *B*, which can be identified with the subring of *K*[t_1, \ldots, t_m] generated by monomials $t^u := t_1^{u_1} \cdots t_m^{u_m}$, where $u = (u_1, \ldots, u_m) \in B$. Denote by *C*(*B*) and by *G*(*B*) the cone and the group generated by *B*. From now on let $A \subseteq B$ be positive affine semigroups with C(A) = C(B). We will now discuss the decomposition of *K*[*B*] into a direct sum of monomial ideals in *K*[*A*]. Observe that

$$K[B] = \bigoplus_{g \in G} K \cdot \left\{ t^b \mid b \in B \cap g \right\},\$$

where G := G(B)/G(A). We have C(A) = C(B) if and only if K[B] is a finitely generated K[A]-module. From this it follows that *G* is finite, and we can compute the above decomposition since all summands are finitely generated. Moreover, there are shifts $h_g \in G(B)$ such that $I_g := K \cdot \{t^{b-h_g} \mid b \in B \cap g\}$ is a monomial ideal in K[A]. Thus, $K[B] \cong \bigoplus_{g \in G} I_g(-h_g)$ as \mathbb{Z}^m -graded K[A]-modules (with deg $t^b = b$). A detailed formulation of the algorithm computing the ideals I_g and shifts h_g and a more general version of the decomposition in the setup of cancellative abelian semigroup rings over an integral domain can be found in [BEN, Algorithm 1, Theorem 2.1].

Our original motivation for developing this decomposition was to provide a fast algorithm to compute the Castelnuovo-Mumford regularity reg K[B] of a homogeneous semigroup ring in order to test the Eisenbud-Goto conjecture [EG, Conjecture p. 93]. Recall that the *Castelnuovo-Mumford regularity* reg M of a finitely generated graded module M over a standard graded polynomial ring $R = K[x_1, \ldots, x_n]$ is defined as the smallest integer m such that every j-th syzygy module of M is generated by elements of degree at most m + j. Moreover, B is called a *homogeneous semigroup* if there exists a group homomorphism deg: $G(B) \to \mathbb{Z}$ with deg $b_i = 1$ for $i = 1, \ldots, n$, where

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MonomialAlgebras version 2.3.

Hilb(*B*) = { $b_1, ..., b_n$ } is the minimal generating set of *B*; by reg*K*[*B*] we mean its regularity with respect to the *R*-module structure which is given by the *K*-algebra homomorphism $R \rightarrow K[B], x_i \mapsto t^{b_i}$.

The toric Eisenbud-Goto conjecture can be formulated as follows: let *K* be a field and *B* a homogeneous semigroup, then reg $K[B] \leq \deg K[B] - \operatorname{codim} K[B]$, where deg K[B] denotes the degree and codim $K[B] := \dim_K K[B]_1 - \dim K[B]$ the codimension. Even this special case of the Eisenbud-Goto conjecture is largely open; for references on known results see [BEN, §4]. The regularity of K[B] is usually computed from a minimal graded free resolution. If *n* is large this computation is very expensive, and hence it is impossible to test the conjecture systematically in high codimension using this method. However, choosing *A* to be generated by minimal generators e_1, \ldots, e_d of C(B) of degree 1 the regularity can be computed as reg $K[B] = \max \{ \operatorname{reg} I_g + \deg h_g \mid g \in G \}$, where $\operatorname{reg} I_g$ denotes the regularity of I_g with respect to the canonical $T = K[x_1, \ldots, x_d]$ -module structure given by $T \to K[A], x_i \mapsto t^{e_i}$. Since the free resolution of every ideal I_g appearing has length at most d - 1, this computation is typically much faster than the traditional approaches. This enabled us to test the conjecture for a large class of homogeneous semigroup rings by using our regularity algorithm. See [BEN, §4] for details.

In first section, we illustrate a step by step decomposition and regularity computation for an explicit example using our *Macaulay2* [M2] package *MonomialAlgebras*. We say that K[B] is a *simplicial semigroup ring* if the cone C(B) is simplicial. In second section, we focus on simplicial semigroup rings K[B]. Based on the characterizations of ring-theoretic properties given in [BEN, Proposition 3.1] we develop explicit algorithms for testing whether K[B] is Buchsbaum, Cohen-Macaulay, Gorenstein, seminormal, or normal. We also discuss that, by known results, all these ring-theoretic properties imply the Eisenbud-Goto conjecture. The algorithms mentioned are implemented in our *Macaulay2* package.

