

Peak reduction and finite presentations for automorphism groups of right-angled Artin groups

MATTHEW B DAY

We generalize the peak reduction algorithm (Whitehead’s theorem) for free groups to a theorem about a general right-angled Artin group A_Γ . As an application, we find a finite presentation for the automorphism group $\text{Aut}A_\Gamma$ that generalizes McCool’s presentation for the automorphism group of a finite rank free group. We also consider a stronger generalization of peak reduction, giving a counterexample and proving a special case.

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1 Introduction

1.1 Background and statement of results

Let Γ be a graph on n vertices, with vertex set X and adjacency relation denoted by $e(-, -)$. Let A_Γ be the *right-angled Artin group* of Γ , defined by

$$A_\Gamma := \langle X \mid R_\Gamma \rangle$$

where the relations are $R_\Gamma = \{[x, y] \mid x, y \in X \text{ and } e(x, y)\}$ (we use the convention that $[x, y] = xyx^{-1}y^{-1}$). If Γ is the edgeless graph (n vertices and no edges), then A_Γ is the free group F_n on n generators. If Γ is the complete graph, then A_Γ is the free abelian group \mathbb{Z}^n . So in a sense, the group A_Γ interpolates between free groups and free abelian groups as we vary Γ . Similarly, automorphism group $\text{Aut}A_\Gamma$ interpolates between $\text{Aut}F_n$ and the integral general linear group $\text{GL}(n, \mathbb{Z})$.

In this paper, we develop a framework for understanding $\text{Aut}A_\Gamma$ in which ideas from the study of linear groups and ideas from the study of $\text{Aut}F_n$ can both be applied. Our first main result is the following.

Theorem A *For any graph Γ , the group $\text{Aut}A_\Gamma$ is finitely presented. Moreover, for each Γ , we have an explicit finite presentation for $\text{Aut}A_\Gamma$.*

See Theorem 2.7 below for a precise statement. Although Bux–Charney–Vogtmann [1, Corollary 13] have shown that $\text{Aut } A_\Gamma$ is finitely presented when Γ is a tree, our result is more explicit and holds for arbitrary Γ .

To prove this, we transfer an idea from the study of automorphism groups of free groups: *peak reduction*. Whitehead’s 1936 theorem [14, Theorem 2] is a result about automorphism groups of free groups with important applications; peak reduction is an algorithmic approach used by Rapaport [11], Higgins–Lyndon [5] and others to reprove and extend this theorem. Whitehead’s theorem states that there is a finite generating set for $\text{Aut } F_n$ that has a special property concerning factorizations of elements of $\text{Aut } F_n$ and the lengths of elements of F_n . One corollary of this theorem is that for any k -tuple W of elements of F_n , the stabilizer $(\text{Aut } F_n)_W$ is finitely generated; in fact, McCool [9] used peak reduction methods to prove that $(\text{Aut } F_n)_W$ is finitely presented. Another corollary is that there is an algorithm that determines whether two k -tuples of elements of F_n are in the same orbit under the action of $\text{Aut } F_n$ (see Lyndon–Schupp [7, Proposition I.4.19]). Peak-reduction methods were also used by Culler–Vogtmann [3] to study the structure of the outer space of F_n . In the present paper, we give a peak reduction theorem for right-angled Artin groups (Theorem B below). Theorem A is the only application of peak reduction we pursue in this paper; however, we discuss other potential applications in Section 6.

We proceed to define the notion of a peak-reduced factorization. Define the length of a conjugacy class of A_Γ to be the minimum of the lengths of its representative elements (with respect to X), and define the length of a k -tuple of conjugacy classes of A_Γ to be the sum of the lengths of its elements (for any $k \geq 1$). For W a k -tuple of conjugacy classes in A_Γ , we say that a string $\alpha_m \cdots \alpha_1$ of elements of $\text{Aut } A_\Gamma$ is *peak-reduced* with respect to W if for each $i = 1, \dots, m-1$, we do not have both

$$\begin{aligned} |(\alpha_{i+1} \cdots \alpha_1) \cdot W| &\leq |(\alpha_i \cdots \alpha_1) \cdot W| \\ \text{and} \quad |(\alpha_i \cdots \alpha_1) \cdot W| &\geq |(\alpha_{i-1} \cdots \alpha_1) \cdot W| \end{aligned}$$

unless all three lengths are equal.

Suppose $G < \text{Aut } A_\Gamma$, S is a finite generating set for G , and W is a k -tuple of conjugacy classes in A_Γ . We say that G has *peak reduction* with respect to W by elements of S if every $\alpha \in G$ has a factorization by elements of S that is peak-reduced with respect to W . The peak reduction theorem for a free group F_n states that there is a finite generating set Ω for $\text{Aut } F_n$ (called the *Whitehead automorphisms*) such that $\text{Aut } F_n$ has peak reduction with respect to any k -tuple of conjugacy classes W in F_n by elements of Ω . See Lyndon–Schupp [7, Proposition I.4.20] for a proof. Whitehead’s theorem is an important special case of the peak reduction theorem.

The following definition of a Whitehead automorphism in $\text{Aut } A_\Gamma$ generalizes the definition of a Whitehead automorphism of a free group (see page 31 of Lyndon–Schupp [7]).

Definition 1.1 A *Whitehead automorphism* is an element $\alpha \in \text{Aut } A_\Gamma$ of one of the following two types:

Type (1): α restricted to $X \cup X^{-1}$ is a permutation of $X \cup X^{-1}$.

Type (2): There is an element $a \in X \cup X^{-1}$, called the *multiplier* of α , such that $\alpha(a) = a$ and for each $x \in X$, the element $\alpha(x)$ is in $\{x, xa, a^{-1}x, a^{-1}xa\}$.

Let Ω be the set of all Whitehead automorphisms of A_Γ .

We use the graph Γ to define two subsets of Ω that generate relatively well-behaved subgroups of $\text{Aut } A_\Gamma$.

Definition 1.2 Let α be a type (2) Whitehead automorphism α with multiplier a . Call α *short-range* if α fixes the elements of X not adjacent to a (so that α acts nontrivially only on elements near a , ie adjacent to a). Call α *long-range* if α fixes the elements of X adjacent to a (so that α acts nontrivially only on elements far away from a , ie not adjacent to a). We also consider type (1) Whitehead automorphisms to be long-range. Let Ω_s and Ω_ℓ denote the sets of short-range and long-range Whitehead automorphisms respectively.

It is easy to see that Ω is finite, and it is a consequence of the work of Laurence [6] (see Section 2.2) that $\Omega_\ell \cup \Omega_s$ generates $\text{Aut } A_\Gamma$. Our second main result further clarifies the roles of Ω_s and Ω_ℓ as follows.

Theorem B (Peak-reduction theorem) *The finite generating set $\Omega_\ell \cup \Omega_s$ for $\text{Aut } A_\Gamma$ has the following properties:*

- (1) Each $\alpha \in \text{Aut } A_\Gamma$ can be written as $\alpha = \beta\gamma$ for some $\beta \in \langle \Omega_s \rangle$ and some $\gamma \in \langle \Omega_\ell \rangle$.
- (2) The usual representation $\text{Aut } A_\Gamma \rightarrow \text{Aut } H_1(A_\Gamma)$ to the automorphism group of the abelianization $H_1(A_\Gamma)$ of A_Γ restricts to an embedding $\langle \Omega_s \rangle \hookrightarrow \text{Aut } H_1(A_\Gamma)$.
- (3) The subgroup $\langle \Omega_\ell \rangle$ has peak reduction by elements of Ω_ℓ with respect to any k -tuple W of conjugacy classes in A_Γ .

The proof of Theorem B is effective: there is an algorithm that splits an automorphism into its $\langle \Omega_\ell \rangle$ -part and its $\langle \Omega_s \rangle$ -part and an algorithm that peak-reduces an element of $\langle \Omega_\ell \rangle$. Further, the theorem implies that we can analyze an element of $\langle \Omega_s \rangle$ by using row reduction methods in $\text{Aut } H_1(A_\Gamma) \cong \text{GL}(n, \mathbb{Z})$. Note that if A_Γ is a free group, Ω_ℓ is Whitehead's generating set Ω , our Ω_s contains only the identity, and Theorem B restricts to the classical peak-reduction theorem.

At this point, we can say a little more about Theorem A. In Section 2.3, we define a finite set R of relations among the Whitehead automorphisms Ω . These relations tell us when one element of Ω is the inverse of another, when one is a product of two others, when two elements of Ω commute, when an element of Ω is the commutator of two other elements of Ω and how type (1) Whitehead automorphisms interact with type (2) Whitehead automorphisms. Our presentation is then $\text{Aut } A_\Gamma = \langle \Omega \mid R \rangle$. These relations are based on the relations McCool gives for the automorphism group of the free group in [8] and our presentation restricts to McCool's presentation for $\text{Aut } F_n$ in the case of the free group. The idea behind Theorem A is that we can use the methods of Theorem B (using all three parts of the theorem) to put any word in Ω representing the identity in $\text{Aut } A_\Gamma$ into a standard form. We also use the fact that the inner automorphism group of a right-angled Artin group is also a right-angled Artin group.

We also investigate the possibility of a stronger generalization of the peak reduction theorem to right-angled Artin groups. For most right-angled Artin groups, $\langle \Omega_\ell \rangle$ is a proper subgroup of $\text{Aut } A_\Gamma$, and part (3) of Theorem B is seemingly weaker than the peak reduction theorem for free groups. In fact, we cannot hope for a straightforward generalization of peak reduction that applies to all of $\text{Aut } A_\Gamma$ for arbitrary Γ . We show the following in Section 4.1.

