

Topological complexity of n points on a tree

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The topological complexity of a path-connected space X , denoted by $\text{TC}(X)$, can be thought of as the minimum number of continuous rules needed to describe how to move from one point in X to another. The space X is often interpreted as a configuration space in some real-life context. Here, we consider the case where X is the space of configurations of n points on a tree Γ . We will be interested in two such configuration spaces. In the first, denoted by $C^n(\Gamma)$, the points are distinguishable, while in the second, $\text{UC}^n(\Gamma)$, the points are indistinguishable. We determine $\text{TC}(\text{UC}^n(\Gamma))$ for any tree Γ and many values of n , and consequently determine $\text{TC}(C^n(\Gamma))$ for the same values of n (provided the configuration spaces are path-connected).

[57M15](#); [55R80](#), [57Q05](#)

1 Introduction

For any topological space X , let $P(X)$ be the space of continuous paths $\sigma: [0, 1] \rightarrow X$ equipped with the compact–open topology. There is a fibration $p: P(X) \rightarrow X \times X$ which sends a path σ to its endpoints: $p(\sigma) = (\sigma(0), \sigma(1))$. When studying the problem of motion planning within a topological space, one often wishes to find sections of this fibration. That is, one wishes to find functions $s: X \times X \rightarrow P(X)$ such that $p \circ s$ is the identity. Such a function takes a pair of points as input and produces a path between those points, hence the relation to motion planning. The continuity of a section s at a point $(x, y) \in X \times X$ means that if (x', y') is “close” to (x, y) , then the path $s(x', y')$ is “close” to the path $s(x, y)$. Unfortunately it is a rarity that such a function can be continuous over all of $X \times X$ (in fact, such a continuous section exists if and only if the space X is contractible; see [3]). This leads to the definition of topological complexity introduced by Farber in [3]:

Definition 1.1 For any path-connected space X , the topological complexity of X , denoted by $\text{TC}(X)$, is the smallest integer $k \geq 1$ such that there is a cover of $X \times X$ by open sets U_1, U_2, \dots, U_k and continuous sections $s_i: U_i \rightarrow P(X)$. If there is no such k , set $\text{TC}(X) = \infty$.

Such a collection of sets U_i and sections s_i is called a *motion planning algorithm*. Thus, $\text{TC}(X)$ is in some sense the smallest number of continuous rules required to describe how to move between any two points in X . The space X is often viewed as the space of configurations of some real-world system. One example is when X is the space of configurations of n robots which move around a factory along a system of one-dimensional tracks. Such a system of tracks can be interpreted as a graph Γ (a one-dimensional CW complex). There are two types of these configuration spaces: in the first, denoted by $C^n(\Gamma)$, the robots are distinguishable, and in the second, denoted by $\text{UC}^n(\Gamma)$, the robots are indistinguishable. In other words, in $C^n(\Gamma)$, both the points of Γ occupied by robots, and the specific robots which occupy those points are of interest, while in $\text{UC}^n(\Gamma)$, it is only required that each specified point in Γ is occupied by *some* robot, but it is irrelevant which specific robot occupies each point. There are different real-world situations in which one configuration space is preferable over the other, depending on whether or not the robots are to perform different tasks. Our main goal is to study the topological complexity of the configuration spaces $C^n(\Gamma)$ and $\text{UC}^n(\Gamma)$ when Γ is a tree (a tree is a connected graph which has no cycles). The topological complexity of these configuration spaces is related to the number of vertices of degree greater than 2 (the degree of a vertex is defined in Section 2). These vertices are called *essential vertices*, and $m(\Gamma)$ is the number of essential vertices in Γ . Here, an *arc* in Γ is a subspace homeomorphic to a nontrivial closed interval. We will be interested in certain collections of arcs in Γ , which are called *allowable* and will be defined in Definition 2.3. In Section 3, we establish the following:

Theorem 1.2 *Let Γ be a tree with $m := m(\Gamma) \geq 1$.*

(1) *Let k be the smallest integer such that there is a collection of oriented arcs $\{A_i\}_{i=1}^k$ which is allowable for the collection of all vertices of degree 3 in Γ . If there are no vertices of degree 3, let $k = 0$. Let $n \geq 2m + k$ be an integer. Then $\text{TC}(C^n(\Gamma)) = \text{TC}(\text{UC}^n(\Gamma)) = 2m + 1$.*

(2) *Let $n = 2q + \epsilon < 2m$, with $\epsilon \in \{0, 1\}$ and $q \geq 1$, let r be the number of vertices of degree greater than 3, and let s be the number of vertices of degree 3. Suppose one of these three cases hold:*

(a) $s \geq 2(q - r)$.

(b) (i) $s < 2(q - r)$, $\epsilon = 0$, and there is some $k \geq 1$ such that there exists a collection of oriented arcs $\{A_i\}_{i=1}^k$ with the following properties:

- (A) The endpoints of each A_l are (distinct) essential vertices, neither of which is an endpoint of any other $A_{l'}$.
 - (B) There are $r' \leq r$ vertices of degree greater than 3 which are not the endpoints of any A_l .
 - (C) There is a collection \mathcal{V} of degree-3 vertices, with $|\mathcal{V}| \geq q - r' - k$ such that $\{A_i\}_{i=1}^k$ is allowable for \mathcal{V} .
- (ii) $s < 2(q - r)$, $\epsilon = 1$, and there is an arc A_0 whose endpoints have no restrictions and whose interior includes a collection \mathcal{W}' of $s' \leq q$ distinct vertices of degree 3, and if $s' < q - r$, there are arcs A_1, \dots, A_k as above whose endpoints are also not vertices in \mathcal{W}' , and there is another collection of degree-3 vertices, \mathcal{W} , such that $\mathcal{W} \cap \mathcal{W}' = \emptyset$, $|\mathcal{W}| \geq q - r' - k - s'$ and $\{A_i\}_{i=1}^k$ is allowable for \mathcal{W} , where r' is as above.

Then $\text{TC}(C^n(\Gamma)) = \text{TC}(\text{UC}^n(\Gamma)) = 2q + 1$.

In [5], Farber proves a similar statement for the spaces $C^n(\Gamma)$, showing that

$$\text{TC}(C^n(\Gamma)) = 2m(\Gamma) + 1 \quad \text{for } m(\Gamma) \geq 1 \text{ and } n \geq 2m(\Gamma),$$

with an additional assumption that Γ is not homeomorphic to the letter Y if $n = 2$. His proof uses methods which differ from the ones used here, and only address the spaces $C^n(\Gamma)$ for n large. [Theorem 1.2](#) in some sense improves this, since it addresses both the spaces $C^n(\Gamma)$ and $\text{UC}^n(\Gamma)$, and includes some values of n less than $2m(\Gamma)$. Furthermore, [Corollary 3.7](#) determines the topological complexity of both configuration spaces of a tree with no vertices of degree 3 for all values of $n \geq 1$, provided the configuration spaces are connected. On the other hand, with $k \geq 1$ as above, Farber's results determine $\text{TC}(C^n(\Gamma))$ for $n = 2m, 2m + 1, \dots, 2m + k - 1$, while our result does not. We discuss in [Proposition 3.8](#) the extent to which the results of [Theorem 1.2](#) are the best one can achieve with the methods used here.

2 Configuration spaces of points on graphs

Consider a graph Γ , where, as above, a graph is a 1-dimensional CW complex. The zero-dimensional cells of Γ are the vertices, and the closures of the 1-dimensional cells are the edges. The degree of a vertex v is the number of edges which have v as exactly one of their endpoints plus twice the number of edges which have v as both endpoints. We will deal exclusively with finite graphs, so that the number of

vertices and edges is finite. An essential vertex is a vertex of degree equal to or greater than 3, and $m(\Gamma)$ is the number of essential vertices in Γ . Let $C^n(\Gamma)$ be the space of n -tuples of distinct points in Γ . That is,

$$C^n(\Gamma) = \overbrace{\Gamma \times \cdots \times \Gamma}^{n \text{ times}} - \Delta,$$

where $\Delta = \{(x_1, \dots, x_n) \in \Gamma \times \cdots \times \Gamma : x_i = x_j \text{ for some } i \neq j\}$. The space $C^n(\Gamma)$ will be called the *topological configuration space of n ordered points on Γ* . Similarly, let

$$D^n(\Gamma) = \overbrace{\Gamma \times \cdots \times \Gamma}^{n \text{ times}} - \tilde{\Delta},$$

where $\tilde{\Delta}$ consists of all products of cells in $\Gamma \times \cdots \times \Gamma$ whose closures intersect Δ . Thus, $D^n(\Gamma)$ consists of all products of cells $c_1 \times \cdots \times c_n$ such that $\bar{c}_i \cap \bar{c}_j = \emptyset$ whenever $i \neq j$. In what follows, the word ‘‘cell’’ will always refer to the closure of a cell. A point in $D^n(\Gamma)$ is then an ordered n -tuple of points (x_1, \dots, x_n) in Γ such that there is at least a full open edge between two distinct points x_i and x_j . The space $D^n(\Gamma)$ will be called the *discrete configuration space of n ordered points on Γ* .

There is a free action of the symmetric group S_n on $C^n(\Gamma)$ and $D^n(\Gamma)$ which permutes the coordinates. The quotients of these two spaces under this action are denoted by $UC^n(\Gamma)$ and $UD^n(\Gamma)$, and are called the *unordered topological and discrete configuration spaces*. Given a point $\mathbf{y} = (y_1, \dots, y_n) \in C^n(\Gamma)$, we may find neighborhoods U_i in Γ which contain y_i and satisfy $U_i \cap U_j = \emptyset$ for $i \neq j$. Then $U = U_1 \times \cdots \times U_n$ is a neighborhood of \mathbf{y} in $C^n(\Gamma)$ and for each $\alpha \in S_n$, we have $\alpha(U) = U_{\alpha(1)} \times \cdots \times U_{\alpha(n)}$. So, if $\alpha \neq \alpha' \in S_n$, then there is some i such that $\alpha(i) \neq \alpha'(i)$, and then $U_{\alpha(i)} \cap U_{\alpha'(i)} = \emptyset$, so $\alpha(U) \cap \alpha'(U) = \emptyset$. This implies that the quotient map $C^n(\Gamma) \rightarrow UC^n(\Gamma)$ is a covering space projection (see [10, Proposition 1.40]). If \mathbf{y} is a point in $D^n(\Gamma)$, by letting $U' = U \cap D^n(\Gamma)$ we see that the same is true of the quotient map $D^n(\Gamma) \rightarrow UD^n(\Gamma)$. The topological and discrete spaces are related by the following:

Theorem 2.1 [1; 2; 11] *Let Γ be a graph with at least n vertices. Suppose*

- (1) *each path between distinct vertices of degree not equal to 2 in Γ contains at least $n - 1$ edges, and*
- (2) *each loop at a vertex in Γ which is not homotopic to a constant map contains at least $n + 1$ edges.*

Then $C^n(\Gamma)$ and $UC^n(\Gamma)$ deformation retract onto $D^n(\Gamma)$ and $UD^n(\Gamma)$, respectively.

In [1], Abrams proves a slightly weaker version of [Theorem 2.1](#), assuming that each path as in item (1) contains $n + 1$ edges, and conjectures the stronger version given here. Kim, Ko, and Park prove this conjecture in [11]. A graph which satisfies these conditions is called *sufficiently subdivided for n* . Any graph can be made sufficiently subdivided for any n by adding enough degree-2 vertices; this has no effect on the topology of the graph or either of the topological configuration spaces. The following relates the connectivity of a graph Γ and the connectivity of the topological configuration spaces:

Theorem 2.2 [1] *Suppose Γ is a graph with at least one edge or at least $n + 1$ vertices. Then:*

- (1) $C^n(\Gamma)$ is path-connected if and only if Γ is connected and either
 - (a) $n = 1$,
 - (b) $n = 2$ and Γ is not homeomorphic to a closed interval, or
 - (c) $n \geq 3$ and Γ is not homeomorphic to a closed interval or a circle.
- (2) $UC^n(\Gamma)$ is path-connected if and only if Γ is connected.

It follows that if Γ is sufficiently subdivided, then in [Theorem 2.2](#), $C^n(\Gamma)$ and $UC^n(\Gamma)$ can be replaced with $D^n(\Gamma)$ and $UD^n(\Gamma)$. In fact, Abrams also shows that if Γ has at least $n + 1$ vertices, then $UD^n(\Gamma)$ is connected if and only if Γ is connected, regardless of subdivision. If Γ has exactly n vertices, one can easily find examples in which Γ is disconnected, but $UD^n(\Gamma)$ is a single point (so $UD^n(\Gamma)$ is connected). There are less trivial examples in [1] in which Γ is connected while $D^n(\Gamma)$ is disconnected.

Daniel Farley and Lucas Sabalka have studied extensively the homotopy and homology groups and the cohomology rings of the spaces $UD^n(\Gamma)$ for a sufficiently subdivided tree Γ using Forman's discrete Morse theory [9]. Recall a tree is a simply connected graph. We summarize some of their results which will be relevant here. From here on, Γ is a tree which is sufficiently subdivided for n .

First, an ordering on the vertices is constructed as follows. Embed the tree Γ in the plane, and let $*$ be a vertex of degree 1. Assign $*$ the number 0, travel away from $*$, and number the remaining vertices in order (starting with 1) when they are first encountered. Whenever an essential vertex is encountered, take the leftmost edge, and turn around when a vertex of degree 1 is encountered. For each edge e , let $\iota(e)$ and $\tau(e)$ be the two endpoints of e , with $\iota(e) < \tau(e)$. There is also a notion of *directions* from a vertex $v \neq *$ of degree d . These directions are a numbering of the edges

incident to v from 0 to $d - 1$, in increasing order clockwise around the vertex, with 0 being the direction on the geodesic segment from v to $*$. An example is given in Figure 1. The 0 -direction at each essential vertex is marked with an arrow; to avoid clutter, the other directions are not marked on the graph. Note this graph is only sufficiently subdivided for $n = 2$, since there is only one edge along the geodesic from vertex 12 to vertex 15.

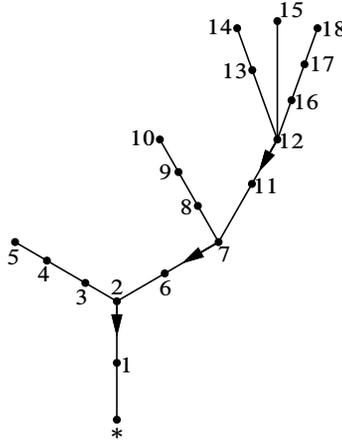


Figure 1: The ordering of the vertices on a tree Γ

Recall an arc in Γ is a subspace homeomorphic to a nontrivial closed interval. Farley and Sabalka’s notion of directions enables us to define the notion of an allowable collection of oriented arcs. Given a finite collection of oriented arcs $\{A_i\}_{i=1}^k$ in Γ and a vertex v in Γ of degree d , we will define integers $\eta_0(v), \eta_1(v), \dots, \eta_{d-1}(v)$ as follows. First, for $i = 1, 2, \dots, k$ and $j = 0, 1, \dots, d - 1$, if v falls on the arc A_i and A_i intersects the interior of the edge e_j incident to v in direction j , let $\eta_{j,i}(v) = 1$ if A_i is oriented towards v on e_j and $\eta_{j,i}(v) = -1$ if A_i is oriented away from v on e_j . If v does not fall on A_i or if A_i does not intersect the interior of e_j , let $\eta_{j,i}(v) = 0$. Then let

$$\eta_j(v) = \sum_{i=1}^k \eta_{j,i}(v).$$

This leads to the following definition:

Definition 2.3 Suppose $\mathcal{A} = \{A_i\}_{i=1}^k$ is a collection of oriented arcs in Γ and \mathcal{V} is a collection of vertices of Γ . The collection \mathcal{A} is said to be allowable for \mathcal{V} if every vertex $v \in \mathcal{V}$ of degree d has the property that v is not an endpoint of any A_i , and at least one of $\eta_0(v), \eta_1(v), \dots, \eta_{d-1}(v)$ is nonzero.