DECOMPOSITION AND REGULARITY. Our Macaulay2 package can be loaded by

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Macaulay2, version 1.6
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with packages: ConwayPolynomials, Elimination, IntegralClosure, LLLBases,
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PrimaryDecomposition, ReesAlgebra, TangentCone

i1 : needsPackage "MonomialAlgebras";

--loading configuration for package "MonomialAlgebras" from file .../init-MonomialAlgebras.m2 --loading configuration for package "FourTiTwo" from file .../init-FourTiTwo.m2

We discuss the decomposition and computation of the regularity at the example of the homogeneous semigroup $B \subset \mathbb{N}^3$ specified by a list of generators

i2 : B = {{4,0,0},{2,2,0},{2,0,2},{0,2,2},{0,3,1},{3,1,0},{1,1,2}};

As an input for our algorithms we encode this data in a multigraded polynomial ring

i3 : K = ZZ/101;

i4 : S = K[x_1..x_7, Degrees => B];

The usual approach to computing $\operatorname{reg} K[B]$ is to obtain it from a minimal graded free resolution of the toric ideal I_B with respect to the standard grading:

```
i5 : IB = binomialIdeal S;
o5 : Ideal of S
```

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i6 : R = newRing(ring IB, Degrees => {7:1});
```

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i7 : betti res sub(IB,R)
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```
o7 : BettiTally
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Hence, we observe that $\operatorname{reg} K[B] = 3$. On the other hand, using the decomposition of K[B], we can proceed as follows: The command

```
o8 : HashTable
```

decomposes K[B] over K[A] where $A \subseteq B$ is generated by minimal generators of C(B) with minimal coordinate sum; so in the example $A = \langle (4,0,0), (2,2,0), (2,0,2), (0,2,2), (0,3,1) \rangle$. The keys of the hash table represent the elements of *G* and the values are the tuples (I_g, h_g) , hence

(†)
$$K[B] \cong \langle \overline{x}_1, \overline{x}_3 \rangle \left(-(-1, 1, 0) \right) \oplus K[A]$$

as \mathbb{Z}^3 -graded K[A]-modules; here we write $K[A] \cong T/J$ with $T = K[x_1, x_2, x_3, x_4, x_5]$ and \overline{x}_i for the class of x_i . The on-screen output of *Macaulay2* does not distinguish between the class and the representative. With respect to the standard grading deg $u = (u_1 + u_2 + u_3)/4$, Equation (†) yields reg $K[B] = \max \{ \operatorname{reg} \langle \overline{x}_1, \overline{x}_3 \rangle + (-1+1+0)/4, \operatorname{reg} K[A] + 0 \}$. We compute reg K[A]:

```
i9 : KA = ring first first values dc;
   i10 : T = newRing(ring ideal KA, Degrees => {5:1});
   i11 : J = sub(ideal KA,T);
   oll : Ideal of T
   il2 : betti res J
                  012
   o12 = total: 1 3 2
              0:1..
              1: . 1 .
              2: . 2 2
   o12 : BettiTally
Hence \operatorname{reg} K[A] = 2. We can compute \operatorname{reg} \langle \overline{x}_1, \overline{x}_3 \rangle as follows:
   i13 : I1 = first (values dc)#0
   o13 = ideal(x, x)
                   1
                        3
   ol3 : Ideal of KA
   i14 : g = matrix entries sub(gens I1, T);
                            2
                   1
   o14 : Matrix T <--- T
   i15 : betti res image map(coker gens J, source g, g)
```