Proposition C *There is a graph Γ such that for every finite generating set S of $\text{Aut } A_\Gamma$, there is a conjugacy class w in A_Γ and an automorphism $\alpha \in \text{Aut } A_\Gamma$ such that α cannot be peak-reduced with respect to w by members of S .*

In spite of this disappointing fact, there are still special cases where peak-reduction works. As an example of such a special case, we prove the following:

Proposition D *Let $W = (w_1, \dots, w_k)$ be a k -tuple of conjugacy classes such that for each i , $|w_i| = 1$. If $\alpha \in \text{Aut } A_\Gamma$ and $|\alpha \cdot W| = |W|$, then α can be peak-reduced with respect to W by elements of $\Omega_\ell \cup \Omega_s$.*

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2 Generating sets for right-angled Artin groups

2.1 Combinatorial group theory of A_Γ

Let the set of letters L be $X \cup X^{-1}$. For $x \in L$, let $v(x) \in X$, the *vertex of x* , be the unique element of $X \cap \{x, x^{-1}\}$. We will use $e(x, y)$ as a shorthand for $e(v(x), v(y))$ and we will use $\text{st}_L(x)$ and $\text{lk}_L(x)$ as notation for $\text{st}(v(x)) \cup \text{st}(v(x))^{-1}$ and $\text{lk}(v(x)) \cup \text{lk}(v(x))^{-1}$ respectively.

As usual, a word in L represents an element in A_Γ . A *cyclic word* is a string of elements of L indexed cyclically (or alternatively, an equivalence class of words under cyclic permutation of the indices). Any two noncyclic indexings of a cyclic word w represent group elements that are conjugate to each other, so w represents a well-defined conjugacy class. If w is a cyclic word, we will use $[w]$ to denote the conjugacy class it represents. If w is a noncyclic word, we will sometimes use $[w]$ to denote the cyclic word or conjugacy class it represents.

We will denote the length of a word or cyclic word w by $|w|$. The length of a group element or conjugacy class is the minimum length of any of its representative words or cyclic words, respectively. A word or cyclic word w on L is *graphically reduced* if it contains no subsegments of the form aua^{-1} , where $a \in L$ and u is a word in $\text{lk}_L(a)$. The *support* $\text{supp } w$ of a word or cyclic word w is the subset of X consisting of all generators that appear (or whose inverses appear) in w , and the support $\text{supp } W$ of a k -tuple $W = (w_1, \dots, w_k)$ of conjugacy classes is $\bigcup_{i=1}^k \text{supp } w_i$.

According to Servatius [13, Section I] any graphically reduced word can be transformed into any other graphically reduced representative of the same element by repeated

application of commutation moves (replacing a subsegment ab with ba when $e(a, b)$). The same is true for cyclic words and conjugacy classes. Therefore, we take the support $\text{supp } w$ of a group element or conjugacy class to be the support of any graphically reduced representative. The number of instances of a given generator in a group element or conjugacy class can be defined in the same way.

Servatius's centralizer theorem from [13, Section III] finds all the centralizers of elements in A_Γ . We restate a special case here:

Theorem 2.1 (Special case of Servatius's centralizer theorem) *For $x \in X$, the centralizer of x in A_Γ is $\langle \text{st}_L(x) \rangle$.*

2.2 Laurence's generators for $\text{Aut } A_\Gamma$

There is a reflexive, transitive, binary relation on X called the *domination relation*: say $x \geq y$ (x dominates y) if $\text{lk}(y) \subset \text{st}(x)$. Domination is clearly reflexive and transitive. For $x, y \in L$, say $x \geq y$ if $v(x) \geq v(y)$. Write $x \sim y$ when $x \geq y$ and $y \geq x$; the relation \sim is called the *domination equivalence* relation. We will also consider the *adjacent domination* relation, which holds for x and y if $e(x, y)$ and $x \leq y$, and the *nonadjacent domination* relation, which holds if $x \leq y$ and not $e(x, y)$. Each of these relations has a corresponding equivalence relation. We say that x strictly dominates y if $x \geq y$ and $x \not\sim y$ (other authors have used "strict domination" to refer to what we mean by "nonadjacent domination").

The following classes of automorphisms were defined by Servatius in [13], where he conjectured that they generate $\text{Aut } A_\Gamma$.

Definition 2.2 The Laurence–Servatius generators are the following four classes of automorphisms:

- (Dominated transvections) For $x, y \in L$ with $x \geq y$ and $v(x) \neq v(y)$, the *dominated transvection* (or simply *transvection*) $\tau_{x,y}$ is the automorphism that sends

$$y \mapsto yx$$

and fixes all generators not equal to $v(y)$.

- (Partial conjugations) For $x \in L$ and Y a union of connected components of $\Gamma - \text{st}(v(x))$, the *partial conjugation* $c_{x,Y}$ is the automorphism that sends

$$y \mapsto x^{-1}yx \quad \text{for } y \in Y$$

and fixes all generators not in Y .

- (Inversions) For $x \in X$, the *inversion* of x is the automorphism that sends

$$x \mapsto x^{-1}$$

and fixes all other generators.

- (Graphic automorphisms) For π an automorphism of the graph Γ , the *graphic automorphism* of π is the automorphism that sends

$$x \mapsto \pi(x)$$

for each generator $x \in X$.

It is a potential point of confusion that an ordinary conjugation automorphism is an example of a partial conjugation automorphism.

The following is Theorem 6.9 of Laurence [6].

Theorem 2.3 (Laurence) *The group $\text{Aut } A_\Gamma$ is generated by the finite set consisting of all dominated transvections, partial conjugations, inversions and graphic automorphisms of A_Γ .*

2.3 Whitehead automorphisms for right-angled Artin groups

We start with some comments on the Whitehead automorphisms Ω defined the introduction. There is a special notation for type (2) Whitehead automorphisms. Let $A \subset L$ and $a \in L$, such that $a \in A$ and $a^{-1} \notin A$. The symbol (A, a) denotes the Whitehead automorphism satisfying

$$(A, a)(a) = a$$

and for $x \in X - v(a)$,

$$(A, a)(x) = \begin{cases} x & \text{if } x \notin A \text{ and } x^{-1} \notin A \\ xa & \text{if } x \in A \text{ and } x^{-1} \notin A \\ a^{-1}x & \text{if } x \notin A \text{ and } x^{-1} \in A \\ a^{-1}xa & \text{if } x \in A \text{ and } x^{-1} \in A \end{cases}$$

if such an automorphism exists. Say that (A, a) is *well defined* if the formula given above defines an automorphism of A_Γ . For $\alpha \in \Omega$ of type (2), one can always find a multiplier $a \in L$ and a subset $A \subset L$ such that $\alpha = (A, a)$. There is a little ambiguity in choosing such a representation that comes from the following fact: if $a, b \in L$ with $e(a, b)$, then $(\{a, b, b^{-1}\}, a)$ is the trivial automorphism.

Note that the set of type (1) Whitehead automorphisms is the finite subgroup of $\text{Aut } A_\Gamma$ generated by the graphic automorphisms and inversions.

Claim 2.4 *The sets $\Omega_\ell \cup \Omega_s$ and Ω are both finite generating sets for $\text{Aut } A_\Gamma$.*

Proof It is easy to see that any dominated transvection is either in Ω_s or Ω_ℓ , and that any partial conjugation, inversion or graphic automorphism is in Ω_ℓ . So $\Omega_\ell \cup \Omega_s$ contains the Laurence–Servatius generators from Theorem 2.3. There are only finitely many permutations of L and finitely many subsets of L , so Ω is finite. Since $\Omega_\ell \cup \Omega_s \subset \Omega$, this proves the claim. \square

Lemma 2.5 *For $A \subset L$ with $a \in A$ and $a^{-1} \notin A$, the automorphism (A, a) is well defined if and only if both of the following hold:*

- (1) *The set $X \cap A \cap A^{-1} - \text{lk}(v(a))$ is a union of connected components of $\Gamma - \text{st}(v(a))$.*
- (2) *For each $x \in (A - A^{-1})$, we have $a \geq x$.*

Alternatively, (A, a) is well defined if and only if for each $x \in A - \text{st}_L(a)$ with $a \not\geq x$, (A, a) acts on the entire component of $v(x)$ in $\Gamma - \text{st}(v(a))$ by conjugation.

Proof The alternate statement follows immediately from the first one. For the “only if” direction of the first statement, note that if both conditions hold, one can write (A, a) as a product of the Laurence–Servatius generators. For the other direction, assume either condition fails and (A, a) defines an automorphism. One can then find elements $x, y \in X$ such that $[x, y] = 1$, but $[(A, a)(x), (A, a)(y)] \neq 1$ by Theorem 2.1. This is a contradiction. \square

2.4 Relations among Whitehead automorphisms

In this section we define the set of relations for our presentation for $\text{Aut } A_\Gamma$. Note that we use function composition order and automorphisms act on the left. With sets, we use the notation $A + B$ for $A \cup B$ when $A \cap B = \emptyset$. Note the shorthands $A - a$ for $A - \{a\}$ and $A + a$ for $A + \{a\}$.

Let Φ be the free group generated by the set Ω . We understand the relation “ $w_1 = w_2$ ” to correspond to $w_1 w_2^{-1} \in \Phi$. Note that if $(A, a) \in \Omega$ and $B \subset \text{lk}(v(a))$ such that $(B \cup B^{-1}) \cap A = \emptyset$, then (A, a) and $(A + B + B^{-1}, a)$ represent the same element of Ω and therefore the same element of Φ . This is why we do not list “ $(A, a) = (A + B + B^{-1}, a)$ ” in the relations below.

Definition 2.6 The relations of type (R1) are

$$(R1) \quad (A, a)^{-1} = (A - a + a^{-1}, a^{-1})$$

for $(A, a) \in \Omega$.

The relations of type (R2) are

$$(R2) \quad (A, a)(B, a) = (A \cup B, a)$$

for (A, a) and $(B, a) \in \Omega$ with $A \cap B = \{a\}$.