Intuitively, a collection of oriented arcs is allowable for a given collection \mathcal{V} of vertices if at each vertex in \mathcal{V} , there is at least one direction in which the orientations of the arcs don't "cancel out", and no vertex in \mathcal{V} is an endpoint of any arc. Also, if $\mathcal{V} = \emptyset$, then any collection of arcs is allowable for \mathcal{V} . We will be most interested in the case in which \mathcal{V} is a subset of the vertices of degree 3. **Figure 2** shows an example of two collections of oriented arcs in a graph Γ . The arcs are shown with a dashed line and orientations indicated with arrows, and have been moved away from Γ so that they are distinguishable in the figure. If \mathcal{V} is the collection of all essential vertices of Γ (which are all of degree 3), and \mathcal{W} is the collection of all essential vertices of Γ except the vertex labeled v , then both collections of arcs are allowable for \mathcal{W} , but the collection of arcs on the left is not allowable for \mathcal{V} , while the collection on the right is.

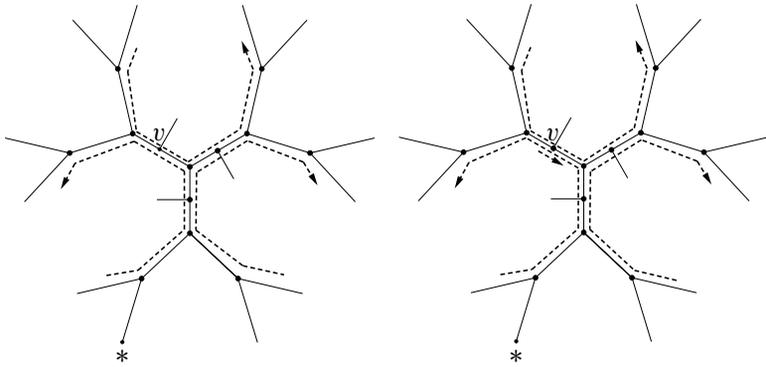


Figure 2: Two collections of arcs in a graph Γ

Note that given a collection of oriented arcs $\mathcal{A} = \{A_i\}_{i=1}^k$ which is allowable for \mathcal{V} and a vertex $v \in \mathcal{V}$ of degree d , if j is a direction satisfying $\eta_j(v) \neq 0$, then there must be some other direction $j' \neq j$ with $\eta_{j'}(v) \neq 0$. Indeed if v falls on an arc A_i , then v must fall on the interior of A_i , so that $\eta_{j_1,i}(v) = 1$ for some j_1 , and $\eta_{j_2,i}(v) = -1$ for some $j_2 \neq j_1$. Since Γ is a tree, for all other directions j , we have $\eta_{j,i}(v) = 0$, so A_i contributes 0 to the sum

$$\sum_{j=0}^{d-1} \sum_{i=1}^k \eta_{j,i}(v) = \sum_{j=0}^{d-1} \eta_j(v).$$

On the other hand, if v does not fall on A_i , then $\eta_{j,i}(v) = 0$ for all j , so again, A_i contributes 0 to the sum above. So, we have $\eta_0(v) + \eta_1(v) + \dots + \eta_{d-1}(v) = 0$, showing that there cannot be exactly one value j for which $\eta_j(v) \neq 0$. Furthermore,

for the same reason, if $\eta_j(v) \neq 0$, then there must be some direction $j' \neq j$ such that $\eta_j(v)$ and $\eta_{j'}(v)$ have opposite signs.

A cell c of $UD^n(\Gamma)$ can be described as a collection of vertices and edges, $c = \{c_1, \dots, c_n\}$, where each c_i is a vertex or an edge, and $c_i \cap c_j = \emptyset$ for $i \neq j$. Note the order in which the vertices and edges appear in c does not matter. The dimension of the cell c is the number of edges in c . Consider a cell $c = \{v, c_2, \dots, c_n\}$ containing some vertex $v \neq *$. If e_v is the unique edge in Γ which has $\tau(e_v) = v$, and $\{e_v, c_2, \dots, c_n\}$ is a valid cell in $UD^n(\Gamma)$, then v is said to be *unblocked* in c . Otherwise, v is *blocked* in c . The vertex $*$ is also said to be blocked in any cell containing it. In other words, v is unblocked in c if and only if $v \neq *$ and v can be replaced with the edge which contains v and is on the geodesic segment from v to $*$.

Now consider a cell $c = \{e, c_2, \dots, c_n\}$ which contains some edge e . If there is some vertex v in c which has the property that $\iota(e_v) = \iota(e)$ and $\iota(e) < v < \tau(e)$, then e is said to be *order-disrespecting* in c . Otherwise, the edge e is *order-respecting* in c . An example of 3 different cells in $UD^2(\Gamma)$ with Γ as in Figure 1 is given in Figure 3. The vertices/edges which are to be included in a cell c are labeled; all unlabeled vertices and edges are not included in the cell. In the left cell, the vertex 3 is blocked, since e_3 intersects the edge labeled e . Also, the edge e is order-disrespecting since $\iota(e_3) = 2 = \iota(e)$ and $\iota(e) = 2 < 3 < \tau(e) = 6$. In the middle cell, the vertex 10 is blocked, but 9 is unblocked. In the cell on the right, the vertex 16 is blocked, but f is order-respecting, since, although $\iota(e_{16}) = \iota(f)$, we have $\tau(f) = 13 < 16$.

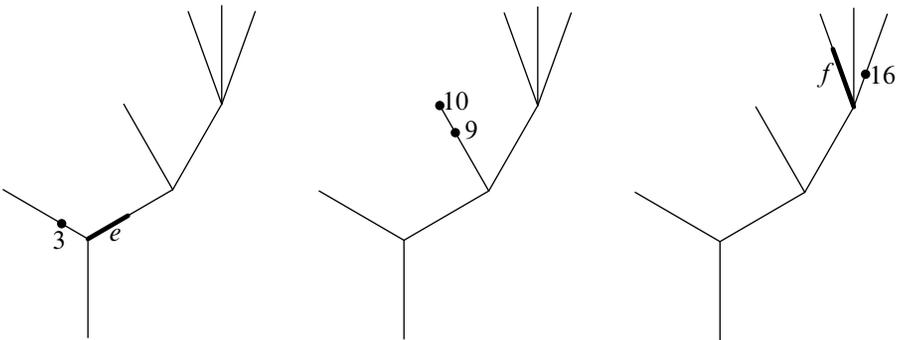


Figure 3: Three different cells in $UD^2(\Gamma)$

Farley and Sabalka construct a discrete vector field W and prove the following classification of the critical, collapsible, and redundant cells in $UD^n(\Gamma)$ with respect to W .

The terms *critical*, *collapsible*, and *redundant* come from discrete Morse theory, but their definitions will not be needed here.

Theorem 2.4 [7] *A cell c of $UD^n(\Gamma)$ is critical if and only if each vertex in c is blocked and each edge in c is order-disrespecting. A cell c is collapsible if and only if it contains some order-respecting edge e with the property that any unblocked vertex v in c satisfies $v > \tau(e)$. All other cells are redundant.*

Since we will be focused primarily on critical cells, we describe a procedure to construct a critical k -cell in $UD^n(\Gamma)$. First, notice that if c is any k -cell, then c must consist of k edges and $n - k$ vertices. If c is to be critical, each edge must be order-disrespecting, and each vertex must be blocked. Consider an edge e in c . If e is to be order-disrespecting, then $\iota(e)$ must be an essential vertex, or else there would be no possible vertex v that could satisfy $\iota(e_v) = \iota(e)$ other than $v = \tau(e)$, but $\tau(e)$ cannot be included in a cell which contains e (by this, we mean that $\tau(e)$ cannot appear as a vertex in the list of vertices and edges which define c). Furthermore, the direction d from $\iota(e)$ on which e falls must be at least 2, and the direction d' from $\iota(e)$ on which v falls must satisfy $0 < d' < d$. Note also that if v causes e to be order-disrespecting, then v is automatically blocked. This also implies that if there is a critical k -cell in $UD^n(\Gamma)$, then $k \leq m(\Gamma)$ and $n \geq 2k$. The remaining $n - 2k$ vertices of c can be easily chosen so that they are blocked in c . The cell on the left in Figure 3 is a critical 1-cell in $UD^2(\Gamma)$. In Figure 4 is an example of a critical 3-cell in $UD^8(\Gamma)$. Note that more vertices of degree 2 must be added to the tree Γ above so that it is sufficiently subdivided for $n = 8$. In this example and what follows, we make no indication of the total number of vertices in a sufficiently subdivided tree.

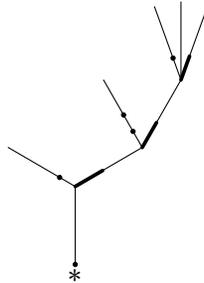


Figure 4: A critical 3-cell c in $UD^8(\Gamma)$

The discussion above is similar to the proof of the following theorem:

Theorem 2.5 [7] *Let Γ be a tree, and let $k = \min\{\lfloor \frac{n}{2} \rfloor, m(\Gamma)\}$. If c is any critical cell in $\text{UD}^n(\Gamma)$, then $\dim(c) \leq k$. Furthermore, $\text{UD}^n(\Gamma)$ deformation retracts onto the space $(\text{UD}^n(\Gamma))'_k$, which consists of the k -skeleton of $\text{UD}^n(\Gamma)$ with the redundant k -cells removed.*

The second statement of this theorem follows from results from discrete Morse theory. Now, before discussing the cohomology ring $H^*(\text{UD}^n(\Gamma); \mathbb{Z})$, we describe the equivalence relation on cells given in [8]. Given two cells c and c' of $\text{UD}^n(\Gamma)$, define \sim by $c \sim c'$ if and only if c and c' share the same edges (so in particular c and c' are of the same dimension), and if E is the set of edges in c (and in c'), and $\mathcal{C} = \Gamma - E$, then for every connected component C of \mathcal{C} , the number of vertices of c in C equals the number of vertices of c' in C . Here and in what follows, we use E to denote both the set of edges and the union of the edges in E . Context should make the desired interpretation of E clear. Let $[c]$ denote the equivalence class of c . Now, given two equivalence classes $[c]$ and $[d]$, write $[d] \leq [c]$ if there are representatives $c \in [c]$ and $d \in [d]$ such that d is obtained from c by removing some (possibly zero) edges of c and replacing each of these edges with one of its endpoints. Farley and Sabalka show the following:

Lemma 2.6 [8] *The relation \leq is a well-defined partial order with the following properties:*

- (1) *If a collection of distinct equivalence classes of 1-cells $\{[c_1], \dots, [c_k]\}$ has an upper bound, then it has a least upper bound and if e_i is the unique edge in $[c_i]$ (that is, every cell in $[c_i]$ contains the edge e_i), then $e_{i_1} \cap e_{i_2} = \emptyset$ for $i_1 \neq i_2$.*
- (2) *For any k -cell c of $\text{UD}^n(\Gamma)$, there is a unique collection of equivalence classes of 1-cells $\{[c_1], [c_2], \dots, [c_k]\}$ having $[c]$ as its least upper bound.*

Farley and Sabalka also introduced the idea of a “cloud diagram” to represent an equivalence class. These diagrams consist of a collection E of edges and an indication of the number of vertices in each connected component of $\Gamma - E$. The components of $\Gamma - E$ are called *clouds*. If $f(C)$ is the number of vertices in a cloud C , then

$$\sum_{\text{clouds } C} f(C) = n - |E|.$$

Figure 5 gives three examples of cloud diagrams. The cloud diagram on the left represents the class $[c]$, with c as in Figure 4.

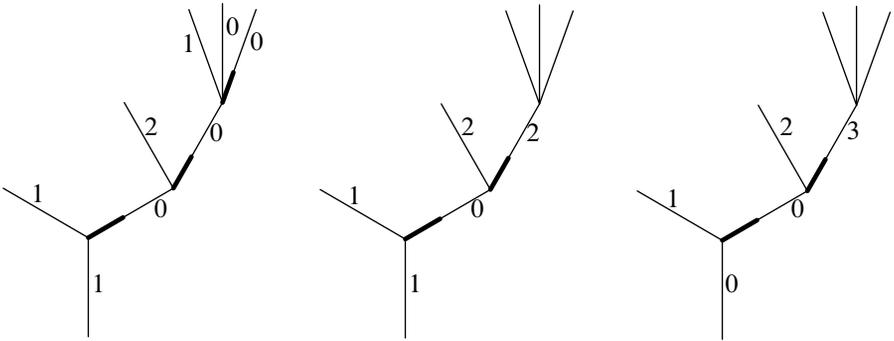


Figure 5: Cloud diagrams for classes $[c]$ with c as in Figure 4 (left), $[d]$ (middle) and $[d']$ (right)

We will sometimes call the number of vertices in a cloud C the *value of C* . Cloud diagrams also provide a convenient way to determine if $[d] \leq [c]$. If this is the case, then the set of edges in the diagram for $[d]$ must be a subset of the set of edges in the diagram for $[c]$, which implies that each cloud in the diagram for $[c]$ must be contained in some cloud in the diagram for $[d]$. For each cloud D in the diagram for $[d]$, the number of edges of $[c]$ which are contained in D plus the sum of the values of the clouds of $[c]$ which are contained in D must equal the value of D . In particular, if the set of edges in $[c]$ equals the set of edges in $[d]$, then the diagrams for $[c]$ and $[d]$ have the same clouds, and $[c]$ and $[d]$ are comparable if and only if the values of each cloud are the same in both diagrams, in which case $[c] = [d]$. In Figure 5, the middle diagram is a cloud diagram for a class $[d]$ with $[d] \leq [c]$ and the diagram on the right is a cloud diagram for a class $[d']$ which is comparable to neither $[c]$ nor $[d]$.