```
0 1 2 3
o15 = total: 2 5 4 1
1: 2 2 . .
2: . . . .
3: . 3 4 1
o15 : BettiTally
```

Thus, reg $\langle \overline{x}_1, \overline{x}_3 \rangle = 3$, and therefore we see again that reg K[B] = 3. Observe that the resolution of K[B] has length 5, whereas the ideals I_g have resolutions of length at most 3. The command

provides an implementation of this approach, also returning the tuples (I_g, h_g) where the maximum is achieved. By [BEN, Proposition 4.1] we have deg $K[B] = #G \cdot \text{deg } K[A] = 10$ since

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i17 : degree J
o17 = 5
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Moreover, $\operatorname{codim} K[B] = 4$ since $\dim K[B] = \dim C(B) = 3$. Hence the ring K[B] satisfies the Eisenbud-Goto bound.

ALGORITHMS FOR RING THEORETIC PROPERTIES. In this section, we focus on simplicial semigroup rings K[B]. Based on the characterizations given in [BEN, Proposition 3.1], we develop and prove explicit algorithms for testing whether K[B] is Buchsbaum, Cohen-Macaulay, Gorenstein, seminormal, or normal. In the simplicial case, all these properties are independent of K and they imply the Eisenbud-Goto conjecture by results of [SV, Corollary p. 307], [T, Corollary 2.3 and Proposition 2.2], and [N, Theorem 3.16]. As an example, consider the following homogeneous simplicial semigroup $B \subset \mathbb{N}^3$ specified by the generators

i18 : B = {{4,0,0}, {0,4,0}, {0,0,4}, {1,0,3}, {0,2,2}, {3,0,1}, {1,2,1}};

We compute the decomposition of K[B] over K[A], where $A = \langle (4,0,0), (0,4,0), (0,0,4) \rangle \subset B$ is generated again by minimal generators of C(B) with minimal coordinate sum.

}

o20 : HashTable

Hence, $K[B] \cong K[A] \oplus K[A](-1)^4 \oplus K[A](-2)^2 \oplus \langle x_1, x_2, x_3 \rangle (-1)$ with respect to the standard grading induced by deg $u = (u_1 + u_2 + u_3)/4$. It follows that depth K[B] = 1; thus, K[B] is not Cohen-Macaulay. Hence, K[B] is also not normal by [H, Theorem 1]. We can test seminormality via Algorithm 1.

Algorithm 1 (Seminormality test).

Input: A simplicial semigroup $B \subseteq \mathbb{N}^m$.

Output: true if K[B] is seminormal, false otherwise.

- Let e₁,..., e_d ∈ B be minimal generators of C(B) with minimal coordinate sum, and set A := ⟨e₁,...,e_d⟩.
- 2: Compute $B_A := \{x \in B \mid x \notin B + (A \setminus \{0\})\}$ as described in [BEN, Algorithm 1, Step 1].
- 3: for all $x \in B_A$ do
- 4: Solve the linear system of equations $\sum_{i=1}^{d} \lambda_i e_i = x$ for $\lambda = (\lambda_1, \dots, \lambda_d) \in \mathbb{Q}^d$.
- 5: **if** $\|\lambda\|_{\infty} > 1$ **then return** false.
- 6: **end for**
- 7: return true.

Here, $\|-\|_{\infty}$ denotes the maximum norm. Note that all λ_i are non-negative since C(B) is a simplicial cone. Verifying in Step 5 the condition $\|\lambda\|_{\infty} \ge 1$ instead, results in an algorithm which tests normality. Using our package we observe that, in the example above, K[B] is not seminormal:

i21 : isSeminormalMA B
o21 = false

The Buchsbaum property can be tested by Algorithm 2. We denote by $K[A]_+$ the homogeneous maximal ideal of K[A].

Algorithm 2 (Buchsbaum test).

Input: A simplicial semigroup $B = \langle b_1, \ldots, b_n \rangle \subseteq \mathbb{N}^m$.

Output: true if K[B] is Buchsbaum, false otherwise.

- Let *e*₁,...,*e*_d ∈ *B* be minimal generators of *C*(*B*) with minimal coordinate sum, and set *A* := ⟨*e*₁,...,*e*_d⟩.
- 2: Using the (minimal) generators e_1, \ldots, e_d of A decompose $K[B] \cong \bigoplus_{g \in G} I_g(-h_g)$, where $I_g \subseteq K[A]$, $h_g \in G(B)$ and G = G(B)/G(A) by [BEN, Algorithm 1].
- 3: if there exists $g \in G$ with $I_g \neq K[A]$ and $I_g \neq K[A]_+$ then return false.

4: $H := \{h_g \mid g \in G \text{ with } I_g = K[A]_+\}.$

5: $C := \{b_1, \ldots, b_n\} \setminus \{0, e_1, \ldots, e_d\}.$

6: $H + C := \{h_g + b_i \mid h_g \in H, b_i \in C\}.$

7: **return** true if $(H+C) \cap H = \emptyset$ and false otherwise.