The relations of type (R3) are

$$(R3) \quad (B, b)(A, a)(B, b)^{-1} = (A, a)$$

for (A, a) and $(B, b) \in \Omega$ such that $a \notin B$, $b \notin A$, $a^{-1} \notin B$, $b^{-1} \notin A$ and at least one of (a) $A \cap B = \emptyset$ or (b) $b \in \text{lk}_L(a)$ holds. We refer to this relation as (R3)(a) if condition (a) holds and (R3)(b) if condition (b) holds.

The relations of type (R4) are

$$(R4) \quad (B, b)(A, a)(B, b)^{-1} = (A, a)(B - b + a, a)$$

for (A, a) and $(B, b) \in \Omega$ such that $a \notin B$, $b \notin A$, $a^{-1} \notin B$, $b^{-1} \in A$ and at least one of (a) $A \cap B = \emptyset$ or (b) $b \in \text{lk}_L(a)$ holds. We refer to this relation as (R4)(a) if condition (a) holds and (R4)(b) if condition (b) holds.

The relations of type (R5) are

$$(R5) \quad (A - a + a^{-1}, b)(A, a) = (A - b + b^{-1}, a)\sigma_{a,b}$$

for $(A, a) \in \Omega$ and $b \in A$ with $b^{-1} \notin A$, $b \neq a$ and $b \sim a$, where $\sigma_{a,b}$ is the type (1) Whitehead automorphism with $\sigma_{a,b}(a) = b^{-1}$, $\sigma_{a,b}(b) = a$ and which fixes the other generators.

The relations of type (R6) are

$$(R6) \quad \sigma(A, a)\sigma^{-1} = (\sigma(A), \sigma(a))$$

for $(A, a) \in \Omega$ of type (2) and $\sigma \in \Omega$ of type (1).

The relations of type (R7) are

(R7) the entire multiplication table of the type (1) Whitehead automorphisms,

which form a finite subgroup of $\text{Aut } A_\Gamma$.

The relations of type (R8) are

$$(R8) \quad (A, a) = (L - a^{-1}, a)(L - A, a^{-1})$$

for $(A, a) \in \Omega$.

The relations of type (R9) are

$$(R9) \quad (A, a)(L - b^{-1}, b)(A, a)^{-1} = (L - b^{-1}, b)$$

for $(A, a) \in \Omega$ and $b \in L$ with $b, b^{-1} \notin A$.

The relations of type (R10) are

$$(R10) \quad (A, a)(L - b^{-1}, b)(A, a)^{-1} = (L - a^{-1}, a)(L - b^{-1}, b)$$

for $(A, a) \in \Omega$ and $b \in L$ with $b \in A$ and $b^{-1} \notin A$.

Let R be the set of elements of Φ corresponding to all relations of the forms (R1), (R2), (R3)(a), (R3)(b), (R4)(a), (R4)(b), (R5), (R6), (R7), (R9) and (R10).

Note that R is a finite set. Now we can state the following, which we prove in Section 5.

Theorem 2.7 (Restatement of Theorem A) *The group $\text{Aut } A_\Gamma$ has the presentation $\langle \Omega \mid R \rangle$, which is a finite presentation.*

Remark 2.8 Relations (R1), (R2), (R3)(a), (R4)(a) and (R5)–(R10) appear for the automorphism group of the free group in McCool [8, Section 3] (McCool uses reverse composition order for his statements). We have renamed McCool's (R3) as (R3)(a) and (R4) as (R4)(a). Note that in Lyndon–Schupp [7], these relations also appear, but (R7) is unnamed, and (R8)–(R10) are relabeled as (R7)–(R9).

The relations (R3)(b) and (R4)(b) are new here. Our statement of (R4)(a) varies from McCool's by an application of (R2); this allows us to give (R4)(b) as a relation of the same form. Our statement of (R10) varies from McCool's by applications of (R1) and (R2), and our statement of (R9) varies from McCool's by an application of (R1). This restatement should make it easier to see what relations (R9) and (R10) do.

Remark 2.9 For $a \in L$, the automorphism $(L - a^{-1}, a)$ is the inner automorphism given by conjugating by a . Relation (R8) states that (A, a) and $(L - A, a^{-1})$ represent the same element of $\text{Out } A_\Gamma$. Relations (R9) and (R10) are cases of the following familiar fact about groups: if for g in a group G , C_g denotes conjugation by g , then for any $\phi \in \text{Aut } G$, we have $\phi C_g \phi^{-1} = C_{\phi(g)}$.

In the case of the free group, Relations (R8), (R9) and (R10) follow from relations (R1)–(R7). However, in the case of a general right-angled Artin group, this is only true of Relation (R8), which follows immediately from Relations (R1) and (R2). This is why we leave relations of type (R8) out of R . However, we leave Relation (R8) in the list of relations for convenience and to keep with McCool's numbering system. Relations (R9) and (R10) follow from the other relations only if conjugation automorphisms can be factored into products of dominated transvections; this is not possible for general A_Γ .

Proposition 2.10 *For each relation w in any of the classes of relations (R1)–(R10), all the symbols appearing in w denote well-defined Whitehead automorphisms. Furthermore, these relations are true identities in $\text{Aut } A_\Gamma$.*

Proof If w is a relation of type (R3) or (R7), then it is vacuously true that all the terms appearing in w are well defined (since the instances of these relations are indexed over well-defined terms).

If w is a relation of type (R1), (R2), (R5) or (R6), then type (2) Whitehead automorphisms in w are clearly well defined by Lemma 2.5.

If w is a relation of type (R4), then since $b \notin A$ but $b^{-1} \in A$, we know $a \geq b$ (by Lemma 2.5). Since $a \geq b$, every component of $\Gamma - \text{st}(v(b))$ is a union of components of $\Gamma - \text{st}(v(a))$ and elements of $\text{st}(a)$. Then by Lemma 2.5, $(B - b + a, a)$ is well defined.

If w is a relation of type (R5), then we have $a \sim b$ with $a \neq b$, which implies that there is an automorphism π of Γ switching $v(a)$ and $v(b)$ but fixing the other vertices. Then $\sigma_{a,b}$ is the composition of the automorphism of A_Γ induced from π with the inversion of b .

If w is a relation of type (R8), then $(L - A, a^{-1})$ is well defined by Lemma 2.5. Since for any $c \in L$, $(L - c^{-1}, c)$ represents conjugation by c , we know the terms in w are well defined if w is a relation of type (R8), (R9) or (R10).

Each identity can then be verified by computing actions on X . □

Remark 2.11 At this point it is easy to see that $\Omega = \Omega^{-1}$. This is because of Equation (R1) and the fact that the set of type (1) Whitehead automorphisms is closed under taking inverses.

3 The structure of $\text{Aut } A_\Gamma$

3.1 Sorting automorphisms by their scope

Using the special notation for type (2) Whitehead automorphisms, we can restate the definitions of Ω_s and Ω_ℓ more succinctly. A Whitehead automorphism α is in Ω_ℓ if it is of type (1), or if α is of type (2) and we can write $\alpha = (A, a)$ for some A with $A \cap \text{lk}_L(a) = \emptyset$. A Whitehead automorphism α is in Ω_s if $\alpha = (A, a)$ is of type (2) and $A \subset \text{st}_L(a)$.

Whenever we declare an element $(A, a) \in \Omega_\ell$, we will assume that $A \cap \text{lk}_L(a) = \emptyset$. This is necessary since if $x \in \text{lk}_L(a)$, then $(A \cup \{x, x^{-1}\}, a) = (A, a)$.

The goal of this subsection is to prove part (1) of Theorem B. We proceed by describing a series of identities that allow us to rewrite a product of a long-range and a short-range automorphism. We will then show that by a finite number of applications of these identities, we can express any automorphism as a product of a single element of $\langle \Omega_s \rangle$ and a single element of $\langle \Omega_\ell \rangle$.

Lemma 3.1 *Every Whitehead automorphism $\alpha \in \Omega$ has a unique decomposition as a product $\alpha = \alpha_s \alpha_\ell$, where $\alpha_s \in \Omega_s$ and $\alpha_\ell \in \Omega_\ell$.*

Proof If α is of type (1), then $\alpha_s = 1$ and $\alpha_\ell = \alpha$. So assume $\alpha = (A, a)$. Set $A_1 = A \cap \text{st}_L(a)$ and set $A_2 = A - \text{lk}_L(a)$. By Lemma 2.5, both (A_1, a) and (A_2, a) are well defined. So set $\alpha_s = (A_1, a)$ and set $\alpha_\ell = (A_2, a)$. By Equation (R2), we have $\alpha = \alpha_s \alpha_\ell$. This decomposition is unique since $\Omega_s \cap \Omega_\ell = \{1\}$. \square

We call α_s the *short-range part* of α and α_ℓ the *long-range part* of α . Let $\ell: \Omega \rightarrow \Omega_\ell$ be given by $\ell(\alpha) = \alpha_\ell$ and let $s: \Omega \rightarrow \Omega_s$ be given by $s(\alpha) = \alpha_s$.

Definition 3.2 Suppose $\alpha \in \Omega_\ell$ and $\beta = (B, b) \in \Omega_s$. Of course, we may assume that for $x \in \text{lk}_L(b)$, not both x and x^{-1} are in B . Let the *sorting substitution* of $\alpha\beta$ be the word in Ω defined as follows.