Farley and Sabalka determine the structure of the cohomology ring $H^*(UD^n(\Gamma); \mathbb{Z})$ by first constructing a space $\widehat{UD^n(\Gamma)}$ as follows: For each equivalence class of 1-cells $[c]$, let $S_{[c]}^1$ denote a circle with the usual cell structure consisting of a single open 1-cell $e_{[c]}^1$ and a single 0-cell. Then, each open k -cell of the product $\prod_{[c]} S_{[c]}^1$, taken over all equivalence classes of 1-cells of $UD^n(\Gamma)$, is of the form $e_{[c_1]}^1 \times \cdots \times e_{[c_k]}^1$, where we refrain from writing factors corresponding to 0-cells, and such a cell corresponds to a collection of equivalence classes of 1-cells $\{[c_1], \dots, [c_k]\}$. The space $\widehat{UD^n(\Gamma)}$ is obtained from $\prod_{[c]} S_{[c]}^1$ by removing open k -cells of the form $e_{[c_1]}^1 \times \cdots \times e_{[c_k]}^1$ such that the corresponding collection $\{[c_1], \dots, [c_k]\}$ does not have an upper bound. Then, each k -cell σ in $\widehat{UD^n(\Gamma)}$ corresponds to a collection $\{[c_1], \dots, [c_k]\}$ which has an upper bound, and therefore a least upper bound $[c]$, and the cell σ can be labeled

by $[c]$. For each distinct equivalence class $[c]$ of cells in $\text{UD}^n(\Gamma)$, there is exactly one cell labeled $[c]$ in $\widehat{\text{UD}^n(\Gamma)}$.

Now, for a k -cell labeled $[c]$ in $\widehat{\text{UD}^n(\Gamma)}$, let $\hat{\phi}_{[c]}$ be the k -cocycle in $C^*(\widehat{\text{UD}^n(\Gamma)}; \mathbb{Z})$ defined by

$$\hat{\phi}_{[c]}([c']) = \begin{cases} 1 & \text{if } [c'] = [c], \\ 0 & \text{if } [c'] \neq [c]. \end{cases}$$

We have the following description of $H^*(\widehat{\text{UD}^n(\Gamma)}; \mathbb{Z})$:

Theorem 2.7 [8; 6] *Let $\{[c_1], \dots, [c_M]\}$ be the collection of all equivalence classes of 1-cells in $\text{UD}^n(\Gamma)$. The cohomology ring $H^*(\widehat{\text{UD}^n(\Gamma)}; \mathbb{Z})$ is isomorphic to the quotient ring*

$$\Lambda[[c_1], \dots, [c_M]]/I,$$

where $\Lambda[[c_1], \dots, [c_M]]$ is the integral exterior ring generated by the collection of all equivalence classes of 1-cells and I is the ideal generated by products $[c_{i_1}] \cdot [c_{i_2}] \cdots [c_{i_k}]$ such that the collection $\{[c_{i_1}], [c_{i_2}], \dots, [c_{i_k}]\}$ does not have an upper bound.

The isomorphism $H^*(\widehat{\text{UD}^n(\Gamma)}) \rightarrow \Lambda[[c_1], \dots, [c_M]]/I$ sends $\hat{\phi}_{[c]}$ to $[c_1] \cdot [c_2] \cdots [c_k]$, where $\{[c_1], [c_2], \dots, [c_k]\}$ is the unique collection of equivalence classes of 1-cells which has $[c]$ as its least upper bound, arranged so that $\iota(e_i) < \iota(e_{i+1})$ for each i , where e_i the unique edge in $[c_i]$.

The isomorphism in [Theorem 2.7](#) depends on a choice of orientations of the cells and an ordering of the factors in $\widehat{\text{UD}^n(\Gamma)}$. The details are given in [8] and [6], but will be omitted here.

Similarly, for each equivalence class $[c]$ of k -cells of $\text{UD}^n(\Gamma)$, define a cellular cocycle $\phi_{[c]} \in C^*(\text{UD}^n(\Gamma); \mathbb{Z})$ by

$$\phi_{[c]}(c') = \begin{cases} 1 & \text{if } c' \sim c, \\ 0 & \text{otherwise.} \end{cases}$$

These cocycles will be called standard cocycles, and if there is a (unique) critical cell in $[c]$, then $\phi_{[c]}$ is called a critical cocycle. Since standard cocycles are determined by equivalence classes, cloud diagrams can also be used to describe standard cocycles.

Theorem 2.8 [8; 6] (1) *There is a well-defined map $q: \text{UD}^n(\Gamma) \rightarrow \widehat{\text{UD}^n(\Gamma)}$, and the induced homomorphism $q^*: C^*(\widehat{\text{UD}^n(\Gamma)}) \rightarrow C^*(\text{UD}^n(\Gamma))$ sends the cocycle $\hat{\phi}_{[c]}$ to the standard cocycle $\phi_{[c]}$.*

- (2) The collection of critical cocycles represents a basis for $H^*(UD^n(\Gamma); \mathbb{Z})$.
- (3) For any cell c , we have $\phi_{[c]}^2 = 0$.
- (4) If c is a k -cell and $\{[c_1], \dots, [c_k]\}$ is the unique collection of equivalence classes of 1-cells with $[c]$ as its least upper bound, arranged so that $\iota(e_i) < \iota(e_{i+1})$ for each i , where e_i is the unique edge in $[c_i]$, then

$$\phi_{[c_1]} \cdots \phi_{[c_k]} = \phi_{[c]}.$$

- (5) If $\{[c_1], \dots, [c_j]\}$ is any collection of equivalence classes of cells with no upper bound, then

$$\phi_{[c_1]} \cdots \phi_{[c_j]} = 0.$$

Here, for any equivalence class $[c]$, we use $\phi_{[c]}$ to denote both the standard cocycle $\phi_{[c]}$ and the cohomology class it represents. It follows from the universal coefficient theorem that the same statements hold true for rational cohomology, where we identify $\phi_{[c]} \in H^*(UD^n(\Gamma); \mathbb{Z})$ with the corresponding class in $H^*(UD^n(\Gamma); \mathbb{Q})$. Again, [Theorem 2.8](#) depends on a choice of orientations of cells, but we will omit these details.

If $\{[c_1], \dots, [c_k]\}$ is a collection of equivalence classes of 1-cells which has a least upper bound $[\omega]$, then $\phi_{[c_1]} \cdots \phi_{[c_k]} = \phi_{[\omega]}$, and if $[\omega]$ contains a critical cell, then $\phi_{[\omega]}$ is a critical cocycle, and therefore represents a basis element in $H^*(UD^n(\Gamma))$. If $[\omega]$ does not contain a critical cell, then $\phi_{[\omega]}$ is cohomologous to a linear combination of critical cocycles. In [\[6\]](#), Farley gives a procedure to rewrite the cohomology class of $\phi_{[\omega]}$ in terms of critical cocycles, which we recall here.

If $[c]$ is an equivalence class of k -cells and e is an edge of $[c]$ with $\iota(e) = v$ and $\deg(\tau(e)) \leq 2$, and \mathcal{C} is the cloud diagram for the standard k -cocycle $\phi_{[c]}$, then a $(k-1)$ -dimensional cochain $\mathcal{R}_{\mathcal{C},v}$ is defined as follows. The support of $\mathcal{R}_{\mathcal{C},v}$ consists of $(k-1)$ -cells c_0 such that

- (i) $E(c_0) = E(c) - \{e\}$ (where for any cell σ of $UD^n(\Gamma)$, $E(\sigma)$ denotes the set of edges in σ),
- (ii) if C is any component of $\Gamma - E(c)$, other than the component C_l which falls in the 0-direction from v or the component C_τ which is adjacent to $\tau(e)$, then the number of vertices of c_0 in C equals the number of vertices of c in C ,
- (iii) the number of vertices of c_0 in $C_l \cup \{\iota(e)\}$ equals the number of vertices of c in C_l , and
- (iv) the number of vertices of c_0 in $C_\tau \cup \{\tau(e)\}$ is one more than the number of vertices of c in C_τ .

For each cell c_0 that satisfies these conditions, put $\mathcal{R}_{C,v}(c_0) = 1$. Let $E(\mathcal{R}_{C,v})$ denote the set of edges in any cell in the support of $\mathcal{R}_{C,v}$. The cochain $\mathcal{R}_{C,v}$ can be described with a cloud diagram, where the union of the clouds around v forms a connected component of $\Gamma - E(\mathcal{R}_{C,v})$. For example, the left side of Figure 6 gives a cloud diagram \mathcal{C} for a standard cocycle $\phi_{[c]}$ (which is not a critical cocycle), and the right side gives the cloud diagram for the cochain $\mathcal{R}_{C,v}$, where $v = \iota(e)$. Here, we have emphasized the clouds by indicating them with dotted lines. For each cloud C in the diagram for $\mathcal{R}_{C,v}$, let $f_{C,v}(C)$ denote the number of vertices in C .

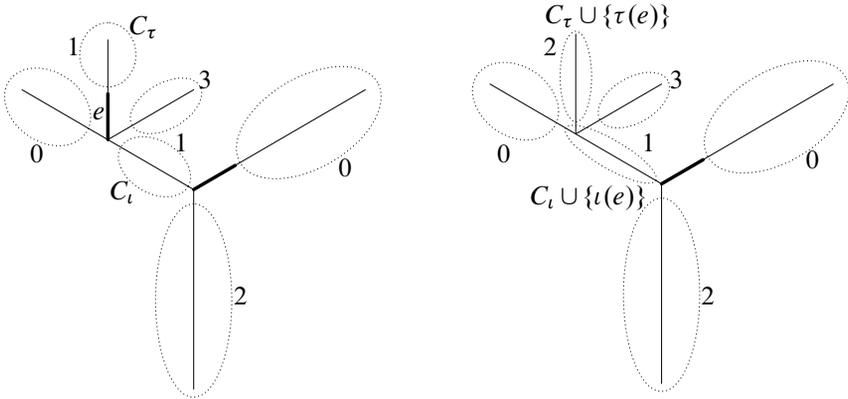


Figure 6: The cloud diagram \mathcal{C} for a standard cocycle $\phi_{[c]}$ (left) and the cloud diagram for the cochain $\mathcal{R}_{C,v}$ (right)

Farley shows that, up to sign, the coboundary $\delta(\mathcal{R}_{C,v})$ is given by the following, where $d(v)$ denotes the degree of v :

$$\delta(\mathcal{R}_{C,v}) = \sum_{i=1}^{d(v)-1} \Theta_{C,v,\tau,i} - \sum_{i=1}^{d(v)-1} \Theta_{C,v,\iota,i},$$

where $\Theta_{C,v,\tau,i}$ is the standard k -cocycle $\phi_{[c']}$ such that $E(c') = E(\mathcal{R}_{C,v}) \cup \{e'_i\}$ and e'_i is the edge in direction i from v , and if C_i is the cloud in direction i from v in the cloud diagram for $\phi_{[c']}$, then the number of vertices in any cloud C' in the cloud diagram for $\phi_{[c']}$ is given by

$$f_{[c']}(C') = \begin{cases} f_{C,v}(C') & \text{if } C' \neq C_i, \\ f_{C,v}(C') - 1 & \text{if } C' = C_i. \end{cases}$$

Note there is a slight abuse of notation, in the sense that the clouds in the diagram for $\mathcal{R}_{C,v}$ are slightly different than those in the diagram for $\phi_{[c']}$. For example, the cloud in direction 0 from v in the diagram for $\mathcal{R}_{C,v}$ includes $\iota(e)$, whereas the cloud

in direction 0 from v in the diagram for $\phi_{[c']}$ does not include this vertex. This should cause no confusion. It is possible that $f_{[c']}(C')$ is negative; if this is the case, then $\phi_{[c']}$ is defined to be zero. For example, Figure 7 gives the sum $\sum_{i=1}^{d(v)-1} \Theta_{C,v,\tau,i}$ with C as in Figure 6; here, $\Theta_{C,v,\tau,1}$ is zero.

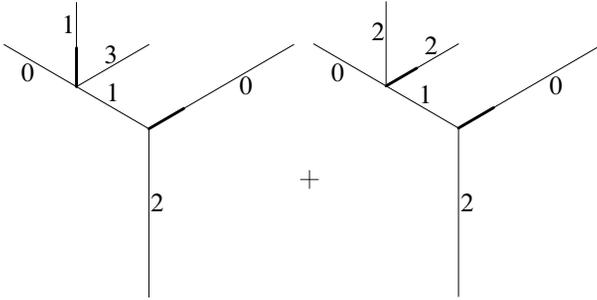


Figure 7: The sum $\sum_{i=1}^{d(v)-1} \Theta_{C,v,\tau,i}$ with C as in Figure 6

Similarly, $\Theta_{C,v,t,i}$ is the standard cocycle $\phi_{[c'']}$, where $E(c'') = E(\mathcal{R}_{C,v}) \cup \{e'_i\}$, and

$$f_{[c'']}(C') = \begin{cases} f_{C,v}(C') & \text{if } C' \neq C_0, \\ f_{C,v}(C') - 1 & \text{if } C' = C_0. \end{cases}$$

Here C_0 is the cloud in direction 0 from v in the diagram for $\phi_{[c'']}$. Figure 8 gives the sum $\sum_{i=1}^{d(v)-1} \Theta_{C,v,t,i}$ with C as in Figure 6.

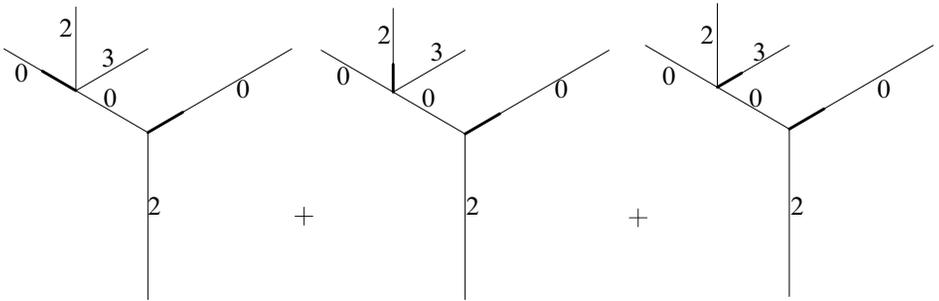


Figure 8: The sum $\sum_{i=1}^{d(v)-1} \Theta_{C,v,t,i}$ with C as in Figure 6

Denote by $\widehat{\delta\mathcal{R}_{C,v}}$ the unique cochain in $C^*(\widehat{UD^n(\Gamma)})$ which maps to $\delta\mathcal{R}_{C,v}$. If $[\omega]$ is an equivalence class of k -cells which does not contain a critical cell, then, by the classification of critical cells, there must be some edge e such that e is order-respecting in every cell in $[\omega]$. In this case, call e a *bad edge* of $[\omega]$.