Proof of Algorithm 2. By [BEN, Proposition 3.1], the ring K[B] is Buchsbaum iff each ideal I_g is either equal to K[A], or to $K[A]_+$ and $h_g + b \in B$ for all $b \in \text{Hilb}(B)$. So, Step 3 is correct and we may now assume that $I_g = K[A]$ or $I_g = K[A]_+$ for all $g \in G$. Recall that $I_g = \{t^{\nu-h_g} \mid v \in \Gamma_g\}K[A]$ where $\Gamma_g = \{x \in B_A \mid x \in g\}$. Moreover, $\{t^{\nu-h_g} \mid v \in \Gamma_g\}$ is always a minimal generating set of I_g and $h_g = \sum_{k=1}^d \min\{\lambda_k^{\nu} \mid v \in \Gamma_g\}e_k$ where $\nu = \sum_{k=1}^d \lambda_k^{\nu}e_k$ with $\lambda_k^{\nu} \in \mathbb{Q}$. Since $h + e_k \in B_A$ for all $h \in H$ and all $k = 1, \ldots, d$, we have $H \cap B = \emptyset$. In case that $(H + C) \cap H \neq \emptyset$, we obtain $h + b \notin B$ for some $h \in H$ and some $b \in B \setminus \{0\}$, that is $h + \text{Hilb}(B) \not\subseteq B$. Hence, K[B] is not Buchsbaum.

In case that K[B] is not Buchsbaum, there is an $h \in H$ and some $b \in Hilb(B)$ such that $h + b \notin B$. It is now sufficient to show that $b \in C$ and $h + b \in H$. By the above argument, $b \in C$. Let $m_k = h + b + e_k$ for k = 1, ..., d. Suppose that $m_i \notin B_A$ for some $i \in \{1, ..., d\}$. Since $m_k - e_k \notin B$ for all k = 1, ..., d, necessarily $m_i - e_j \in B$ for some $j \neq i$. Consider $y = m_j - \sum_{k=1}^d n_k e_k \in B$ with $n_k \in \mathbb{N}$ such that $\sum_{k=1}^d n_k$ is maximal. By construction $y \in B_A$, moreover, $n_j = 0$ since $m_j - e_j \notin B$. In the same way if $x = m_i - e_j - \sum_{k=1}^d n_k e_k \in B$ with $\sum_{k=1}^d n_k$ maximal, then $x \in B_A$. Since $m_i, m_j \in g$ for some $g \in G$, we also have $x, y \in g$. Since $e_1, ..., e_d$ are linearly independent, we have $\lambda_j^y - \lambda_j^x \ge 2$. Moreover, since $t^{y-h_g}, t^{x-h_g} \in K[A]$ we get that t^{y-h_g} is not a linear form. Hence $I_g \neq K[A]$ and $I_g \neq K[A]_+$, thus, $m_k \in B_A$ for all k = 1, ..., d. We have $\#\Gamma_g \in \{1, d\}$ by minimality, hence $\Gamma_g = \{m_1, ..., m_d\}$. By construction, $h_g = h + b$ and $I_g = K[A]_+$, therefore $h + b \in H$.

In Step 2 of Algorithm 2, the shifts h_g and hence the ideals I_g are uniquely determined since e_1, \ldots, e_d are linearly independent. This is not true for arbitrary generating sets. Continuing the example, by [SV, Corollary p. 307] and

i22 : isBuchsbaumMA B
o22 = true

it follows that K[B] satisfies the Eisenbud-Goto conjecture. We can also read off from the decomposition the regularity and the Eisenbud-Goto bound: we have reg K[A] = 0 and reg $\langle x_1, x_2, x_3 \rangle = 1$, therefore reg $K[B] = \max\{0, 1, 2, 1+1\} = 2$. Moreover, deg K[B] is the number of ideals which occur in the decomposition, hence deg $K[B] - \operatorname{codim} K[B] = 8 - 4 = 4$.

When *B* is Buchsbaum, the regularity of K[B] is independent of the field *K* since all ideals in the decomposition are equal to the homogeneous maximal ideal or to K[A].

We finish this section by providing Algorithm 3 for testing the Gorenstein property.

Algorithm 3 (Gorenstein test).

Input: A simplicial semigroup $B \subseteq \mathbb{N}^m$.

Output: true if K[B] is Gorenstein, false otherwise.

- Let *e*₁,..., *e*_d ∈ *B* be minimal generators of *C*(*B*) with minimal coordinate sum, and set *A* := ⟨*e*₁,...,*e*_d⟩.
- 2: Using the (minimal) generators e_1, \ldots, e_d of A decompose $K[B] \cong \bigoplus_{g \in G} I_g(-h_g)$ where $I_g \subseteq K[A]$, $h_g \in G(B)$ and G = G(B)/G(A) by [BEN, Algorithm 1].
- 3: if there exists $g \in G$ with $I_g \neq K[A]$ then return false.
- 4: $H := \{h_g \mid g \in G\}.$
- 5: if $h \in H$ with maximal coordinate sum is not unique then return false.

- 6: Let $h \in H$ with maximal coordinate sum.
- 7: while $H \neq \emptyset$ do
- 8: Let $h_g \in H$.
- if $h h_g \notin H$ then return false. 9:
- $H := H \setminus \{h_g, h h_g\}.$ 10:
- 11: end while
- 12: return true.

Proof of Algorithm 3. By [BEN, Proposition 3.1] the ring K[B] is Gorenstein if and only if $I_g = K[A]$ for all $g \in G$ and H has a unique maximal element with respect to \leq given by $x \leq y$ if there is a $z \in B$ such that x + z = y. Note that $H = B_A$ since $I_g = K[A]$ for all $g \in G$. If there is a maximal element $h \in H$, then this element has maximal coordinate sum. If H has more than one element with maximal coordinate sum, then H does not have a unique maximal element. To complete the proof we need to show that an element $h_g \in H$ satisfies $h_g \leq h$ iff $h - h_g \in H$. But this follows from the fact that if $x \notin B_A$ then $x + y \notin B_A$ for all $x, y \in B$.

Performing Steps 1–3 of Algorithm 3 (and returning true afterwards) also gives a test for the Cohen-Macaulay property.

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