If α is given by a permutation of L , then $(\alpha(B), \alpha(b)) \in \Omega_s$, and the substitution is:

$$(3-1) \quad \alpha\beta \mapsto (\alpha(B), \alpha(b))\alpha.$$

Now suppose $\alpha = (A, a)$. If $v(a) = v(b)$, then the substitution is given by:

$$(3-2) \quad \alpha\beta \mapsto \beta\alpha.$$

If $a \in \text{lk}_L(b)$, then we know $b, b^{-1} \notin A$. As we assumed earlier, not both $a \in B$ and $a^{-1} \in B$. The substitution is given by:

$$(3-3) \quad \alpha\beta \mapsto \begin{cases} \beta\alpha & \text{if } a \notin B, a^{-1} \notin B \\ \beta s((A - a + b, b)) \ell((A - a + b, b)) \alpha & \text{if } a \notin B, a^{-1} \in B \\ \beta s((A - a + b^{-1}, b^{-1})) \ell((A - a + b^{-1}, b^{-1})) \alpha & \text{if } a \in B, a^{-1} \notin B. \end{cases}$$

If $a \notin \text{st}_L(b)$, the substitution is given by:

$$(3-4) \quad \alpha\beta \mapsto \beta\alpha.$$

Sublemma 3.3 *Suppose $b \in L$, $c \in \text{lk}_L(b)$ and $b \geq c$. If $a \in L$ and $a \geq b$, then $b \in \text{st}_L(a)$.*

Proof By transitivity, $a \geq c$. Then by the definition of domination, $b \in \text{st}_L(a)$. \square

Lemma 3.4 *All of the elements substituted for $\alpha\beta$ in Definition 3.2 are equal to $\alpha\beta$ in $\text{Aut } A_\Gamma$.*

Proof The substitution in Equation (3-1) is valid by Equation (R6). The substitution in Equation (3-2) is valid by Equation (R2) (if $a = b$, then we know $A \cap B = \{a\}$ and therefore $\alpha\beta = (A \cup B, a) = \beta\alpha$) and Equation (R1) (if $a^{-1} = b$, then $\alpha^{-1} = (L - A - \text{lk}_L(a), a^{-1})$ and Equation (R2) implies α^{-1} and β commute). The first substitution in Equation (3-3) is valid by Equation (R3)(b). The second substitution in Equation (3-3) is valid by Equation (R4)(b), and Equation (R2) to split $(A - a + b, b)$ into long-range and short-range parts. To get the third substitution in Equation (3-3), note that since $\beta \in \Omega_s$, we have $\beta = (\text{st}_L(b) - B, b^{-1})$. Then the third substitution is simply the second substitution, after representing β in a different way.

Now suppose $a \notin \text{st}_L(b)$. By Sublemma 3.3 (assuming β is nontrivial), we know that $a \not\geq b$. Then if $b \notin A$, each element of the component of $v(b)$ in $\Gamma - \text{st}(v(a))$ is fixed by α . Then $A \cap B = \emptyset$ since $\beta \in \Omega_s$. So if $b \notin A$, then Equation (3-4) is valid by Equation (R3). If $b \in A$, we apply Equation (R8) to replace (A, a) by $(L - a^{-1}, a)(L - A, a^{-1})$. Then $b \notin L - A$, so $(L - A, a^{-1})$ commutes with β by Equation (R3), and $(L - a^{-1}, a)$ commutes with β by Equation (R9) ($a, a^{-1} \notin B$ since $\beta \in \Omega_s$). After commuting both automorphisms past β , we recombine them by Equation (R8). \square

Remark 3.5 Equation (3-3) indicates that there are many examples of graphs Γ such that neither $\langle \Omega_\ell \rangle$ nor $\langle \Omega_s \rangle$ is a normal subgroup of $\text{Aut } A_\Gamma$.

Lemma 3.6 *Suppose we have $(A, a) \in \Omega_\ell$. Suppose w is a product (in any order) of long-range automorphisms of the form $\ell((A - a + x, x))$ for $x \in \text{st}_L(a)$ with $x \geq a$, together with short-range automorphisms with multiplier $b^{\pm 1}$. Then we can apply sorting substitutions to rewrite w as a word w' satisfying the same hypotheses as w (for the same (A, a)), and such that all the short-range automorphisms in w' appear to the left of any long-range automorphisms in w' .*

Proof We argue by induction on the number k of long-range automorphisms in w with multipliers other than $b^{\pm 1}$. It is clear that by applying the substitution from Equation (3-2), we can move a short-range automorphism appearing in the word to the left across a long-range automorphism with multiplier $b^{\pm 1}$. In the base case $k = 0$, we only need to move short-range automorphisms across long-range automorphisms with multiplier $b^{\pm 1}$, so we are done. Now suppose $k > 0$. We break up w as $w_1 w_2$, where there is only one long-range automorphism with multiplier not equal to $b^{\pm 1}$ in w_2 , say $\ell((A - a + x, x))$. If $x \notin \text{st}_L(b)$, then by applying the substitution from

Equation (3-4) (and possibly the one from Equation (3-2)), we can move all the short-range automorphisms in w_2 to the left across $\ell((A - a + x, x))$. If $x \in \text{lk}_L(b)$, then by applying the substitution from Equation (3-3), we can move any short-range automorphism to the left across $\ell((A - a + x, x))$. In doing so, we may introduce a new short-range automorphism with multiplier $b^{\pm 1}$ to the left of $\ell((A - a + x, x))$, as well as a new long-range automorphism $\ell((A - a + y, y))$ where $y = b^{\pm 1}$. It is then clear that by applying substitutions from Equation (3-3) and Equation (3-2), we can move all the short-range automorphisms to the left across $\ell((A - a + x, x))$. In either case, we can rewrite w_2 as a word w_3v satisfying the hypotheses of the lemma, where w_3 contains no long-range elements with multipliers other than $b^{\pm 1}$ and v contains no short-range elements. Then by induction, w_1w_3 can be rewritten as a word u satisfying the conclusions of the lemma, and we have rewritten w as a word uv satisfying the conclusions of the lemma. \square

Lemma 3.7 *Suppose $\alpha = (A, a) \in \Omega_\ell$ and $\beta_1, \dots, \beta_k \in \Omega_s$. Then we can apply finitely many sorting substitutions to the word $w_0 = \alpha\beta_k \cdots \beta_1$ to get a word where all of the long-range elements are of the form $\ell((A - a + x, x))$ for various $x \in \text{st}_L(a)$ with $x \geq a$, and all the short-range elements are to the left of any long-range elements.*

Proof We prove the lemma by induction on k . If $k = 0$, it is true. Now assume the lemma holds for $\alpha\beta_k \cdots \beta_2$. Then we rewrite w_0 as $u_1\delta_m \cdots \delta_1\beta_1$, where u_1 is a word in Ω_s and $\delta_i = \ell((A - a + x_i, x_i))$ for some $x_1, \dots, x_m \in \text{st}_L(a)$ with $x_i \geq a$. We apply Lemma 3.6 to the subsequence $\delta_m \cdots \delta_1\beta_1$ and rewrite it as u_2v , where u_2 is a word in Ω_s and v is a product of automorphisms of the form $\ell((A - a + x, x))$ for various $x \geq a$. Then we have rewritten w_0 as u_1u_2v , which is in the desired form. \square

Proof of part (1) of Theorem B We induct on the length of the word w in $\Omega_\ell \cup \Omega_s$ that we wish to rewrite. If $|w| \leq 1$, we are done. Now assume the theorem is true for words of length $|w| - 1$. Let w' be a word in $\Omega_\ell \cup \Omega_s$ and let $\alpha \in \Omega_\ell \cup \Omega_s$ such that $w = \alpha w'$ is a reduced factorization. Then the theorem applies to w' , so $w' = w_s w_\ell$ where w_s is a word in Ω_s and w_ℓ is a word in Ω_ℓ . If $\alpha \in \Omega_s$, we are done, so assume $\alpha \in \Omega_\ell$. If α a type (1) Whitehead automorphism, we can move it across w_s by $|w_s|$ applications of the substitution from Equation (3-1) and we are done. Otherwise, $\alpha = (A, a)$, and we apply Lemma 3.7 to rewrite αw_s as $w'_s w'_\ell$, with w'_s a word in Ω_s and w'_ℓ a word in Ω_ℓ . Then $w = w'_s w'_\ell w_\ell$ and we are done. \square

3.2 The homology representation and short-range automorphisms

Let H_Γ denote the abelianization $H_1(A_\Gamma) \cong \mathbb{Z}^n$ of A_Γ . Since the commutator subgroup of A_Γ is a characteristic subgroup of A_Γ , every automorphism of A_Γ

induces an automorphism of H_Γ . This defines a map $\text{Aut } A_\Gamma \rightarrow \text{Aut } H_\Gamma \cong \text{GL}(n, \mathbb{Z})$, which we call the *homology representation*. For $\alpha \in \text{Aut } A_\Gamma$, we denote its image under the homology representation by $\alpha_* \in \text{Aut } H_\Gamma$. In this section, we prove part (2) of Theorem B, and we examine the structure of $\langle \Omega_s \rangle$.

Lemma 3.8 *Let $\beta \in \langle \Omega_s \rangle$ and let $c \in X$. Then $\text{supp } \beta(c)$ is a clique contained in $\text{st}(c)$.*

Proof We proceed by induction on the length of β as a product of members of Ω_s . If $|\beta| = 0$, the lemma holds. Now suppose that w is a word such that $\text{supp } w$ is a clique contained in $\text{st}(c)$ and that $(B, b) \in \Omega_s$. Of course, we may assume that $B \cap B^{-1} = \emptyset$. If $(B - \{b\}) \cap (\text{supp } w)^{\pm 1} = \emptyset$, then $\beta(w) = w$. So suppose we have $a \in (B - \{b\}) \cap (\text{supp } w)^{\pm 1}$. Then $a \in \text{st}_L(c)$ and by Lemma 2.5, $b \geq a$. Since $\beta \in \Omega_s$, b is adjacent to a , and since $\text{supp } w$ is a clique, it follows from the definition of domination that b is adjacent to every other member of $\text{supp } w$. Since $\text{supp}(B, b)(w) \subset \{v(b)\} \cup \text{supp } w$, this proves the lemma. \square

Proof of part (2) of Theorem B Suppose $\beta \in \langle \Omega_s \rangle$ and $\beta_* \in \text{Aut } H_\Gamma$ is trivial. Then for any $a \in X$, it follows from Lemma 3.8 that we can commute the elements of $\text{supp } \beta(a)$ past each other. But since β_* is trivial, the sum exponent of a in $\beta(a)$ is 1, and the sum exponent of any other x in $\beta(a)$ is 0. So $\beta(a) = a$ for any a , and β is trivial. \square

Now we examine the structure of $\langle \Omega_s \rangle$. The argument below tells us the structure of the images of many subgroups of $\text{Aut } A_\Gamma$ under the homology representation, so we phrase it in greater generality. In particular, we prove an intermediate result that is quoted in the sequel to the current paper [4].