Theorem 2.9 [6] *Let $J \subset H^*(\widehat{\text{UD}}^n(\Gamma); \mathbb{Z})$ denote the ideal generated by the classes $\delta\mathcal{R}_{\mathcal{C},v}$ where \mathcal{C} is a cloud diagram containing at least one bad edge e and $v = \iota(e)$. Then we have*

$$H^*(\text{UD}^n(\Gamma); \mathbb{Z}) \cong H^*(\widehat{\text{UD}}^n(\Gamma); \mathbb{Z})/J.$$

In the proof of Lemma 3.6, we will be interested in writing the cohomology class of a standard cocycle $\phi_{[\omega]}$ as a linear combination of basis elements (ie cohomology classes of critical cocycles), and comparing $\phi_{[\omega]}$ with these basis elements. If $[\omega]$ contains a critical cell, then $\phi_{[\omega]}$ itself represents a basis element. If $[\omega]$ does not contain a critical cell, the coboundaries $\delta\mathcal{R}_{\mathcal{C},v}$ give a way to rewrite the cohomology class of $\phi_{[\omega]}$ in terms of critical cocycles as follows: Let \mathcal{C} be the cloud diagram for $\phi_{[\omega]}$. The class $[\omega]$ necessarily contains at least one bad edge e which falls in direction d_0 from some vertex v . We will again assume here that $\tau(e)$ has degree less than 3 (the case in which $\text{deg}(\tau(e)) \geq 3$ is addressed in [6], but will not be needed here). Since e is a bad edge, it must be the case that for all $0 < i < d_0$, we have $f_{[\omega]}(C_i) = f_{\mathcal{C},v}(C_i) = 0$, so the first nonzero term in $\sum \Theta_{\mathcal{C},v,\tau,i}$ is $\phi_{[\omega]}$, and any other nonzero term (if there are any) is of the form $\phi_{[\omega']}$, where $[\omega']$ contains an edge in direction $j > d_0$ from v . Any nonzero term in $\sum \Theta_{\mathcal{C},v,\iota,i}$ (if there are any) is of the form $\phi_{[\omega']}$, where $[\omega']$ is an equivalence class of k -cells with the property that if $\{[\omega_1], \dots, [\omega_k]\}$ and $\{[\omega'_1], \dots, [\omega'_k]\}$ are the unique collections of equivalence classes of 1-cells which have $[\omega]$ and $[\omega']$ as their respective least upper bounds, then there are some i and j such that $[\omega_i]$ contains the edge e and $[\omega'_j]$ contains an edge with initial point at $v = \iota(e)$, but $f_{[\omega'_j]}(C_0) < f_{[\omega_i]}(C_0)$. Here, C_0 is the cloud in direction 0 from v in the diagrams for $[\omega_i]$ and $[\omega'_j]$.

Therefore, on the level of cohomology, we have $\phi_{[\omega]} = -A + B$, where A is a sum of standard cocycles $\phi_{[\omega']}$ such that $[\omega']$ has an edge e' whose initial point is v , but e' falls in a direction from v greater than that in which e falls, and B is a sum of standard cocycles $\phi_{[\omega']}$ as above. It is possible that some of the terms in A or B are again standard cocycles corresponding to classes which do not contain critical cells (such as the three standard cocycles described in Figure 8), but if this is the case, we may rewrite each such cocycle using the procedure above. Farley shows that this process eventually terminates and, after repeatedly applying the procedure, we may write $\phi_{[\omega]} = \Sigma$, where Σ is a linear combination of critical cocycles. Suppose $\phi_{[\tilde{\omega}]}$ is a critical cocycle which appears in Σ . Let $\{[\omega_1], \dots, [\omega_k]\}$ and $\{[\tilde{\omega}_1], \dots, [\tilde{\omega}_k]\}$ be the unique collections of equivalence classes of 1-cells which have $[\omega]$ and $[\tilde{\omega}]$ as their respective least upper bounds. We wish to compare the classes in each collection.

First, note that the cloud diagram for any term appearing in $\delta(\mathcal{R}_{C,v})$ must contain an edge whose initial point is $\iota(e) = v$ and every edge $f \neq e$ in $[\omega]$ is also an edge in the cloud diagram for each term in $\delta(\mathcal{R}_{C,v})$. Therefore, for any edge e' in $[\omega]$, the cloud diagram of any term in Σ must contain an edge whose initial point is $\iota(e')$. In other words, for each i , there is some s such that if e' is the unique edge in $[\omega_i]$ and e'' is the unique edge in $[\tilde{\omega}_s]$, then $\iota(e') = \iota(e'')$. Let C_0 be the cloud in direction 0 from v in the cloud diagrams for $[\omega_i]$ and $[\tilde{\omega}_s]$. At no point in the rewriting process do we add more vertices to the cloud C_0 . If e' is a bad edge in $[\omega]$, then either (i) $f_{[\omega_i]}(C_0) = f_{[\tilde{\omega}_s]}(C_0)$ and $\tau(e'') > \tau(e')$, or (ii) $f_{[\omega_i]}(C_0) > f_{[\tilde{\omega}_s]}(C_0)$. If e' is not a bad edge in $[\omega]$, then it may or may not become a bad edge at some stage of the rewriting process. If e' never becomes a bad edge, then we must have $e'' = e'$ and $f_{[\omega_i]}(C_0) = f_{[\tilde{\omega}_s]}(C_0)$. If e' does become bad, then, as above, either (i) $f_{[\omega_i]}(C_0) = f_{[\tilde{\omega}_s]}(C_0)$ and $\tau(e'') > \tau(e')$, or (ii) $f_{[\omega_i]}(C_0) > f_{[\tilde{\omega}_s]}(C_0)$.

Note that if D_0 is the cloud in the 0-direction from v in the diagram for $[\omega]$ (so D_0 is contained in C_0), then it is possible that vertices are added to D_0 in the rewriting process if some edge with an initial point on the geodesic from $\iota(e')$ to $*$ becomes bad at some stage, but for each vertex added to D_0 , there must be some other cloud D_1 contained in C_0 which loses a vertex. The observations in the preceding paragraph are similar to Farley's notion of the *rank* of a cell c defined in [6].

The coboundaries $\delta\mathcal{R}_{C,v}$ illustrate the complicated nature of the cohomology ring $H^*(UD^n(\Gamma))$. The delicacy of the ring structure is studied further in [12], for example, but the above is sufficient for what follows.

3 Motion planning of configuration spaces of trees

Before proving [Theorem 1.2](#), we first mention some of the tools for determining the topological complexity of any space.

Theorem 3.1 [3] *TC(X) is homotopy-invariant. That is, if X and Y are homotopic, then $TC(X) = TC(Y)$.*

Since we assume Γ is sufficiently subdivided, the spaces $C^n(\Gamma)$ and $UC^n(\Gamma)$ are homotopic to $D^n(\Gamma)$ and $UD^n(\Gamma)$, respectively, by [Theorem 2.1](#), so [Theorem 3.1](#) allows us to work with $D^n(\Gamma)$ and $UD^n(\Gamma)$ to determine $TC(C^n(\Gamma))$ and $TC(UC^n(\Gamma))$. The next theorem gives an upper bound for $TC(X)$ based on the dimension of X :

Theorem 3.2 [3] *Let X be any CW complex. Then we have the upper bound $\text{TC}(X) \leq 2 \cdot \dim(X) + 1$.*

Finally, there is a cohomological lower bound for $\text{TC}(X)$. Before stating the theorem, we introduce some definitions. Let k be a field, and consider the cup product

$$\smile: H^*(X; k) \otimes H^*(X; k) \rightarrow H^*(X; k).$$

Let $Z(X) \subset H^*(X; k) \otimes H^*(X; k)$ be the kernel of this homomorphism, called the ideal of *zero-divisors* of $H^*(X; k)$. The tensor product $H^*(X; k) \otimes H^*(X; k)$ has a multiplication given by $(\alpha \otimes \beta)(\alpha' \otimes \beta') = (-1)^{|\alpha'| \cdot |\beta|} \alpha \alpha' \otimes \beta \beta'$, where $|x| = j$ if $x \in H^j(X; k)$. The *zero-divisors-cup-length* of $H^*(X; k)$ is the largest i such there are elements $a_1, \dots, a_i \in Z(X)$ with $a_1 \cdots a_i \neq 0$.

Theorem 3.3 [3] *$\text{TC}(X)$ is greater than the zero-divisors-cup-length of $H^*(X; k)$.*

Now we establish the upper bounds in [Theorem 1.2](#):

Lemma 3.4 *Let $k = \min\{\lfloor \frac{n}{2} \rfloor, m(\Gamma)\}$. Then $\text{TC}(\text{UD}^n(\Gamma)) \leq 2k + 1$.*

Proof This is immediate from the second statement of [Theorem 2.5](#) and [Theorems 3.1](#) and [3.2](#), but we will give a (fairly) explicit motion planning algorithm which realizes this upper bound. This algorithm is similar to the one given by Farber in [\[4\]](#). Let a_1, \dots, a_n be the first n vertices in a sufficiently subdivided tree Γ (so $a_1 = *$). The fact that Γ is sufficiently subdivided implies that only possibly a_n is essential. Consider a point $\mathbf{x} = \{x_1, \dots, x_n\} \in \text{UD}^n(\Gamma)$. If x_i falls on a vertex v , let $f(x_i)$ be the number assigned to v in the ordering above. If x_i falls on the interior of an edge e , let $f(x_i)$ be the number assigned to $\iota(e)$. Since the order in which the x_i appear in \mathbf{x} is irrelevant, we may assume without loss of generality that $f(x_i) < f(x_{i+1})$ for all i . Define a map $\sigma_{\mathbf{x}}: [0, 1] \rightarrow \text{UD}^n(\Gamma)$ as follows: during the interval $[\frac{i-1}{n}, \frac{i}{n}]$, $\sigma_{\mathbf{x}}$ is the path which moves x_i along the geodesic to a_i at constant speed, and keeps all other x_j fixed. The choice of ordering of the vertices makes this a valid path in $\text{UD}^n(\Gamma)$ (ie there is at least a full open edge between any two components of $\sigma_{\mathbf{x}}(t)$ at any given time t). Each $\sigma_{\mathbf{x}}$ is clearly continuous. Define the section $s: \text{UD}^n(\Gamma) \times \text{UD}^n(\Gamma) \rightarrow P(\text{UD}^n(\Gamma))$ by $s(\mathbf{x}, \mathbf{y}) = \sigma_{\mathbf{x}} \bar{\sigma}_{\mathbf{y}}$, the path $\sigma_{\mathbf{x}}$ followed by the reverse of $\sigma_{\mathbf{y}}$.

This is not continuous on $\text{UD}^n(\Gamma) \times \text{UD}^n(\Gamma)$. If some x_i (or y_i) falls on an endpoint $\tau(e)$ of some edge e , a slight perturbation of x_i (or y_i) may cause it to fall on the

interior of e , which will alter the numbering of the elements in \mathbf{x} (or \mathbf{y}), which can lead to a very different path if $\iota(e)$ is essential. So, we wish to examine the sets on which s is continuous. For a collection E of edges, let S_E be the set of points $\mathbf{x} = \{x_1, \dots, x_n\} \in \text{UD}^n(\Gamma)$ with the property that the interior of each edge $e \in E$ contains (exactly) one x_i and no x_j falls on the interior of any edge not in E (so such an x_j falls on a vertex which is not the endpoint of any $e \in E$). The function s is continuous on each $S_E \times S_{E'}$. Now let

$$S_i = \bigcup_{|E|=i} S_E.$$

If $|E| = |E'|$ and $E \neq E'$, then a sequence of points in S_E cannot converge to a point in $S_{E'}$, so that $\bar{S}_E \cap S_{E'} = \emptyset$, and similarly $S_E \cap \bar{S}_{E'} = \emptyset$. In other words, S_i is a topologically disjoint union of the sets S_E , and then for each fixed i and j , the set $S_i \times S_j$ is a topologically disjoint union of sets on which s is continuous, so s is continuous on $S_i \times S_j$. Now, a sequence of points in S_i may converge to a point in $S_{i'}$ for some $i' < i$, but no sequence of points in S_i can converge to a point in $S_{i'}$ if $i' > i$, so the sets

$$U_l = \bigcup_{i+j=l} S_i \times S_j$$

are again topologically disjoint unions of sets on which s is continuous, so s is continuous on U_l .

The sets U_0, \dots, U_{2k} cover $(\text{UD}^n(\Gamma))'_k \times (\text{UD}^n(\Gamma))'_k$, since at most k points can fall on the interior of an edge in either factor (see [Theorem 2.5](#)). They are not necessarily open, but each U_l can be replaced with an open set U'_l which allows each x_i which falls on a vertex v (and x_i appears in a point \mathbf{x} in the first component of U_l) to vary slightly away from v (while keeping the point \mathbf{x} in $\text{UD}^n(\Gamma)$), and defining $\sigma_{\mathbf{x},l}$ which is as above, except each of these x_i is given the number for v . This is well-defined, since each v does not fall on any of the l edges whose interiors are occupied by some x_j in \mathbf{x} , so if x_i falls on v , a small perturbation of x_i will not cause it to fall on the interior of any of those l edges. Similar modifications are made in the second component. Define $s'_l: U'_l \rightarrow P(\text{UD}^n(\Gamma))$ by

$$s'_l(\mathbf{x}, \mathbf{y}) = \sigma_{\mathbf{x},l} \bar{\sigma}_{\mathbf{y},l}$$

This is continuous.

If the map $H: [0, 1] \times \text{UD}^n(\Gamma) \rightarrow \text{UD}^n(\Gamma)$ is a deformation retraction from $\text{UD}^n(\Gamma)$ to $(\text{UD}^n(\Gamma))'_k$ with $H(0, -)$ equaling the identity map, then $H_{\mathbf{x}} := H(-, \mathbf{x})$ is a path

from a point $x \in \text{UD}^n(\Gamma)$ to some point in $(\text{UD}^n(\Gamma))'_k$, which varies continuously with x . If $V_l = (H(1, -) \times H(1, -))^{-1}(U'_l)$, then $\{V_0, \dots, V_{2k}\}$ is an open cover of $\text{UD}^n(\Gamma) \times \text{UD}^n(\Gamma)$, and the section $s_l: V_l \rightarrow P(\text{UD}^n(\Gamma))$ given by

$$s_l(x, y) = H_x s'_l(x_1, y_1) \bar{H}_y,$$

where $x_1 = H_x(1)$ and $y_1 = H_y(1)$ is continuous on each V_l for $l = 0, \dots, 2k$, completing the proof. □

An additional step can be added to this algorithm to give the same upper bound for the ordered configuration spaces. Again, this approach is essentially the one described in [4].

Corollary 3.5 *Let Γ be a tree with $m(\Gamma) \geq 1$, and let $k = \min\{\lfloor \frac{n}{2} \rfloor, m(\Gamma)\}$. Then $\text{TC}(D^n(\Gamma)) \leq 2k + 1$.*

Proof Let a_1, \dots, a_n be the first n vertices of Γ , and let $\mathbf{a} = \{a_1, \dots, a_n\} \in \text{UD}^n(\Gamma)$. For $l = 0, \dots, 2k$, let $V_l, \sigma_{x,l}$ and H_x be as in the proof of Lemma 3.4. Let $p_1: \text{UD}^n(\Gamma) \times \text{UD}^n(\Gamma) \rightarrow \text{UD}^n(\Gamma)$ be the projection of the first component, and let $V_l^1 = p_1(V_l)$, so that if $x \in V_l^1$, then $H_x \sigma_{x_1,l}$ (with $x_1 = H_x(1)$) is a path from x to \mathbf{a} , and this path varies continuously as x varies in V_l^1 . If π is the covering space projection $\pi: D^n(\Gamma) \rightarrow \text{UD}^n(\Gamma)$, and $\tilde{x} \in \pi^{-1}(x) \subset D^n(\Gamma)$, with $x \in V_l^1$, let $\tilde{\sigma}_{\tilde{x},l}$ be the unique lift of $H_x \sigma_{x_1,l}$ which satisfies $\tilde{\sigma}_{\tilde{x},l}(0) = \tilde{x}$. This is a path from \tilde{x} to some point $\tilde{\mathbf{a}} \in \pi^{-1}(\mathbf{a})$ which varies continuously as \tilde{x} varies in $\tilde{V}_l^1 = \pi^{-1}(V_l^1)$. Define $\tilde{\sigma}_{\tilde{y},l}$ similarly if $\tilde{y} \in \pi^{-1}(y)$ for some y in the second component of V_l .