Let \leq' be a transitive, reflexive relation on X , such that $a \leq' b$ implies $a \leq b$ for $a, b \in X$ (for example, the adjacent domination relation). Write $a \sim' b$ when $a \leq' b$ and $b \leq' a$; then \sim' is an equivalence relation on X . Let $G < \text{Aut } H_\Gamma$ be generated by $\{(\tau_{a,b})_* \mid a \geq' b\}$. Let $C_1 \cup \dots \cup C_m = X$ be the \sim' -classes of X . Let $N = \{(\tau_{a,b})_* \mid a, b \in X, a \geq' b \text{ and } a \not\sim' b\}$ and for each $i = 1, \dots, m$, let $G_i = \{(\tau_{a,b})_* \mid a, b \in C_i\}$.

Proposition 3.9 *The group N is nilpotent, each $G_i \cong \text{SL}(|C_i|, \mathbb{Z})$, and the inclusion maps of N and the G_i into G give the decomposition*

$$(3-5) \quad G \cong (G_1 \times \dots \times G_m) \ltimes N.$$

Proof We can pick an indexing $X = \{x_1, \dots, x_n\}$ such that if $x_i \geq' x_j$ with $x_i \not\sim' x_j$, then $j > i$. Taking the image of X in H_Γ to be an ordered basis under this indexing, the homology representation takes the transvection τ_{x_i, x_j} to the elementary matrix $E_{i,j}$. Then the group N is then a group of upper-triangular unipotent matrices and is therefore nilpotent.

We make a further demand on our indexing of X : if $i < j < k$ and $x_i \sim' x_k$, then $x_j \sim' x_i$. Under such an indexing, the elements of C_i are an unbroken string of elements of X , so say $C_i = \{x_{r_i}, \dots, x_{s_i}\}$. So G_i is generated by the elementary matrices $E_{j,k}$ with $r_i \leq j, k \leq s_i$; in particular, it is an embedded copy of $\text{SL}(|C_i|, \mathbb{Z})$.

The generators of G_i and G_j commute for $i \neq j$ by Equation (R3), so the subgroup generated by the $\{G_i\}_i$ is a direct product.

It is obvious that N and G_1, \dots, G_m generate G . Now suppose that $(\tau_{a,b})_*$ is a generator of N and $(\tau_{c,d})_*$ is a generator of one of the G_i . Then $(\tau_{c,d})_*^{-1}(\tau_{a,b})_*(\tau_{c,d})_*$ is an element of N ; if $d \neq b$, this follows from Equations (R3) and (R4), and if $d = b$, then this element is $(\tau_{a,b})_*$ since we are working in H_Γ . Since N is normal, we get the decomposition of G in Equation (3-5). \square

Proposition 3.10 *The group G has a presentation in which the generators S_G are the row operations $E_{a,b} = (\tau_{a,b})_*$ for $a, b \in X$ with $a \geq' b$, and with the relations R_G being all relations among the S_G of the following forms:*

- (1) $[E_{a,b}, E_{c,d}] = 1$ if $b \neq c$ and $a \neq d$.
- (2) $[E_{a,b}, E_{b,d}]E_{a,d}^{-1} = 1$ if $a \neq d$.
- (3) $(E_{a,b}E_{b,a}^{-1}E_{a,b})^4 = 1$, if $a \sim' b$ and $a \neq b$.
- (4) $(E_{a,b}E_{b,a}^{-1}E_{a,b})^2(E_{a,b}E_{b,a}^{-1}E_{a,b}E_{b,a})^{-3} = 1$, if $a, b \in C_i$, $a \neq b$, for some i with $|C_i| = 2$.

Proof By Proposition 3.9, each $G_i \cong \text{SL}(|C_i|, \mathbb{Z})$ and N is a nilpotent group. For each i , it follows from classical presentations for $\text{SL}(n, \mathbb{Z})$ that G_i has a presentation with generators $S_G \cap G_i$ and whose relations are those elements of R_G only involving the generators in $S_G \cap G_i$ (see Milnor [10, Corollary 10.3] for $n \geq 3$ and Serre [12, Example I.4.2(c)] for $n = 2$). Since N is a unipotent matrix group and $S_G \cap N$ is a generating set for N that is closed under taking commutators, we know that N has a presentation with generators $S_G \cap N$ and whose relations are those elements of R_G only involving generators in $S_G \cap N$.

Relation (1) implies that the group generated by $S_G \cap (G_1 \cup \dots \cup G_m)$ subject to these relations is isomorphic to the product $G_1 \times \dots \times G_m$, and Relation (1) and

Relation (2) encode the semidirect product action of $G_1 \times \cdots \times G_m$ on N , so that the group $\langle S_G \mid R_G \rangle \cong G$. \square

Corollary 3.11 *Suppose \tilde{G} is a subgroup of $\text{Aut } A_\Gamma$ generated by a set of dominated transvections. Then there is a relation \leq' such that the image of \tilde{G} under the homology representation is the group G generated by $\{(\tau_{a,b})_* \mid a \geq' b\}$ and the conclusions of Proposition 3.9 and Proposition 3.10 hold for G . In particular, the group $\langle \Omega_S \rangle$ has a decomposition of the form in Equation (3-5) and a presentation of the form given in Proposition 3.10.*

Proof Let S be a set of dominated transvections such that $\tilde{G} = \langle S \rangle$. Let \leq' be the reflexive relation that is the transitive closure of \leq'' , where $a \geq'' b$ whenever $\tau_{a,b} \in S$. Then \tilde{G} is generated by $\{\tau_{a,b} \mid a \geq' b\}$, and its image under the homology representation is the group G generated by $\{(\tau_{a,b})_* \mid a \geq' b\}$. Then we can apply Proposition 3.9 and Proposition 3.10. If $S = \Omega_S$, then by part (2) of Theorem B, the homology representation restricted to $\langle \Omega_S \rangle = \tilde{G}$ maps isomorphically to G . So the conclusions of Proposition 3.9 and Proposition 3.10 apply to $\langle \Omega_S \rangle$, where \leq' is adjacent domination. \square

3.3 Peak-reducing products of long-range automorphisms

The goal of this subsection is to prove part (3) of Theorem B. Our proof is similar to the proof of the peak reduction theorem for free groups given in Higgins–Lyndon [5].

Let $k \geq 1$. For a k -tuple $W = (w_1, \dots, w_k)$ of cyclic words, we denote the k -tuple of conjugacy classes by $[W] = ([w_1], \dots, [w_k])$. We proceed by putting the definition of peak reduction from the introduction in context.

Definition 3.12 Suppose $\alpha, \beta \in \Omega$ and $[W]$ is a k -tuple of conjugacy classes in A_Γ . Then $\beta\alpha$, the word of length 2, is called a *peak* with respect to $[W]$ if

$$|\alpha \cdot [W]| \geq |[W]|$$

and

$$|\alpha \cdot [W]| \geq |\beta\alpha \cdot [W]|$$

and at least one of these inequalities is strict.

Suppose $\gamma \in \text{Aut } A_\Gamma$ and we have a factorization $\gamma = \alpha_k \cdots \alpha_1$ with $\alpha_1, \dots, \alpha_k$ in Ω . We say α_i is a *peak* of this factorization, with respect to $[W]$, if $1 < i < k$ and $\alpha_{i+1}\alpha_i$ is a peak with respect to $(\alpha_{i-1} \cdots \alpha_1) \cdot [W]$. The *height* of a peak α_i is simply $|(\alpha_i \cdots \alpha_1) \cdot [W]|$.

Then the factorization $\gamma = \alpha_k \cdots \alpha_1$ is peak-reduced with respect to $[W]$ (as defined in the introduction) if and only if it has no peaks with respect to $[W]$.

It is important to note that for general right-angled Artin groups, the automorphism group $\text{Aut } A_\Gamma$ does not act on the set of graphically reduced words; rather it only acts on the set of group elements. This means we need to take care to distinguish words from the elements they represent. These measures were unnecessary in the original proof for free groups, since for a free group the set of reduced words is the set of group elements.

Definition 3.13 If $(A, a) \in \Omega_\ell$ and w is a graphically reduced cyclic word, define the *obvious representative* of $(A, a)([w])$ based on w to be the cyclic word w' gotten from w by the following replacements:

- For every subsegment of w of the form buc^{-1} or cub^{-1} , with u any word in $\text{lk}_L(a)$, $b \in A - a$ and $c \in L - A - \text{lk}_L(a)$, replace this subsegment with $bauc^{-1}$ or $cua^{-1}b^{-1}$ respectively in w' .
- For every subsegment of w of the form bua^{-1} or aub^{-1} , with u any word in $\text{lk}_L(a)$ and $b \in A - a$, replace this subsegment with bu or ub^{-1} respectively in w' .

The *obvious representative* of $(A, a) \cdot [W]$ based on W is the k -tuple (w'_1, \dots, w'_k) , where each w'_i is the obvious representative of $(A, a)([w_i])$ based on w_i .

Claim 3.14 *The obvious representative w' of $(A, a)([w])$ based on w is a graphically reduced representative of $(A, a)([w])$.*

Proof First we show that w' is graphically reduced. Note that we have only added or removed instances of $a^{\pm 1}$. Since w is graphically reduced, w' can only fail to be graphically reduced on subsegments where we added or removed instances of $a^{\pm 1}$. Those replacements that introduce an instance of $a^{\pm 1}$ introduce it in a way where it cannot cancel (a does not commute with $b \in A - a$ or with $c \in L - A - \text{lk}_L(a)$). Suppose a replacement that removes an instance of $a^{\pm 1}$ results in w' not being graphically reduced. Then we have a subsegment dvd^{-1} of w being replaced by $dv'd^{-1}$ in w' , where $d \in L - \text{st}_L(a)$, v' is a word in $\text{lk}_L(d)$, and v is v' with some instances a or a^{-1} inserted. Then dvd^{-1} contains an instance of some $b \in (A - a)^{\pm 1}$. If $a \not\geq b$, then $v(b)$ is in a component of $\Gamma - \text{st}(v(a))$ that is conjugated by (A, a) . It follows that for each x in $\text{supp } dv'd^{-1}$, either $x \in \text{st}_L(a)$ or $v(x)$ is in the same component of $\Gamma - \text{st}(v(a))$ as $v(b)$. Then $v = v'$, a contradiction. So suppose $a \geq b$. If $v(b) \neq v(d)$, then d commutes with a , a contradiction. If $v(b) = v(d)$, then $\text{supp } v \subset \text{st}(v(a))$,

and our substitutions never would have removed instances of $a^{\pm 1}$ from v . So w' is graphically reduced.

Observe that w' represents $(A, a)([w])$: because (A, a) is long-range, it has no effect on the subsegments of w that are words in $\text{lk}_L(a)$, and the substitutions in the definition of w' capture all those changes that (A, a) makes to w that do not cancel each other out. \square

Definition 3.15 Let w be a graphically reduced cyclic word and let $a \in L$. Then for $b, c \in L - \text{lk}_L(a)$, we define the *adjacency counter* of w relative to a , written as $\langle b, c \rangle_{w, a}$, to be the number of subsegments of w of the form $(buc^{-1})^{\pm 1}$, where u is any (possibly empty) word in $\text{lk}_L(a)$.

For a k -tuple of graphically reduced cyclic words $W = (w_1, \dots, w_k)$, define the adjacency counter of W relative to a as:

$$\langle b, c \rangle_{W, a} = \sum_{i=1}^k \langle b, c \rangle_{w_i, a}$$

For $B, C \subset L$, we define:

$$\langle B, C \rangle_{W, a} = \sum_{b \in (B - \text{lk}_L(a))} \sum_{c \in (C - \text{lk}_L(a))} \langle b, c \rangle_{W, a}$$

For $\alpha \in \text{Aut } A_\Gamma$, we define

$$D_{[W]}(\alpha) = |\alpha \cdot [W]| - |[W]|.$$

When W is clear, we leave it out, writing $\langle B, C \rangle_a$ and $D(\alpha)$.

With W and a as above, note that for any $B, C \subset L$, the number $\langle B, C \rangle_a \geq 0$. Further, we have $\langle B, C \rangle_a = \langle C, B \rangle_a$. If $D \subset L$ with $D \cap C = \emptyset$, then we have:

$$\langle B, C + D \rangle_a = \langle B, C \rangle_a + \langle B, D \rangle_a$$

Also note that $\langle a, a \rangle_a = 0$ (since each w_i is graphically reduced).

Lemma 3.16 If W is a k -tuple of graphically reduced cyclic words, $(A, a) \in \Omega_\ell$, and W' is the obvious representative of $(A, a) \cdot [W]$, then:

$$D_{[W]}((A, a)) = |W'| - |W| = \langle A - a, L - A \rangle_{W, a} - \langle a, A - a \rangle_{W, a}$$

Proof This is immediate from counting the letters removed and added in the definition of W' . \square

Note that if W' and W are different k -tuples of graphically reduced cyclic words representing the same k -tuple of conjugacy classes, we may have different adjacency counters with respect to W and W' . But the function $D_{[W]}$ depends only on $[W]$.

Lemma 3.17 *Let W be a k -tuple of graphically reduced cyclic words. If $(A, a) \in \Omega_\ell$, then:*

$$D_{[W]}((A, a)) = \langle A, L - A \rangle_{W, a} - \langle a, L \rangle_{W, a}$$

Proof From Lemma 3.16:

$$\begin{aligned} D((A, a)) &= \langle A - a, L - A \rangle_a - \langle a, A - a \rangle_a \\ &= \langle A, L - A \rangle_a - (\langle a, L - A \rangle_a + \langle a, A - a \rangle_a + \langle a, a \rangle_a) \\ &= \langle A, L - A \rangle_a - \langle a, L \rangle_a \end{aligned} \quad \square$$

The following lemma is the machine that makes peak reduction possible. This is an extension of a parallel lemma for free groups that appears in Higgins–Lyndon [5].

Lemma 3.18 *Suppose $\alpha, \beta \in \Omega_\ell$ and $[W]$ is a k -tuple of conjugacy classes of A_Γ . If $\beta\alpha^{-1}$ forms a peak with respect to $[W]$, there exist $\delta_1, \dots, \delta_k \in \Omega_\ell$ such that $\beta\alpha^{-1} = \delta_k \cdots \delta_1$ and for each $i, 1 \leq i < k$, we have*

$$|(\delta_i \cdots \delta_1) \cdot [W]| < |\alpha^{-1} \cdot [W]|.$$

A factorization of $\beta\alpha^{-1}$ is *peak-lowering* if it satisfies the conclusions of the lemma, so Lemma 3.18 states that every peak has a peak-lowering factorization. Such a factorization might not be peak-reduced, but the height of its highest peak is lower than the height of the peak in $\beta\alpha^{-1}$. We postpone the proof of Lemma 3.18 to show how it implies part (3) of Theorem B.

Proof of part (3) of Theorem B Let $\gamma \in \langle \Omega_\ell \rangle$ and write $\gamma = \alpha_k \cdots \alpha_1$ with $\alpha_1, \dots, \alpha_k \in \Omega_\ell$. Let h be

$$h = \sup \{ |(\alpha_i \cdots \alpha_1) \cdot [W]| \mid \alpha_i \text{ is a peak} \}$$

which is the height of the highest peak in the factorization, and let m be the number of maximal-height steps between peaks

$$m = \left| \left\{ i \mid h = |(\alpha_i \cdots \alpha_1) \cdot [W]| \text{ and } \alpha_i \text{ is a peak or between two peaks} \right\} \right|.$$

If the factorization is not peak-reduced, then there is a peak α_i of maximal height. Apply Lemma 3.18 to the peak $\alpha_{i+1}\alpha_i$ with respect to $(\alpha_{i-1} \cdots \alpha_1) \cdot W$ to get

$$\alpha_{i+1}\alpha_i = \delta_j \cdots \delta_1$$

satisfying the conclusions of the lemma, and therefore a new factorization of γ :

$$\gamma = \alpha_k \cdots \alpha_{i+2} \delta_j \cdots \delta_1 \alpha_{i-1} \cdots \alpha_1.$$

If m of the old factorization was not 1, then m of the new factorization is one less. If m of the old factorization was 1, then h of the new factorization is strictly lower than h of the old factorization. By repeating this process, we eventually obtain a factorization with $h < 1$. This can only mean that there are no peaks, so we have a factorization which is peak-reduced. \square

Sublemma 3.19 *Let α , β and $[W]$ be as in Lemma 3.18. Then we have*

$$(3-6) \quad 2|\alpha^{-1} \cdot [W]| > |[W]| + |\beta\alpha^{-1} \cdot [W]|.$$

Proof Since $\beta\alpha^{-1}$ is a peak with respect to $[W]$, we can sum the two inequalities in the definition of a peak; by the fact that one of them is strict, we obtain this new inequality. \square

Sublemma 3.20 *Suppose we have $(A, a), (B, b) \in \Omega_\ell$ with $a \notin B$ and $v(a)$ not adjacent to $v(b)$ in Γ (possibly $a = b^{-1}$). Then $\text{lk}_L(a) \cap B = \emptyset$.*

Proof If $x \in \text{lk}_L(a) \cap B$, then $x \in B$ and by Lemma 2.5, either $b \geq x$ or (B, b) acts on the connected component of $v(x)$ in $\Gamma - \text{st}(v(b))$ by conjugation. If the latter were true, since $v(a)$ is adjacent to $v(x)$ and not $v(b)$, we would have that $a \in B$, a contradiction. So $b \geq x$, in which case $v(a)$ is adjacent to $v(b)$, a contradiction. \square

Sublemma 3.21 *Suppose α , β and $[W]$ are as in Lemma 3.18, and also that $\alpha = (A, a)$, $\beta = (B, b)$, and that either $e(a, b)$ or that $A \cap B = \emptyset$ with $a^{-1} \notin B$. Then $|\beta \cdot [W]| < |\alpha^{-1} \cdot [W]|$.*

Proof Take W' to be a representative of $\alpha^{-1} \cdot [W]$ and take W to be the obvious representative of $\alpha \cdot [W']$ based on W' (this doesn't change our original $[W]$). We will show the sublemma by analyzing adjacency counters. First we claim that:

$$(3-7) \quad \langle B, L - B \rangle_{W, b} \geq \langle B, L - B \rangle_{W', b}$$

We will show this by showing that every subsegment of W' (meaning a subsegment of an element of W') that is counted by the adjacency counter on the right above is also counted by the one on the left. So suppose $(cud^{-1})^{\pm 1}$ is counted on the right in Equation (3-7), so that $(cud^{-1})^{\pm 1}$ is a subsegment of W' with $c \in B$, $d \in L - B - \text{lk}_L(b)$ and with u a word in $\text{lk}_L(b)$.