Now, given $\tilde{\mathbf{a}}, \tilde{\mathbf{a}}' \in \pi^{-1}(\mathbf{a})$, let $\rho_{\tilde{\mathbf{a}}, \tilde{\mathbf{a}}'}$ be any path from $\tilde{\mathbf{a}}$ to $\tilde{\mathbf{a}}'$. Such a path exists by Theorem 2.2 since Γ has at least one essential vertex. The function

$$r: \pi^{-1}(\mathbf{a}) \times \pi^{-1}(\mathbf{a}) \rightarrow P(D^n(\Gamma))$$

given by $r(\tilde{\mathbf{a}}, \tilde{\mathbf{a}}') = \rho_{\tilde{\mathbf{a}}, \tilde{\mathbf{a}}'}$ is continuous since the domain is a discrete space. For each $l = 0, \dots, 2k$, let $\tilde{V}_l = (\pi \times \pi)^{-1}(V_l)$, and define the section $\tilde{s}_l: \tilde{V}_l \rightarrow P(D^n(\Gamma))$ by

$$\tilde{s}_l(\tilde{x}, \tilde{y}) = \tilde{\sigma}_{\tilde{x},l} r(\tilde{\sigma}_{\tilde{x},l}(1), \tilde{\sigma}_{\tilde{y},l}(1)) \tilde{\sigma}_{\tilde{y},l}.$$

This is continuous on \tilde{V}_l , and $\{\tilde{V}_0, \dots, \tilde{V}_{2k}\}$ is an open cover of $D^n(\Gamma) \times D^n(\Gamma)$, completing the proof. □

In establishing the lower bound, the following will be useful:

Lemma 3.6 Suppose $m(\Gamma) \geq 1$, $UD^n(\Gamma)$ has the homotopy type of a k -dimensional CW complex, and Φ and Ψ are critical k -cells of $UD^n(\Gamma)$. If $\{[c_1], \dots, [c_k]\}$ and $\{[d_1], \dots, [d_k]\}$ are the unique collections of equivalence classes of 1-cells having least upper bounds $[\Phi]$ and $[\Psi]$, respectively, and $[c_i] \neq [d_j]$ for all i and j , then $TC(UD^n(\Gamma))$ and $TC(D^n(\Gamma))$ are greater than $2k$.

Proof If necessary, rearrange the equivalence classes in $\{[c_1], \dots, [c_k]\}$ so that $\iota(e_i) < \iota(e_{i+1})$, as above, and arrange the classes in $\{[d_1], \dots, [d_k]\}$ similarly. Consider the zero divisors

$$\bar{\phi}_{[c_i]} = \phi_{[c_i]} \otimes 1 - 1 \otimes \phi_{[c_i]}, \quad \bar{\phi}_{[d_j]} = \phi_{[d_j]} \otimes 1 - 1 \otimes \phi_{[d_j]}$$

in $Z(UD^n(\Gamma)) \subset H^*(UD^n(\Gamma); \mathbb{Q}) \otimes H^*(UD^n(\Gamma); \mathbb{Q})$ and their product

$$\begin{aligned} (3-1) \quad & \left(\prod_{i=1}^k \bar{\phi}_{[c_i]} \right) \left(\prod_{j=1}^k \bar{\phi}_{[d_j]} \right) \\ &= \pm (\phi_{[c_1]} \cdots \phi_{[c_k]} \otimes \phi_{[d_1]} \cdots \phi_{[d_k]}) \pm (\phi_{[d_1]} \cdots \phi_{[d_k]} \otimes \phi_{[c_1]} \cdots \phi_{[c_k]}) \\ & \qquad \qquad \qquad + \text{other terms} \\ &= \pm \phi_{[\Phi]} \otimes \phi_{[\Psi]} \pm \phi_{[\Psi]} \otimes \phi_{[\Phi]} + \text{other terms.} \end{aligned}$$

Since $UD^n(\Gamma)$ has the homotopy type of a k -dimensional complex, all products of more than k 1-dimensional classes in $H^*(UD^n(\Gamma))$ are zero, so any nonzero term in *other terms* must be of the form $\alpha \otimes \beta$ where α and β are both degree- k monomials in $\phi_{[c_1]}, \dots, \phi_{[c_k]}, \phi_{[d_1]}, \dots, \phi_{[d_k]}$. For a collection $I = \{[c_{i_1}], \dots, [c_{i_p}]\}$, we denote $\phi_{[c_{i_1}]} \cdots \phi_{[c_{i_p}]}$ by ϕ_I and denote the set $\{[c_i] : [c_i] \notin I\}$ by \bar{I} , and use analogous definitions for a collection $J = \{[d_{j_1}], \dots, [d_{j_r}]\}$. Then we can write $\alpha \otimes \beta = \phi_I \phi_J \otimes \phi_{\bar{I}} \phi_{\bar{J}}$ for some collections I and J , with $I \neq \emptyset$ and $J \neq \emptyset$ (so that neither $I \cup J$ nor $\bar{I} \cup \bar{J}$ is either $\{[c_1], \dots, [c_k]\}$ or $\{[d_1], \dots, [d_k]\}$, by assumption), and $|I| + |J| = k$.

If either $I \cup J$ or $\bar{I} \cup \bar{J}$ does not have an upper bound, then $\alpha \otimes \beta = 0$. If both $I \cup J$ and $\bar{I} \cup \bar{J}$ have least upper bounds $[\omega]$ and $[\omega']$, respectively, which both contain critical cells, then $\alpha \otimes \beta = \phi_{[\omega]} \otimes \phi_{[\omega']}$ is a basis element. Since the collections of equivalence classes of 1-cells which have $[\Phi]$ and $[\Psi]$ as their least upper bounds are unique, and $I \cup J$ is neither $\{[c_1], \dots, [c_k]\}$ nor $\{[d_1], \dots, [d_k]\}$, we have $[\omega] \notin \{[\Phi], [\Psi]\}$ and, similarly, $[\omega'] \notin \{[\Phi], [\Psi]\}$, so $\alpha \otimes \beta \notin \{\phi_{[\Phi]} \otimes \phi_{[\Psi]}, \phi_{[\Psi]} \otimes \phi_{[\Phi]}\}$.

If both $I \cup J$ and $\bar{I} \cup \bar{J}$ have upper least bounds $[\omega]$ and $[\omega']$, but exactly one, say $[\omega]$, contains a critical cell, then $\alpha = \phi_{[\omega]}$, a basis element which is equal to neither

$\phi_{[\Phi]}$ nor $\phi_{[\Psi]}$, as above. We use the procedure following [Theorem 2.9](#) to rewrite the cocycle $\phi_{[\omega']}$ in terms of critical cocycles, arriving at $\alpha \otimes \beta = \phi_{[\omega]} \otimes \Sigma'$, where Σ' is a linear combination of critical cocycles. Note that since each edge e in $[\omega']$ is an order-disrespecting edge in either Φ or Ψ , the endpoint $\iota(e)$ must be essential, so this does not violate our assumption that $\tau(e)$ has degree less than 3 (we may need to further subdivide Γ if $n = 2$). If Σ' is nonzero, then $\alpha \otimes \beta$ can be written as a linear combination of basis elements of the form $\phi_{[\omega]} \otimes \phi_{[\tilde{\omega}]}$, none of which are $\phi_{[\Phi]} \otimes \phi_{[\Psi]}$ or $\phi_{[\Psi]} \otimes \phi_{[\Phi]}$. Similar statements hold if $[\omega']$ contains a critical cell.

Finally, suppose both $I \cup J$ and $\bar{I} \cup \bar{J}$ have upper least bounds $[\omega]$ and $[\omega']$, but neither contains a critical cell. We may write $\phi_{[\omega]}$ and $\phi_{[\omega']}$ as linear combinations of critical cocycles, as above:

$$\phi_{[\omega]} = \Sigma, \quad \phi_{[\omega']} = \Sigma'.$$

Since $[\omega]$ does not contain a critical cell, it must contain some bad edge e . There is an equivalence class in $I \cup J$ which contains the edge e , and this class is unique since $I \cup J$ has an upper bound (see [Lemma 2.6](#)). Suppose first that $[c_i] \in I$ is this class. Consider a critical cocycle $\phi_{[\tilde{\omega}]}$ appearing in Σ , and let $\{[\tilde{\omega}_1], \dots, [\tilde{\omega}_k]\}$ be the unique collection of equivalence classes of 1-cells which has $[\tilde{\omega}]$ as its least upper bound. As in the discussion following [Theorem 2.9](#), there must be some s such that the unique edge e' in $[\tilde{\omega}_s]$ satisfies $\iota(e') = \iota(e)$, and either

$$(3-2) \quad f_{[c_i]}(C_0) = f_{[\tilde{\omega}_s]}(C_0) \quad \text{and} \quad \tau(e) < \tau(e'), \quad \text{or} \quad f_{[c_i]}(C_0) > f_{[\tilde{\omega}_s]}(C_0).$$

But, in either case, we have $[c_i] \neq [\tilde{\omega}_s]$. Since $[c_i]$ is the only class in $\{[c_1], \dots, [c_k]\}$ which contains an edge whose initial point is $\iota(e)$, we have

$$\{[\tilde{\omega}_1], \dots, [\tilde{\omega}_k]\} \neq \{[c_1], \dots, [c_k]\},$$

so that, as above, $[\tilde{\omega}] \neq [\Phi]$. Therefore the cocycle $\phi_{[\Phi]}$ does not appear in Σ (by this, we mean that $\phi_{[\Phi]}$ does not appear with a nonzero coefficient in Σ).

Now, it is possible that $\phi_{[\tilde{\omega}]} = \phi_{[\Psi]}$. If this is the case, we have

$$\{[\tilde{\omega}_1], \dots, [\tilde{\omega}_k]\} = \{[d_1], \dots, [d_k]\},$$

so there must be some j such that $[d_j] = [\tilde{\omega}_s]$. Then, from (3-2), we have either

$$(3-3) \quad f_{[c_i]}(C_0) = f_{[d_j]}(C_0) \quad \text{and} \quad \tau(e) < \tau(e'), \quad \text{or} \quad f_{[c_i]}(C_0) > f_{[d_j]}(C_0).$$

We wish to show that in this case, $\phi_{[\Phi]}$ cannot appear in Σ' . The edge e is in the class $[c_i]$ and the edge e' is in the class $[d_j] = [\tilde{\omega}_s]$, and since $[c_i] \in I \cup J$ and $I \cup J$ contains an upper bound, it must be the case that $[d_j] \notin J$. So we have $[d_j] \in \bar{J}$, and therefore the edge e' appears in $[\omega']$. Let $\phi_{[\tilde{\omega}']}$ be a critical cocycle appearing in Σ' , and let $\{[\tilde{\omega}'_1], \dots, [\tilde{\omega}'_k]\}$ be the unique collection of equivalence classes of 1-cells which has $[\tilde{\omega}']$ as its least upper bound. Because the edge e' appears in $[\omega']$, there must be some t such that if e'' is the unique edge in $[\tilde{\omega}'_t]$, then $\iota(e'') = \iota(e')$. Since the edge e' may or may not be bad in $[\omega']$, we have either

$$(3-4) \quad f_{[d_j]}(C_0) = f_{[\tilde{\omega}'_t]}(C_0) \quad \text{and} \quad \tau(e') \leq \tau(e''), \quad \text{or} \quad f_{[d_j]}(C_0) > f_{[\tilde{\omega}'_t]}(C_0).$$

For the sake of contradiction, suppose that $\phi_{[\tilde{\omega}']} = \phi_{[\Phi]}$. Then we have

$$\{[\tilde{\omega}'_1], \dots, [\tilde{\omega}'_k]\} = \{[c_1], \dots, [c_k]\},$$

so that $[\tilde{\omega}'_t] = [c_{i_0}]$ for some i_0 . In particular, e'' is the unique edge in $[c_{i_0}]$, but since $\iota(e) = \iota(e') = \iota(e'')$, we must have $e'' = e$ and $i_0 = i$. From (3-4), we have either

$$f_{[d_j]}(C_0) = f_{[c_i]}(C_0) \quad \text{and} \quad \tau(e') \leq \tau(e), \quad \text{or} \quad f_{[d_j]}(C_0) > f_{[c_i]}(C_0),$$

but this contradicts (3-3).

So we have shown that if $[c_i]$ is the class in $I \cup J$ which contains the bad edge e , then it is impossible that $\phi_{[\Phi]}$ appears in Σ , and if $\phi_{[\Psi]}$ appears in Σ , it is impossible that $\phi_{[\Phi]}$ appears in Σ' . By a symmetric argument, if $[d_{j'}]$ is the class in $I \cup J$ which contains e , it follows that it is impossible that $\phi_{[\Psi]}$ appears in Σ , and if $\phi_{[\Phi]}$ appears in Σ , it is impossible that $\phi_{[\Psi]}$ appears in Σ' . So, in either case, neither $\phi_{[\Phi]} \otimes \phi_{[\Psi]}$ nor $\phi_{[\Psi]} \otimes \phi_{[\Phi]}$ appears in the expansion of $\Sigma \otimes \Sigma'$ as a linear combination of simple tensors of critical cocycles.

Therefore, the right side of (3-1) may be written as a linear combination of basis elements of the tensor product $H^*(UD^n(\Gamma)) \otimes H^*(UD^n(\Gamma))$, with neither $\phi_{[\Phi]} \otimes \phi_{[\Psi]}$ nor $\phi_{[\Psi]} \otimes \phi_{[\Phi]}$ appearing in *other terms*, and since Φ and Ψ are distinct critical cells,

$$\pm \phi_{[\Phi]} \otimes \phi_{[\Psi]} \pm \phi_{[\Psi]} \otimes \phi_{[\Phi]} \neq 0,$$

so the entire sum is nonzero. So (3-1) is a nonzero product of $2k$ zero-divisors, and Theorem 3.3 establishes that $\text{TC}(UD^n(\Gamma)) \geq 2k + 1$, as desired.

The statement for $\text{TC}(D^n(\Gamma))$ follows from the fact that the map

$$\pi^*: H^*(UD^n(\Gamma)) \rightarrow H^*(D^n(\Gamma))$$

induced by the covering space projection $\pi: D^n(\Gamma) \rightarrow \text{UD}^n(\Gamma)$ is injective (see [10, Proposition 3G.1], which states that any N -sheeted covering space projection given by a group action induces an injection in cohomology with coefficients in a field of characteristic 0). □

Note that the condition $[c_i] \neq [d_j]$ for all i and j is necessary for (3-1) to be nonzero. If $[c_i] = [d_j]$ for some i and j , then $\phi_{[c_i]} = \phi_{[d_j]}$, so (3-1) contains the product

$$\begin{aligned} \bar{\phi}_{[c_i]}\bar{\phi}_{[c_i]} &= (\phi_{[c_i]} \otimes 1 - 1 \otimes \phi_{[c_i]})(\phi_{[c_i]} \otimes 1 - 1 \otimes \phi_{[c_i]}) \\ &= (\phi_{[c_i]} \otimes 1)^2 - (\phi_{[c_i]} \otimes 1)(1 \otimes \phi_{[c_i]}) - (1 \otimes \phi_{[c_i]})(\phi_{[c_i]} \otimes 1) + (1 \otimes \phi_{[c_i]})^2 \\ &= \phi_{[c_i]}^2 \otimes 1 - (-1)^0(\phi_{[c_i]} \otimes \phi_{[c_i]}) - (-1)^1(\phi_{[c_i]} \otimes \phi_{[c_i]}) + 1 \otimes \phi_{[c_i]}^2 \\ &= 0. \end{aligned}$$

Therefore, (3-1) is nonzero if and only if $[c_i] \neq [d_j]$ for any i and j .

We now prove [Theorem 1.2](#).

Proof of Theorem 1.2 Assume without loss of generality that Γ is sufficiently subdivided for n , so that the topological and discrete configuration spaces have the same homotopy type, and hence, the same topological complexity. We will prove the results for the discrete configuration spaces. The upper bounds, $\text{TC}(\text{UD}^n(\Gamma)) \leq 2m + 1$ and $\text{TC}(D^n(\Gamma)) \leq 2m + 1$ in statement (1) and $\text{TC}(\text{UD}^n(\Gamma)) \leq 2q + 1$ and $\text{TC}(D^n(\Gamma)) \leq 2q + 1$ in statement (2), are given in [Lemma 3.4](#) and [Corollary 3.5](#). To establish the lower bounds, we will construct two m -cells for statement (1) and two q -cells for statement (2) that satisfy the conditions of [Lemma 3.6](#).

For statement (1), first note that [Theorem 2.5](#) implies $\text{UD}^n(\Gamma)$ is homotopic to an m -dimensional CW complex. Let v_1, \dots, v_m be (in order) the essential vertices of Γ and let A_1, \dots, A_k be a collection of oriented arcs which is allowable for the set of vertices of degree 3. For each arc A_l , let a_l and b_l be the initial and terminal points of A_l , with respect to the orientation of A_l . The endpoints need not be vertices of Γ .

We will construct an m -cell Φ as follows. At each essential vertex v_i , let e_i be the edge in direction 2 from v_i (so that $\iota(e_i) = v_i$), let f_i be the edge in direction 1 from v_i , and let $u_i = \tau(f_i)$. The labeling of the vertices forces $v_i = \iota(e_i) < u_i < \tau(e_i)$, so that e_i is order-disrespecting in any cell containing u_i . Furthermore, u_i is blocked in any cell containing e_i . Add each edge e_i and each vertex u_i to Φ . The edges e_i determine a system of clouds of Γ .

Now, if $n > 2m$ (which must be the case, by assumption, if there are any vertices of degree 3), we must have $n \geq 3$, so since Γ is sufficiently subdivided, if the endpoint a_l falls on an edge e_i , we may shrink or enlarge A_l slightly so that the new initial point falls just beyond $\tau(e_i)$, without changing the fact that the collection $\{A_i\}_{i=1}^k$ is allowable for the collection of degree-3 vertices. So, we may assume that each endpoint a_l falls in one of the clouds of Γ determined by the edges e_i . Then, inductively, for $l = 1, \dots, k$, let w_l be the minimal vertex in the cloud containing a_l which we have not already included in Φ , and add this vertex to Φ . Then w_1, \dots, w_k are all blocked in Φ . Finally, let w_{k+1}, \dots, w_{n-2m} be the first $n - 2m - k$ vertices of Γ which are not already in Φ , and add them to Φ (so they are all blocked).

So we have

$$\Phi = \{e_1, \dots, e_m, u_1, \dots, u_m, w_1, \dots, w_{n-2m}\},$$

where the vertices w_i do not appear if $n = 2m$. The cell Φ is critical. Now we will construct a critical m -cell Ψ similarly. If v_i is a vertex of degree 3, let $e'_i = e_i$ and $u'_i = u_i$. If v_i is a vertex of degree greater than 3, let e'_i be the edge in direction 3 from v_i , let f'_i be the edge in direction 2, and let $u'_i = \tau(f'_i)$, so u'_i is blocked by e'_i . As above, we have $v'_i = \iota(e'_i) < u'_i < \tau(e'_i)$, so e'_i is order-disrespecting in any cell containing u'_i . Add the edges e'_i and vertices u'_i to the cell Ψ . Similar to the argument above, if $n > 2m$, we may assume each endpoint b_l falls in one of the clouds of Γ determined by the edges e'_i . Inductively for $l = 1, \dots, k$, let w'_l be the minimal vertex in the cloud containing b_l which we have not already included in Ψ , and add this vertex to Ψ . Let $w'_{k+1}, \dots, w'_{n-2m}$, be the first $n - 2m - k$ vertices of Γ which are not already in Ψ , and add them in, so now we have

$$\Psi = \{e'_1, \dots, e'_m, u'_1, \dots, u'_m, w'_1, \dots, w'_{n-2m}\},$$

where the vertices w'_i do not appear if $n = 2m$. The cell Ψ is critical. An example is given in [Figure 9](#) for $m = 11$, $k = 3$, and $n = 27$. The figure on the left is a tree Γ with 3 oriented arcs which are allowable for the collection of degree-3 vertices. The orientations of the arcs are indicated with arrows.

Now let $\{[c_1], \dots, [c_m]\}$ be the unique collection of equivalence classes of 1-cells having $[\Phi]$ as its least upper bound. The equivalence class $[c_i]$ can be represented using a cloud diagram having a single edge (the edge e_i) in direction 2 from the essential vertex v_i , whose degree is t_i , and t_i clouds. Likewise, let $\{[d_1], \dots, [d_m]\}$ be the unique collection of equivalence classes of 1-cells having $[\Psi]$ as its least upper bound.

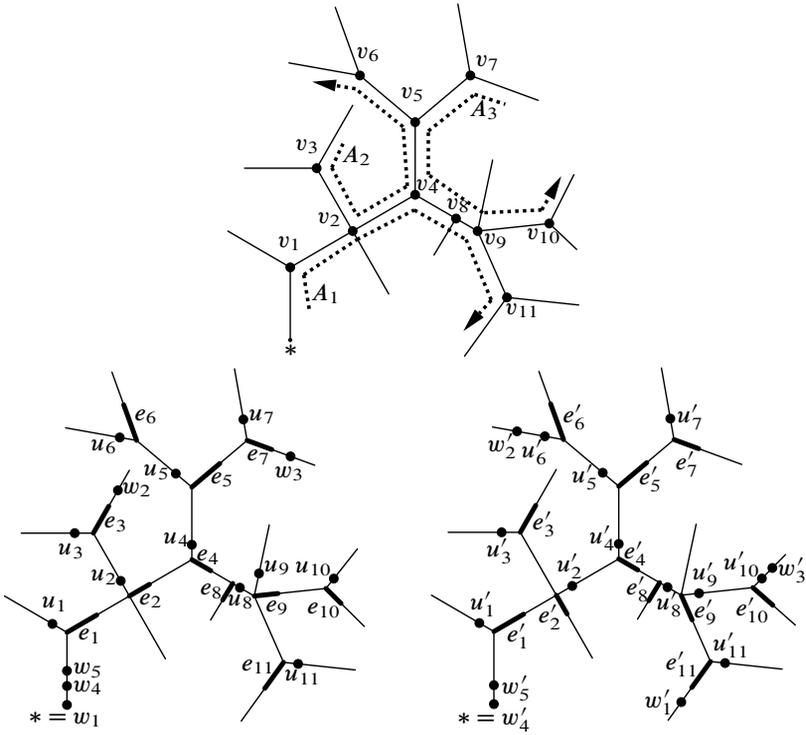


Figure 9: A graph Γ (top) and the critical cells Φ and Ψ in $UD^{27}(\Gamma)$ (bottom left and right)

The equivalence class $[d_i]$ can be represented using a cloud diagram having a single edge (the edge e'_i) and t_i clouds. If $t_i \geq 4$, then $[c_i] \neq [d_i]$, since the edges e_i and e'_i differ.

If $t_i = 3$, then the cloud diagram for $[c_i]$ will contain the single edge e_i and three clouds C_0, C_1 and C_2 , in the 0-, 1- and 2-directions from v_i . The cloud diagram for $[d_i]$ will have the same edge and clouds. Suppose $f(C)$ is the value assigned to cloud C in the diagram for $[c_i]$ and $g(C)$ is the value assigned to cloud C in the diagram for $[d_i]$. For $\delta = 0, 1, 2$, let m_δ be the number of essential vertices in C_δ , and let $W_\delta = |\{p : w_p \in C_\delta\}|$ and $W'_\delta = |\{p : w'_p \in C_\delta\}|$. Then $f(C_\delta) = 2m_\delta + W_\delta + \lambda_\delta$ and $g(C_\delta) = 2m_\delta + W'_\delta + \lambda_\delta$, where $\lambda_0 = \lambda_2 = 0$ and $\lambda_1 = 1$.

Suppose the vertex v_i falls on the interior of an arc A_l (by assumption there is at least one such arc). For each direction δ , if the arc A_l is oriented towards v_i in the direction δ , then a_l must fall in C_δ , and b_l must fall in a different cloud, so the vertex w_l contributes 1 to W_δ , but w'_l contributes 0 to W'_δ . On the other hand, if A_l is

oriented away from v_i in the direction δ , then b_l must fall in C_δ , and a_l must fall in a different cloud so the vertex w'_l contributes 1 to W'_δ , but w_l contributes 0 to W_δ . If v_i does not fall on A_l , then both w_l and w'_l must be in the same cloud, so either w_l and w'_l contribute 1 to W_δ and W'_δ , respectively, or each contributes 0 to W_δ and W'_δ , respectively. Finally, note that each w_i and w'_i for $i > k$ must fall in the same cloud as the basepoint $*$ in both Φ and Ψ since Γ is sufficiently subdivided, so in this case w_i and w'_i contribute 1 to W_0 and W'_0 , respectively.

In other words, the difference $W_\delta - W'_\delta$ is equal to the number of arcs oriented towards v_i in the direction δ minus the number of arcs oriented away from v_i in the direction δ . This is the number $\eta_\delta(v_i)$ in Definition 2.3, so by assumption, there must be at least one direction δ such that $W_\delta \neq W'_\delta$. Furthermore, by the remarks following the same definition, there must be directions δ and δ' such that $W'_\delta - W_\delta > 0$ and $W'_{\delta'} - W_{\delta'} < 0$, and therefore, $g(C_\delta) > f(C_\delta)$ and $g(C_{\delta'}) < f(C_{\delta'})$. In this case, call C_δ and $C_{\delta'}$ positive and negative clouds, respectively, to reflect that $[d_i]$ has more vertices than $[c_i]$ in direction δ (from v_i) and less vertices in direction δ' . It is possible that either all 3 clouds are categorized as either positive or negative or that one cloud remains uncategorized. But, in either case, we have $[c_i] \neq [d_i]$. As an example, Figure 10 shows the classes $[c_4]$ and $[d_4]$ with Γ as above.

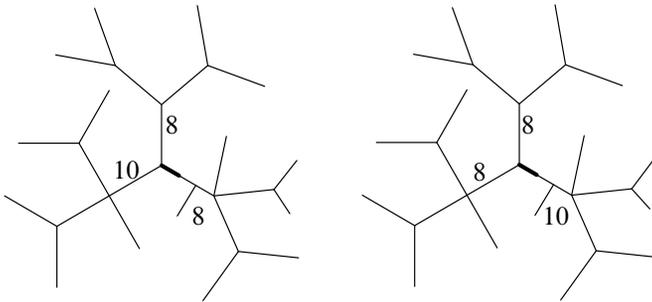


Figure 10: Cloud representations for $[c_4]$ (left) and $[d_4]$ (right) corresponding to the cells Φ and Ψ in $UD^{27}(\Gamma)$

Furthermore, it follows from this description that $[c_i] \neq [d_j]$ for any i and j . Indeed, if $[c_i] = [d_j]$, then both classes must contain a common edge. But, this can only happen if $i = j$ and $t_i = 3$, but we just saw $[c_i] \neq [d_i]$ in this case. So the cells Φ and Ψ satisfy the hypotheses of Lemma 3.6, proving statement (1).

The construction for statement (2) is similar. Here, by Theorem 2.5, $UD^n(\Gamma)$ is homotopic to a complex of dimension q . Let v_1, \dots, v_r be the vertices of degree

greater than 3, and let $\tilde{v}_1, \dots, \tilde{v}_s$ be the vertices of degree 3. We consider the two cases in the statement:

Case (a) ($s \geq 2(q-r)$) Let $R = \min\{r, q\}$, and for $i = 1, \dots, R$, let e_i, u_i, e'_i , and u'_i be as in the proof of statement (1) (so that e_i and e'_i are order-disrespecting edges in directions 2 and 3 from v_i , and u_i and u'_i are blocked vertices in directions 1 and 2). If $r < q$, for $i = 1, \dots, 2(q-r)$, let \tilde{e}_i be the edge in direction 2 from \tilde{v}_i , and let \tilde{u}_i be the vertex on the edge in direction 1 from \tilde{v}_i which forces \tilde{e}_i to be order-disrespecting. If $\epsilon = 0$, let

$$\Phi = \{e_1, \dots, e_R, u_1, \dots, u_R, \tilde{e}_1, \dots, \tilde{e}_{q-r}, \tilde{u}_1, \dots, \tilde{u}_{q-r}\},$$

$$\Psi = \{e'_1, \dots, e'_R, u'_1, \dots, u'_R, \tilde{e}_{q-r+1}, \dots, \tilde{e}_{2(q-r)}, \tilde{u}_{q-r+1}, \dots, \tilde{u}_{2(q-r)}\},$$

where the edges \tilde{e}_i and vertices \tilde{u}_i do not appear if $r \geq q$. If $\epsilon = 1$, add the vertex $*$ to each cell. An example is in Figure 11 for Γ as above with $q = 5$ and $\epsilon = 0$.