If $e(a, b)$, then since $\beta \in \Omega_\ell$, we know $a^{\pm 1} \notin B$. This means that $v(c) \neq v(a)$. Since $a \in \text{lk}_L(b)$, we also have $v(d) \neq v(a)$. This means that the corresponding subsegment of W is $(cu'd^{-1})^{\pm 1}$, where u' is u , possibly with instances of $a^{\pm 1}$ added or removed. This subsegment is then counted by the counter on the right.

If $A \cap B = \emptyset$, then $a \notin B$ and $b \notin A$. Since $a^{\pm 1} \notin B$, we know that $v(c) \neq v(a)$. By Sublemma 3.20, we know $\text{lk}_L(b) \cap A = \emptyset$. Note that it is possible that in passing to W , this letter d could be deleted by a copy of a^{-1} added to its left if $d = a^{-1}$. We consider this case separately.

First suppose d is not deleted. Then the subsegment of W corresponding to $(cud^{-1})^{\pm 1}$ is either $(cud^{-1})^{\pm 1}$ or $(cua^{-1}d^{-1})^{\pm 1}$, depending on whether $d \in A$ or not. In either case, this subsegment is counted once by the counter on the left, in the second case because $a \in L - B - \text{lk}_L(b)$.

If d is deleted, then $d = a^{-1}$, and our original $(cud^{-1})^{\pm 1} = (cua)^{\pm 1}$ was part of a subsegment $(cuavf^{-1})^{\pm 1}$, where $f \in (A - a)$ and v is a word in $\text{lk}_L(a)$. Note that by Sublemma 3.20, we know $\text{lk}_L(a) \cap B = \emptyset$, so the counter on the right counts this segment only once. The subsegment of W corresponding to this $(cuavf^{-1})^{\pm 1}$ is then $(cuvf^{-1})^{\pm 1}$. Write $v = v'v''$ where v' is the maximal initial segment of v that is a word in $\text{lk}_L(b) \cap \text{lk}_L(a)$ and let f' be the unique letter such that $cuv'(f')^{-1}$ is an initial segment of $cuvf^{-1}$. Either $f' = f$ or $f' \in (\text{lk}_L(a) - \text{lk}_L(b))$. In either case, $f' \in L - B - \text{lk}_L(b)$ (since $A \cap B = \emptyset$ and $B \cap \text{lk}_L(a) = \emptyset$), so the corresponding subsegment of W is counted once by the counter on the left. This shows Equation (3-7).

Now we will show:

$$(3-8) \quad \langle b, L \rangle_{W', b} \geq \langle b, L \rangle_{W, b}$$

Suppose $(bud^{-1})^{\pm 1}$ is a subsegment of W counted by the counter on the right above, so $d \in L - \text{lk}_L(b)$ and u is a word in $\text{lk}_L(b)$. If $e(a, b)$, then W came from a subsegment $(bu'd^{-1})^{\pm 1}$ of W' , where u' is u , possibly with an instance of $a^{\pm 1}$ added or removed; this subsegment is counted by the counter on the left.

If $A \cap B = \emptyset$, then either d originally appeared in W or it was added in passing to W' . If it originally appeared in W , then $(bud^{-1})^{\pm 1}$ came from either a $(bud^{-1})^{\pm 1}$ or a $(bu'au''d^{-1})^{\pm 1}$, where in the second case u'' is the maximal terminal segment of u that is a word in $\text{lk}_L(a) \cap \text{lk}_L(b)$; this subsegment of W' is counted by the counter on the left. If it was added, our $(bud^{-1})^{\pm 1}$ in W is part of a $(buavf^{-1})^{\pm 1}$, with v a word in $\text{lk}_L(a)$ and $f \in A$. This subsegment is counted only once and came from a subsegment $(buvf^{-1})^{\pm 1}$ that is counted once (for similar reasons as above). This shows Equation (3-8).

From Lemma 3.17, Equation (3-7) and Equation (3-8), we see that

$$D_{[W]}(\beta) \geq D_{\alpha^{-1} \cdot [W]}(\beta).$$

By the definition of D , this means that $\|[W]\| + |\beta\alpha^{-1} \cdot [W]| \geq |\alpha^{-1} \cdot [W]| + |\beta \cdot [W]|$. Combining this with Equation (3-6), we obtain

$$2|\alpha^{-1} \cdot [W]| > |\alpha^{-1} \cdot [W]| + |\beta \cdot [W]|$$

which immediately implies the sublemma. □

Proof of Lemma 3.18 If we have $\alpha = (A, a)$ and $\beta = (B, b)$, we will set $A' = L - A - \text{lk}_L(a)$ and $B' = L - B - \text{lk}_L(b)$. Let $\bar{\alpha} = (A', a^{-1})$ and $\bar{\beta} = (B', b^{-1})$. By Equation (R8) and the fact that α and β are long-range, these automorphisms describe the same elements of $\text{Out } A_\Gamma$, and therefore $\alpha^{-1} \cdot [W] = \bar{\alpha}^{-1} \cdot [W]$ and $\beta\alpha^{-1} \cdot [W] = \bar{\beta}\bar{\alpha}^{-1} \cdot [W]$. We claim that if the lemma holds with α or β replaced with $\bar{\alpha} = (A', a^{-1})$ or $\bar{\beta} = (B', b^{-1})$, respectively, then it holds as originally stated. Suppose $\delta_k \cdots \delta_1$ is a peak-lowering factorization of $\bar{\beta}\bar{\alpha}^{-1}$ (for example). By Equation (R8), the element $\beta\alpha^{-1}(\bar{\beta}\bar{\alpha}^{-1})^{-1}$ is the conjugation $(L - b^{-1}, b)$ (which is in Ω_ℓ). If $|\beta\alpha^{-1} \cdot [W]| < |\alpha \cdot [W]|$ then

$$\beta\alpha^{-1} = (L - b^{-1}, b)\delta_k \cdots \delta_1$$

is a peak-lowering factorization of $\beta\alpha^{-1}$, since $(L - b^{-1}, b)$ does not change the length of any conjugacy class. Otherwise $|W| < |\alpha \cdot [W]|$. Again by Equation (R8), $\bar{\beta}\bar{\alpha}^{-1}$ is the conjugation $(L - b, b^{-1})$. So $(\bar{\beta}\bar{\alpha}^{-1})^{-1}\beta\alpha^{-1}$ is $\alpha(L - b, b^{-1})\alpha^{-1}$. If $b \notin A$, then by Equation (R9) and Equation (R10), we know $(\bar{\beta}\bar{\alpha}^{-1})^{-1}\beta\alpha^{-1}$ is a product of conjugations. If $b \in A$, then by Equation (R8), we know $(\bar{\beta}\bar{\alpha}^{-1})^{-1}\beta\alpha^{-1}$ is $(L - a^{-1}, a)\bar{\alpha}(L - b, b^{-1})\bar{\alpha}^{-1}(L - a, a^{-1})$, which is then a product of conjugations by Equation (R9) and Equation (R10). In any case, we have a product of conjugations $\gamma'_j \cdots \gamma'_1$ equal to $(\bar{\beta}\bar{\alpha}^{-1})^{-1}\beta\alpha^{-1}$; then

$$\beta\alpha^{-1} = \delta_k \cdots \delta_1 \gamma'_j \cdots \gamma'_1$$

is a peak-lowering factorization of $\beta\alpha^{-1}$, since conjugations do not change the length of conjugacy classes. So we may swap out $\bar{\alpha}$ for α and $\bar{\beta}$ for β as needed in the proof of this lemma. Also, by the symmetry in the definition of a peak, we may switch α and β if needed.

We fix a k -tuple of graphically reduced cyclic words W representing the conjugacy class $[W]$. Throughout this proof, W' will denote the obvious representative of $\alpha^{-1} \cdot [W]$ based on W . We break this proof down into several cases.

Case 1 α is a type (1) Whitehead automorphism. Then $|\alpha \cdot [W]| = |[W]|$. Since (β, α) is a peak, β must shorten $\alpha \cdot [W]$, so $\beta = (B, b)$ for some (B, b) . Then the automorphism $(\alpha^{-1}(B), \alpha^{-1}(b)) \in \Omega_\ell$ is well defined, and by Equation (R6) the following factorization is peak-lowering:

$$\beta\alpha^{-1} = \alpha^{-1}(\alpha^{-1}(B), \alpha^{-1}(b)).$$

In the remaining cases, we assume that $\alpha = (A, a)$ and $\beta = (B, b)$. We will implicitly use Equation (R1) to write α^{-1} as $(A - a + a^{-1}, a^{-1})$ in the following.

Case 2 $a \in \text{lk}_L(b)$. Of course, this implies that $v(a) \neq v(b)$. Since α and β are long-range, we know that $a, a^{-1} \notin B$ and $b, b^{-1} \notin A$. Then by Equation (R3)(b), we have

$$\beta\alpha^{-1} = (B, b)(A - a + a^{-1}, a^{-1}) = (A - a + a^{-1}, a^{-1})(B, b) = \alpha^{-1}\beta.$$

By Sublemma 3.21, we know $|\beta \cdot [W]| < |\alpha^{-1} \cdot [W]|$, so this factorization is peak-lowering.

Case 3 $A \cap B = \emptyset$ and $a \notin \text{lk}_L(b)$. We will break into subcases according to the configuration of a^{-1} and b^{-1} .

Subcase 3a $v(a) = v(b)$. Since $A \cap B = \emptyset$, this implies that $a = b^{-1}$. By Equation (R2), the following factorization is peak-lowering:

$$\beta\alpha^{-1} = (B, b)(A - a + b, b) = (A + B + b, b).$$

Subcase 3b $a^{-1} \notin B$. If $b^{-1} \notin A$, then

$$\beta\alpha^{-1} = (B, b)(A - a + a^{-1}, a^{-1}) = (A - a + a^{-1}, a^{-1})(B, b)$$

by Equation (R3)(a). If $b^{-1} \in A$, then by Equation (R2) and Equation (R4)(a), we have

$$\beta\alpha^{-1} = (B, b)(A - a + a^{-1}, a^{-1}) = (A + B - b - a + a^{-1}, a^{-1})(B, b).$$

In either case, by Sublemma 3.21, $|\beta \cdot [W]| < |\alpha^{-1} \cdot [W]|$, so these factorizations are peak-lowering.