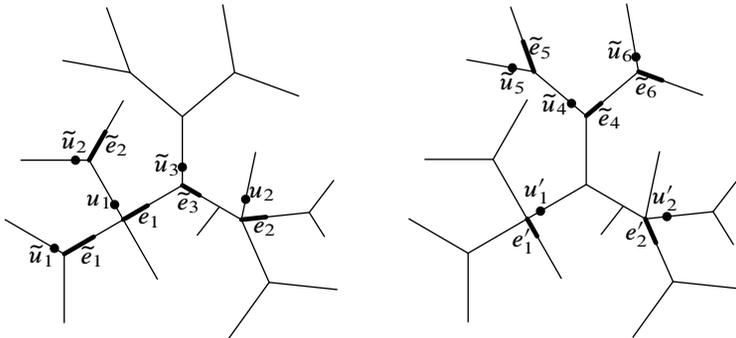


Figure 11: The critical cells Φ and Ψ in $UD^{10}(\Gamma)$

For the same reasons as in the first part of the proof, the cells Φ and Ψ are critical, and if $\{[c_1], \dots, [c_q]\}$ and $\{[d_1], \dots, [d_q]\}$ are the collections of equivalence classes of 1-cells having $[\Phi]$ and $[\Psi]$ as their least upper bounds, here it is clear that $[c_i] \neq [d_j]$ for any i and j , since no edge in Φ is in Ψ .

Case (b) ($s < 2(q-r)$) First, consider the $\epsilon = 0$ case. Let $\{A_i\}_{i=1}^k$ and \mathcal{V} be as in the statement. For $l = 1, \dots, k$, let V_l and V'_l be the initial and terminal vertices of the arc A_l (again with respect to the orientations of A_l). Let E_l be the edge in direction 2 from V_l , and let U_l be the blocked vertex in direction 1 from V_l which forces E_l to be order-disrespecting. Define E'_l and U'_l similarly (so that E'_l is in direction 2 from V'_l and U'_l is in direction 1). Note we must have $2k \leq m - r'$ so $s + r = m \geq 2k + r'$, and

The cells Φ and Ψ are critical. As above, if $\{[c_1], \dots, [c_q]\}$ and $\{[d_1], \dots, [d_q]\}$ are the collections of equivalence classes of 1-cells having $[\Phi]$ and $[\Psi]$ as their least upper bounds, and if $[c_i] = [d_j]$ for some i and j , then $[c_i]$ and $[d_j]$ must have a common edge e which can only happen if $e = \tilde{E}_t$ for some t , so in particular, \tilde{V}_t is of degree 3. For such a t and each $\delta = 0, 1, 2$, let C_δ be the cloud in direction δ from \tilde{V}_t , and let

$$M_\delta = |\{p : V_p \in C_\delta\}|, \quad M'_\delta = |\{p : V'_p \in C_\delta\}|,$$

$$\widehat{M}_\delta = |\{p : \widehat{V}_p \in C_\delta\}|, \quad \widetilde{M}_\delta = |\{p : p \neq t, \tilde{V}_p \in C_\delta\}|.$$

Then, if $f(C_\delta)$ and $g(C_\delta)$ are the values of C_δ in the diagrams for $[c_i]$ and $[d_j]$, and λ_δ is as above, we have

$$f(C_\delta) = 2(M_\delta + \widehat{M}_\delta + \widetilde{M}_\delta) + \lambda_\delta,$$

$$g(C_\delta) = 2(M'_\delta + \widehat{M}_\delta + \widetilde{M}_\delta) + \lambda_\delta,$$

so that $f(C_\delta) - g(C_\delta) = 2(M_\delta - M'_\delta)$. Arguments similar given to those in the first part of the proof shows that $M_\delta - M'_\delta = \eta_\delta(\tilde{V}_t)$, so again we see that at least two of the clouds around \tilde{V}_t can be categorized as either positive or negative, and each of these clouds has a different value in the diagram for $[c_i]$ than it does in the diagram for $[d_j]$, so $[c_i] \neq [d_j]$.

For $\epsilon = 1$, construct Φ and Ψ as above with the following modifications. First, the vertices $\tilde{V}_1, \dots, \tilde{V}_{q-r'-k}$ are the first vertices in $\mathcal{W}' \cup \mathcal{W}$ (where $k = 0$ and $r' = r$ if $s' \geq q - r$). Next, since $n = 2q + 1 \geq 3$, we can again assume the initial and terminal endpoints of A_0 fall in a cloud in the collection of clouds determined by the edges in Φ and Ψ , respectively. Then let x be a vertex in the same cloud as the initial endpoint of A_0 (in the system of clouds determined by the edges of Φ) such that x is blocked in

$$\Phi = \{E_1, \dots, E_k, U_1, \dots, U_k, \widehat{E}_1, \dots, \widehat{E}_{r'}, \widehat{U}_1, \dots, \widehat{U}_{r'},$$

$$\tilde{E}_1, \dots, \tilde{E}_{q-r'-k}, \tilde{U}_1, \dots, \tilde{U}_{q-r'-k}, x\},$$

and let x' be a vertex in the same cloud as the terminal endpoint of A_0 (in the system of clouds determined by the edges of Ψ) such that x' is blocked in

$$\Psi = \{E'_1, \dots, E'_k, U'_1, \dots, U'_k, \widehat{E}'_1, \dots, \widehat{E}'_{r'}, \widehat{U}'_1, \dots, \widehat{U}'_{r'},$$

$$\tilde{E}_1, \dots, \tilde{E}_{q-r'-k}, \tilde{U}_1, \dots, \tilde{U}_{q-r'-k}, x'\}.$$

Here, we of course assume that x is distinct from the other vertices in Φ and x' is distinct from the other vertices in Ψ . If $s' \geq q - r$, the edges E_i and E'_i and the vertices U_i and U'_i do not appear in Φ and Ψ .

For example, for $n = 17$, we can let the arcs A_1 , A_2 , and A_3 be as in Figure 12, so $\mathcal{W} = \{\tilde{V}_1, \tilde{V}_2, \tilde{V}_3\}$, and trivially let the arc A_0 be the unique edge which has $*$ as an endpoint, so that $\mathcal{W}' = \emptyset$. For a less trivial example, if $q = 9$ and $\epsilon = 1$, we can let A_0 be the arc A_3 in Figure 12 slightly enlarged so that its interior includes its two original endpoints labeled V_3 and V'_3 in Figure 12, so that here now we have $k = 2$ and $r' = 2$. See Figure 13. The set \mathcal{W}' consists of the vertices labeled $\tilde{V}_1, \dots, \tilde{V}_5$, and $\mathcal{W} = \emptyset$ so that $|\mathcal{W}| = 0 = 9 - 2 - 2 - 5 = q - r' - k - s'$. Alternatively, we can let \mathcal{W}' consist of the vertices labeled \tilde{V}_3 and \tilde{V}_5 , and let \mathcal{W} consist of the vertices labeled \tilde{V}_1, \tilde{V}_2 and \tilde{V}_4 in Figure 13, so that $s' = 2$ and $|\mathcal{W}| = 3 = 9 - 2 - 2 - 2 = q - r' - k - s'$.

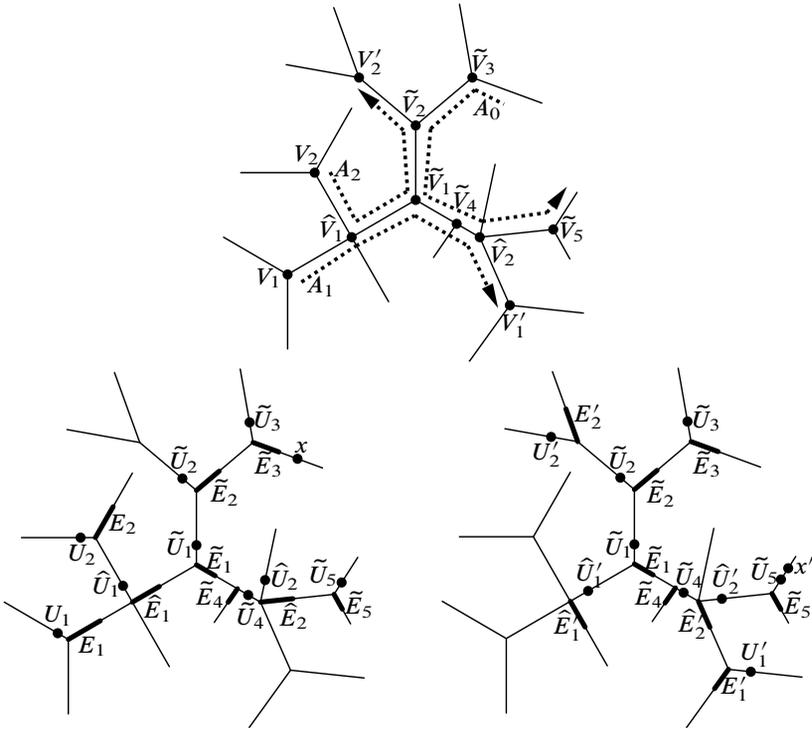


Figure 13: The critical cells Φ and Ψ in $UD^{19}(\Gamma)$ (bottom left and right)

Again, the cells Φ and Ψ are critical. If $\{[c_1], \dots, [c_q]\}$ and $\{[d_1], \dots, [d_q]\}$ are the collections of equivalence classes of 1-cells having $[\Phi]$ and $[\Psi]$ as their least upper bounds, and $[c_i] = [d_j]$ for some i and j , then again it must be the case that $[c_i]$ and $[d_j]$ have a common edge \tilde{E}_t , so \tilde{V}_t is of degree 3. Now, with the notation from the $\epsilon = 0$ case, we have

$$f(C_\delta) = 2(M_\delta + \hat{M}_\delta + \tilde{M}_\delta) + \lambda_\delta + \mu_\delta, \quad g(C_\delta) = 2(M'_\delta + \hat{M}_\delta + \tilde{M}_\delta) + \lambda_\delta + \mu'_\delta,$$

where $\mu_\delta = 1$ if x falls in C_δ , and $\mu_\delta = 0$ if x does not fall in C_δ , and similarly, $\mu'_\delta = 1$ if x' falls in C_δ , and $\mu'_\delta = 0$ if x' does not fall in C_δ . Since the endpoints of A_0 fall in clouds determined by the edges in Φ and Ψ , the vertex \tilde{V}_t cannot be an endpoint, so it either does not fall on A_0 or it falls on the interior of A_0 .

If \tilde{V}_t does not fall on A_0 , so that we have $\tilde{V}_t \in \mathcal{W}$, then x and x' are in the same cloud in the diagrams for $[c_i]$ and $[d_j]$, so that $\mu_\delta = \mu'_\delta$ for each δ , and therefore $f(C_\delta) - g(C_\delta) = 2(M_\delta - M'_\delta) = 2\eta_\delta(\tilde{V}_t)$, and since the collection $\{A_i\}_{i=1}^k$ is allowable for \mathcal{W} , we can categorize at least one cloud as positive and one as negative as above, so that $[c_i] \neq [d_j]$.

If \tilde{V}_t falls on the interior of A_0 , then it is possible that $\tilde{V}_t \in \mathcal{W}$ or $\tilde{V}_t \in \mathcal{W}'$, but in either case, x and x' must fall in different clouds C_{δ_1} and C_{δ_2} . Then we have

$$f(C_{\delta_1}) = 2(M_{\delta_1} + \widehat{M}_{\delta_1} + \widetilde{M}_{\delta_1}) + \lambda_{\delta_1} + 1, \quad g(C_{\delta_1}) = 2(M'_{\delta_1} + \widehat{M}_{\delta_1} + \widetilde{M}_{\delta_1}) + \lambda_{\delta_1} + 0,$$

so now $f(C_{\delta_1}) - g(C_{\delta_1}) = 2(M_{\delta_1} - M'_{\delta_1}) + 1$, which is odd and therefore nonzero. Likewise, $f(C_{\delta_2}) - g(C_{\delta_2}) = 2(M_{\delta_2} - M'_{\delta_2}) - 1$, which is again odd and therefore nonzero. Furthermore, since the sum of the values of the clouds around \tilde{V}_t must equal $n - 1$ in each cell, we have

$$f(C_0) - g(C_0) + f(C_1) - g(C_1) + f(C_2) - g(C_2) = 0,$$

so that $g(C_\delta) - f(C_\delta) > 0$ and $g(C_{\delta'}) - f(C_{\delta'}) < 0$ for two directions $\delta \neq \delta'$, so again at least one cloud is categorized as positive and one as negative, so $[c_i] \neq [d_j]$.

Therefore, in all cases, the conditions of [Lemma 3.6](#) are met, so

$$\text{TC}(\text{UD}^n(\Gamma)), \text{TC}(D^n(\Gamma)) \geq 2q + 1.$$

This, combined with the upper bounds stated at the beginning of the proof, gives the result. □

It is worth noting that if we insist no vertices have degree 3, the statement of [Theorem 1.2](#) becomes much simpler and determines the topological complexity for all n for both configuration spaces, provided they are connected.

Corollary 3.7 *Let Γ be a tree with no vertices of degree 3. Let $k = \min\{\lfloor \frac{n}{2} \rfloor, m(\Gamma)\}$. Then $\text{TC}(\text{UC}^n(\Gamma)) = 2k + 1$. Also, if $n = 1$ or $m(\Gamma) \geq 1$, then $\text{TC}(C^n(\Gamma)) = 2k + 1$.*

Proof Let r and s be as in [Theorem 1.2](#), so that $r = m(\Gamma) =: m$, and $s = 0$. If $m \geq 1$ and $n \geq 2m$, the claim follows from statement (1) of the theorem. If $n = 2q + \epsilon < 2m = 2r$, with $\epsilon \in \{0, 1\}$ and $q \geq 1$, then $r > q$, so $2(q - r) < 0 = s$,

and the claim follows from statement (a). If $m = 0$, so that Γ is homeomorphic to a closed interval, then Lemma 3.4 gives $\text{TC}(\text{UD}^n(\Gamma)) \leq 1$, but $\text{TC}(X) \geq 1$ for any space, so, $\text{TC}(\text{UC}^n(\Gamma)) = 1$. For the remaining case, $n = 1$ and $m \geq 1$, we have $C^n(\Gamma) = \text{UC}^n(\Gamma) = \Gamma$, so all three spaces have topological complexity 1 since Γ is contractible. \square

Now we discuss how in some sense, the results in Theorem 1.2 are the best we can achieve with the methods used here. Consider the case $2m \leq n < 2m + k$, with m and k as in the first part of the theorem. In this case, a construction similar to the one given in the proof shows that there is a critical m -cell, which corresponds to a nonzero m -dimensional cohomology class, so that the space $\text{UD}^n(\Gamma)$ cannot be homotopic to a space of dimension less than m , so the dimensional bound given in Theorem 3.2 cannot improve the bound given by the explicit motion planning algorithm in Lemma 3.4. Likewise, if q, r , and s are as in the second part of the theorem, but the appropriate collection of arcs does not exist, then there will still be a critical q -cell, so again the dimension cannot improve the upper bound. The following shows that Lemma 3.6 cannot be used to get improved lower bounds.

Proposition 3.8 *Let Γ be a tree with $m := m(\Gamma) \geq 1$.*

(1) *Let k be as in Theorem 1.2(1), and assume there is at least one vertex of degree 3, so that $k \geq 1$. Let n satisfy $n \geq 2m$, and consider any two critical m -cells Φ and Ψ of $\text{UD}^n(\Gamma)$. If $\{[c_1], \dots, [c_m]\}$ and $\{[d_1], \dots, [d_m]\}$ are the unique collections of equivalence classes of 1-cells having $[\Phi]$ and $[\Psi]$ as their least upper bounds, and for all i and j we have $[c_i] \neq [d_j]$, then $n \geq 2m + k$.*

(2) *Let q, r, s and ϵ be as in Theorem 1.2(b), and consider critical q -cells Φ and Ψ of $\text{UD}^n(\Gamma)$. If $\{[c_1], \dots, [c_q]\}$ and $\{[d_1], \dots, [d_q]\}$ are the unique collections of equivalence classes of 1-cells having $[\Phi]$ and $[\Psi]$ as their least upper bounds and for all i and j we have $[c_i] \neq [d_j]$, then:*

- (a) *If $\epsilon = 0$, there is some $k \geq 1$ such that there exist oriented arcs A_1, \dots, A_k with the following properties:*
 - (i) *The endpoints of each A_l are (distinct) essential vertices, neither of which is an endpoint of any other $A_{l'}$.*
 - (ii) *There are $r' \leq r$ vertices of degree greater than 3 which are not the endpoints of any A_l .*
 - (iii) *There is a collection \mathcal{V} of degree-3 vertices, with $|\mathcal{V}| \geq q - r' - k$ such that $\{A_i\}_{i=1}^k$ is allowable for \mathcal{V} .*

- (b) If $\epsilon = 1$, there is an arc A_0 whose endpoints have no restrictions and whose interior includes a collection \mathcal{W}' of $s' \leq q$ distinct vertices of degree 3, and if $s' < q - r$, there are arcs A_1, \dots, A_k , as above, whose endpoints are also not vertices in \mathcal{W}' , and there is another collection of degree-3 vertices, \mathcal{W} , such that $\mathcal{W} \cap \mathcal{W}' = \emptyset$, $|\mathcal{W}| \geq q - r' - k - s'$ and $\{A_i\}_{i=1}^k$ is allowable for \mathcal{W} , where r' is as above.

Proof We use the contrapositive for both statements (1) and (2). For the first statement, let Φ and Ψ be critical m -cells, and assume $2m \leq n < 2m + k$. For each essential vertex v , each cell must contain exactly one edge e having $\iota(e) = v$ and a blocked vertex u which makes e order-disrespecting. There is no choice for e or u if v is of degree 3.

If $n > 2m$, let w_1, \dots, w_{n-2m} be the remaining vertices in Φ , and let w'_1, \dots, w'_{n-2m} be the remaining vertices in Ψ . For $l = 1, \dots, n - 2m < k$, let A_l be the geodesic from w_l to w'_l , oriented so that w_l is the initial endpoint and w'_l is the terminal endpoint if $w_l \neq w'_l$. If $w_l = w'_l$, extend A_l slightly so that it is a small arc starting at w_l which doesn't intersect any essential vertices. This gives a collection of less than k oriented arcs in Γ , so by the minimality of k , this collection cannot be allowable for the set of vertices of degree 3 in Γ . It is not possible that a degree-3 vertex v is an endpoint of any A_l , since no w_l or w'_l can be essential. Therefore, there is some degree-3 vertex v which has the property $\eta_0(v) = \eta_1(v) = \eta_2(v) = 0$. Let e be the edge in direction 2 from v . This edge must be in both Φ and Ψ . Suppose i and j have the property that e is the unique edge in $[c_i]$ and $[d_j]$. Then, in the notation from the proof of Theorem 1.2, for $\delta = 0, 1, 2$ and each cloud C_δ in the diagrams for $[c_i]$ and $[d_j]$, we have $f(C_\delta) = 2m_\delta + W_\delta + \lambda_\delta$ and $g(C_\delta) = 2m_\delta + W'_\delta + \lambda_\delta$. But, as above, we have $W_\delta - W'_\delta = \eta_\delta(v) = 0$ for each δ , and therefore $W_\delta = W'_\delta$ and $f(C_\delta) = g(C_\delta)$, for each δ , so $[c_i] = [d_j]$.

If $n = 2m$, then for each vertex v of degree 3, the edge e in direction 2 from v must appear in both cells, so again if $[c_i]$ and $[d_j]$ are the equivalence classes which contain the edge e , we now have $f(C_\delta) = 2m_\delta + \lambda_\delta = g(C_\delta)$ for each $\delta = 0, 1, 2$, so $[c_i] = [d_j]$.

For statement (2), assume that the collections of arcs and vertices in (a) or (b) do not exist for the appropriate value of ϵ . Note first that if $\epsilon = 0$, any critical q -cell must consist solely of edges e_1, \dots, e_q and blocked vertices u_1, \dots, u_q which force the edges e_1, \dots, e_q to be order-disrespecting. In particular, each e_i has the property that $\iota(e_i)$ is

an essential vertex. If $\epsilon = 1$, then the same is true, except that the cell contains one additional blocked vertex. Let Φ and Ψ be critical q -cells, and let \mathcal{S} (resp. \mathcal{S}') be the set of essential vertices v such that $\iota(e) = v$ for some e in Φ (resp. Ψ), so $|\mathcal{S}| = |\mathcal{S}'| = q$.

Let $\mathcal{T} = \mathcal{S} \setminus (\mathcal{S} \cap \mathcal{S}')$ and $\mathcal{T}' = \mathcal{S}' \setminus (\mathcal{S} \cap \mathcal{S}')$, so $|\mathcal{T}| = |\mathcal{T}'| =: k$. The fact that $s < 2(q-r)$ implies that at least one vertex of degree 3 appears in $\mathcal{S} \cap \mathcal{S}'$. Let $\tilde{V}_1, \dots, \tilde{V}_{\bar{s}}$ be the vertices of degree 3 in $\mathcal{S} \cap \mathcal{S}'$, so that $\bar{s} > 0$ and let $\hat{V}_1, \dots, \hat{V}_{\bar{r}}$ be the vertices of degree greater than 3 in $\mathcal{S} \cap \mathcal{S}'$. The edge in direction 2 from each vertex \tilde{V}_t must be included in both Φ and Ψ . In what follows, $[c_i]$ and $[d_j]$ are the equivalence classes which contain the edge in direction 2 from whichever vertex \tilde{V}_t is being discussed.

If $\epsilon = 1$, let x (resp. x') be the additional vertex in Φ (resp. Ψ). Let A_0 be the geodesic from x to x' and extend A_0 slightly if $x = x'$ so that it is a small arc which doesn't intersect any vertex in \mathcal{S} or \mathcal{S}' . Let \mathcal{W}' be the set of degree-3 vertices in $\mathcal{S} \cap \mathcal{S}'$ which fall on the interior of A_0 , and let $s' = |\mathcal{W}'|$. By the assumption of the nonexistence of the appropriate collections of arcs and vertices, we have $s' < q - r$.

First consider the case $k = 0$. Here, we have $\mathcal{S} = \mathcal{S}'$, and $\bar{s} + \bar{r} = q$. If $\epsilon = 0$, then for each vertex $\tilde{V}_t \in \mathcal{S} \cap \mathcal{S}' = \mathcal{S} = \mathcal{S}'$, we have, with the notation from the proof of [Theorem 1.2](#), $f(C_\delta) = 2(\tilde{M}_\delta + \hat{M}_\delta) + \lambda_\delta = g(C_\delta)$, so $[c_i] = [d_j]$. If $\epsilon = 1$, since we have $s' < q - r$, there must be at least one vertex \tilde{V}_t which does not fall on A_0 . Then x and x' must fall in the same cloud in the system of clouds determined by the edge in direction 2 from \tilde{V}_t . Again using the notation from the proof of [Theorem 1.2](#), we have $\mu_\delta = \mu'_\delta$, so $f(C_\delta) = 2(\tilde{M}_\delta + \hat{M}_\delta) + \lambda_\delta + \mu_\delta = 2(\tilde{M}_\delta + \hat{M}_\delta) + \lambda_\delta + \mu'_\delta = g(C_\delta)$, and therefore $[c_i] = [d_j]$.

If $k \geq 1$, let V_1, \dots, V_k and V'_1, \dots, V'_k be the vertices in \mathcal{T} and \mathcal{T}' , respectively, and let A_l be the oriented geodesic from V_l to V'_l . This gives a collection of k arcs with $2k$ distinct essential endpoints if $\epsilon = 0$, and one additional arc A_0 if $\epsilon = 1$. Also note that $\bar{s} + \bar{r} + k = q$ and $\bar{r} \leq r'$, where r' is the number of vertices of degree greater than 3 which are not endpoints of any A_l , so $\bar{s} = q - \bar{r} - k \geq q - r' - k$.

For $\epsilon = 0$, let $\mathcal{V} = \{\tilde{V}_1, \dots, \tilde{V}_{\bar{s}}\}$, so that $|\mathcal{V}| = \bar{s} \geq q - r' - k$. By assumption, the arcs $\{A_i\}_{i=1}^k$ cannot be allowable for \mathcal{V} . Any vertex $v \in \mathcal{V}$ cannot be the endpoint of any A_l since no endpoint is in $\mathcal{S} \cap \mathcal{S}'$. So, there is at least one degree-3 vertex \tilde{V}_t such that $\eta_0(\tilde{V}_t) = \eta_1(\tilde{V}_t) = \eta_2(\tilde{V}_t) = 0$. So, we have $f(C_\delta) - g(C_\delta) = 2(M_\delta - M'_\delta) = 2\eta_\delta(v) = 0$, for each $\delta \in \{0, 1, 2\}$, so each cloud has the same value in the diagram for $[c_i]$ as it does in the diagram for $[d_j]$. So, $[c_i] = [d_j]$.

For $\epsilon = 1$, let \mathcal{W} be the set of degree-3 vertices in $(\mathcal{S} \cap \mathcal{S}') \setminus \mathcal{W}'$, so that $\mathcal{W} \cap \mathcal{W}' = \emptyset$ and $|\mathcal{W}| = \bar{s} - s' \geq q - r' - k - s'$, and therefore, again by assumption, $\{A_i\}_{i=1}^k$ is not allowable for \mathcal{W} . No vertex in \mathcal{W} can be an endpoint of any arc A_l , so there must be some vertex $\tilde{V}_t \in \mathcal{W}$ with $\eta_\delta(\tilde{V}_t) = 0$ for $\delta = 0, 1, 2$. Now we have $f(C_\delta) - g(C_\delta) = 2(M_\delta - M'_\delta) + \mu_\delta - \mu'_\delta$, but since \tilde{V}_t is not on A_0 , we must have both x and x' in the same cloud in the system of clouds determined by the edge in direction 2 from \tilde{V}_t , so $\mu_\delta = \mu'_\delta$, and therefore $f(C_\delta) - g(C_\delta) = 2(M_\delta - M'_\delta) = 2\eta_\delta(\tilde{V}_t) = 0$ for each direction δ , so again we have $[c_i] = [d_j]$. \square

To give a better idea of the values of n for which [Theorem 1.2](#) determines $\text{TC}(\text{UC}^n(\Gamma))$ (and $\text{TC}(C^n(\Gamma))$), we consider all values of $n \geq 2$ with Γ as in the top of [Figure 9](#), where we have $m = 11$, $r = 2$, and $s = 9$. If n is either sufficiently large or sufficiently small, it is easy to determine if the theorem applies. For $n \geq 25$, statement (1) of the theorem applies, and for $2 \leq n \leq 13$, statement (a) applies. For $22 \leq n \leq 24$, the theorem does not apply. For the “middle” values of n , namely $14 \leq n \leq 21$, only statement (b) might apply.

For $n = 14$ statement (b) does apply, since we may chose the arcs A_1, A_2 , and A_3 as in [Figure 12](#), and let \mathcal{V} consist of the vertices \tilde{V}_1, \tilde{V}_2 , and \tilde{V}_3 (as labeled in [Figure 12](#)). Then we have $k = 3, r' = 2, |\mathcal{V}| = 3 \geq q - r' - k = 7 - 2 - 3 = 2$, and the arcs A_1, A_2 , and A_3 are again allowable for \mathcal{V} . Note the vertex \tilde{V}_3 would not be used in the construction of the cells Φ and Ψ for $n = 14$. For $n = 15$, we can again let A_0 be the unique edge which has $*$ as one of its endpoints, and keep A_1, A_2 , and A_3 the same, so the theorem also applies for $n = 15$. The cases $n = 16, n = 17$, and $n = 19$ are covered in the proof.

For $n = 18$, suppose the appropriate collection of arcs A_1, \dots, A_k exists and is allowable for a set \mathcal{V} of vertices of degree 3, so that we have $|\mathcal{V}| \geq q - r' - k$. On the other hand, there must be a total of $2k$ distinct vertices which are the endpoints of the arcs. No vertex in \mathcal{V} can be any of these endpoints, so we must have $|\mathcal{V}| \leq s - 2k + r - r'$, since there are $r - r'$ vertices of degree greater than 3 which are the endpoint of some arc, and the remaining $2k - (r - r')$ endpoints must be of degree 3. So, we have

$$9 - r' - k = q - r' - k \leq |\mathcal{V}| \leq s - 2k + r - r' = 9 - 2k + 2 - r'.$$

Comparing the left and right sides, we see that we must have $k \leq 2$. Since we also must have $r' \leq r = 2$, this gives $|\mathcal{V}| \geq 9 - 2 - 2 = 5$. But, the only vertices of degree 3 which can fall on the interior of any arc with essential endpoints are the vertices labeled

v_4 , v_5 , and v_8 in [Figure 9](#). So $|\mathcal{V}| \leq 3$, arriving at a contradiction, so the theorem does not apply for $n = 18$, and for similar reasons, the theorem does not apply for $n = 20$.

For $n = 21$, suppose the appropriate collection of arcs A_0, A_1, \dots, A_k and collections of vertices \mathcal{W} and \mathcal{W}' exist, so that $|\mathcal{W} \cup \mathcal{W}'| \geq q - r' - k = 10 - r' - k$. Similar to above, since no vertex in $\mathcal{W} \cup \mathcal{W}'$ can be an endpoint of any arc A_1, \dots, A_k , we must have

$$10 - r' - k = q - r' - k \leq |\mathcal{W} \cup \mathcal{W}'| \leq s - 2k + r - r' = 9 - 2k + 2 - r',$$

so that $k \leq 1$. Since again we must have $r' \leq 2$, this gives $|\mathcal{W} \cup \mathcal{W}'| \geq 10 - r' - k \geq 10 - 2 - 1 = 7$. However, again the only vertices of degree 3 which can fall on the interior of any arc A_1 are the vertices labeled v_4 , v_5 , and v_8 in [Figure 9](#), and here it is clear that at most two additional vertices could be included in the interior of the arc A_0 , so that $|\mathcal{W} \cup \mathcal{W}'| \leq 5$, a contradiction, so the theorem does not apply for $n = 21$.

So, the only values of $n \geq 2$ for which [Theorem 1.2](#) does not determine $\text{TC}(\text{UC}^n(\Gamma))$ and $\text{TC}(C^n(\Gamma))$ are $n = 18$ and $20 \leq n \leq 24$.

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