Subcase 3c $v(a) \neq v(b)$, $a^{-1} \in B$ and $b^{-1} \in A$. Note that since we are allowed to switch α and β , if $v(a) \neq v(b)$ and either $a^{-1} \notin B$ or $b^{-1} \notin A$, we are in subcase 3b. Therefore this subcase finishes case 3. Since $a^{-1} \in B$ and $a \notin B$, we see from Lemma 2.5 that $b \geq a$. Similarly, $a \geq b$. So $a \sim b$ and by Lemma 2.5, the automorphisms $\alpha' = (A, b^{-1})$ and $\beta' = (B, a^{-1})$ are well defined.

In the rest of this case, all adjacency counting is done with respect to W' . Since $a \sim b$ and $v(a)$ is not adjacent to $v(b)$ in Γ , note that $\text{lk}_L(a) = \text{lk}_L(b)$ and therefore the

adjacency counters with respect to a and b are the same functions. So by Lemma 3.17,

$$D(\alpha) + D(\beta) = D(\alpha') + D(\beta').$$

Also, by the definition of D and Equation (3-6),

$$D(\alpha) + D(\beta) = -(2|\alpha^{-1} \cdot [W]| - |[W]| - |\beta\alpha^{-1} \cdot [W]|) < 0.$$

So either $D(\alpha') < 0$ or $D(\beta') < 0$. Since we may swap α and β (which swaps α' and β'), we assume $D(\beta') < 0$.

Now we will find our peak-lowering factorization. Let $\sigma_{a,b}$ be the type (1) Whitehead automorphism from Equation (R5). By Equation (R5), we have:

$$\beta(\beta')^{-1} = (B, b)(B - a^{-1} + a, a) = (B - a^{-1} + a - b + b^{-1}, a)\sigma_{a,b}$$

By Equation (R2),

$$\beta'\alpha^{-1} = (B, a^{-1})(A - a + a^{-1}, a^{-1}) = (A + B - a, a^{-1}).$$

Then using $\beta\alpha^{-1} = \beta(\beta')^{-1}\beta'\alpha^{-1}$, we have the factorization:

$$\beta\alpha^{-1} = (B - a^{-1} + a - b + b^{-1}, a)\sigma_{a,b}(A + B - a, a^{-1}).$$

To show this factorization is peak-lowering, note the following:

$$|(A + B - a, a^{-1}) \cdot [W]| = |\beta'\alpha^{-1} \cdot [W]| = D(\beta') + |\alpha^{-1} \cdot [W]| < |\alpha^{-1} \cdot [W]|.$$

This is because $D(\beta') = |\beta'\alpha^{-1} \cdot [W]| - |\alpha^{-1} \cdot [W]|$. Then since $\sigma_{a,b}$ does not change the length of a conjugacy class, this factorization is peak-lowering and we are done with this case.

Case 4 $A \cap B \neq \emptyset$ and $a \notin \text{lk}_L(b)$. All adjacency counting in this case is done with respect to W' . First we show we can assume that we are in one of two subcases: either $a \notin B$ and $b \notin A$, or else $a \notin B$, $a^{-1} \in B$, $b \in A$ and $b^{-1} \notin A$.

Possibly by replacing β with $\bar{\beta}$, we may assume that $a \notin B$. Then if $b \notin A$, then we are in the first subcase, so suppose $b \in A$. First suppose $a^{-1} \in B$; if $b^{-1} \notin A$, then we are in the second subcase, and if $b^{-1} \in A$, we can get to the first subcase by swapping both α with $\bar{\alpha}$ and β with $\bar{\beta}$. Otherwise $a^{-1} \notin B$, and swapping α with $\bar{\alpha}$ puts us in the first subcase.

In both of these subcases we will find that $\alpha^{-1} \cdot [W]$ is shortened by a well-defined Whitehead automorphism $(C \cap D, c)$, where C is A or A' , D is B or B' , and c is an element of $\{a, a^{-1}, b^{-1}, b\} \cap A \cap C$. By swapping α with $\bar{\alpha}$ if necessary, we assume $C = A$; similarly, we assume that $D = B'$. Then c is a or b^{-1} , if it is b^{-1} , we swap α with β , α with $\bar{\alpha}$ and β with $\bar{\beta}$ to get $(C \cap D, c) = (A \cap B', a)$.

Then in any event, we may assume that $(A \cap B', a)$ shortens $\alpha^{-1} \cdot [W]$. We deduce from Lemma 2.5 that $(A - B' + a, a)$ is a well defined Whitehead automorphism. From Equation (R2) we have

$$\alpha = (A - B' + a, a)(A \cap B', a).$$

Then we factor:

$$\beta\alpha^{-1} = \beta(A \cap B', a)^{-1}(A - B' + a, a)^{-1}.$$

Since $(A \cap B', a)$ shortens $\alpha^{-1} \cdot [W]$, we know that

$$|(A - B' + a, a)^{-1} \cdot [W]| < |\alpha^{-1} \cdot [W]|$$

and that $\beta(A \cap B', a)^{-1}$ is a peak with respect to $(A - B' + a, a)^{-1} \cdot [W]$. Then we can apply case 3 of this lemma to the peak $\beta(A \cap B', a)^{-1}$, and obtain a peak-lowering factorization of our original peak.

Subcase 4a $a \notin B$, $a^{-1} \in B$, $b \in A$ and $b^{-1} \notin A$. Then $a \sim b$ by Lemma 2.5, and since $v(a)$ is not adjacent to $v(b)$, we have $\text{lk}_L(a) = \text{lk}_L(b)$. Then adjacency counters taken with respect to a and b are the same. Let $\gamma_1 = (A \cap B, b)$, $\gamma_2 = (A \cap B', a)$, $\gamma_3 = (A' \cap B, a^{-1})$ and $\gamma_4 = (A' \cap B', b^{-1})$. Since $a \sim b$, these $\gamma_1, \gamma_2, \gamma_3$ and γ_4 are all well defined by Lemma 2.5.

Now we will show that one of these automorphisms shortens $\alpha^{-1} \cdot [W]$. Apply Lemma 3.17 twice to get:

$$D(\alpha) + D(\beta) = \langle A, A' \rangle_a - \langle a, L \rangle_a + \langle B, B' \rangle_b - \langle b, L \rangle_b$$

Then by further applications of Lemma 3.17, we obtain:

$$\begin{aligned} \sum_{i=1}^4 D(\gamma_i) &= \langle A \cap B, A' \cup B' \rangle_b - \langle b, L \rangle_b + \langle A' \cap B, A \cup B' \rangle_{a^{-1}} - \langle a^{-1}, L \rangle_{a^{-1}} \\ &\quad + \langle A \cap B', A' \cup B \rangle_a - \langle a, L \rangle_a + \langle A' \cap B', A \cup B \rangle_{b^{-1}} - \langle b^{-1}, L \rangle_{b^{-1}} \end{aligned}$$

Putting these together, it follows from the additivity of adjacency counters that

$$2(D(\alpha) + D(\beta)) = \sum_{i=1}^4 D(\gamma_i) + 2(\langle A \cap B, A' \cap B' \rangle_a + \langle A' \cap B, A \cap B' \rangle_a).$$

Then by Equation (3-6),

$$0 > 2(D(\alpha) + D(\beta)) \geq \sum_{i=1}^4 D(\gamma_i).$$

This shows that for some i , $D(\gamma_i) < 0$, so one of $\gamma_1, \gamma_2, \gamma_3$ or γ_4 shortens $[W']$. We have found an automorphism shortening $[W']$ as described above, so we are done with this subcase.

Subcase 4b $a \notin B$ and $b \notin A$. We claim that $(A \cap B', a)$ is well defined. By Lemma 2.5, $(A \cap B', a)$ is well defined if for every $x \in A \cap B'$ with $a \not\geq x$, $(A \cap B', a)$ acts on component of $v(x)$ in $\Gamma - \text{st}(v(a))$ by conjugation. So suppose $x \in A \cap B'$ with $a \not\geq x$ and let Y_1 denote the component of $v(x)$ in $\Gamma - \text{st}(v(a))$. Then $(A \cap B', a)$ is well defined if for every $y \in Y_1$, we have $y, y^{-1} \in A$ and $y, y^{-1} \in B'$. This first condition is true since (A, a) acts on Y_1 by conjugation (since $a \not\geq x$). So suppose for contradiction that $y \in B \cup \text{lk}_L(b)$ (meaning $y \notin B'$) with $v(y) \in Y_1$. Then $y \in A$ and $a \not\geq y$. By Sublemma 3.20, we know $A \cap \text{lk}_L(b) = \emptyset$. This forces y to be in B . Let Y_2 be the component of $v(y)$ in $\Gamma - \text{st}(v(b))$. Since $y \in B$, either $b \geq y$ or (B, b) conjugates Y_2 . If $b \geq y$, the fact that $\text{st}(v(a))$ separates $v(b)$ from $v(y)$ means that $a \geq y$, a contradiction. Then (B, b) conjugates Y_2 . Since $a \notin B$, this means $a \notin Y_2$, which implies $\text{st}(v(a)) \cap Y_2 = \emptyset$. So since $\text{st}(v(a))$ separates $v(b)$ from $v(y)$ in Γ , this means that Y_2 is also a component of $\Gamma - \text{st}(v(a))$. In that case, however, $Y_1 = Y_2$, which implies $x \in B$, a contradiction. So $(A \cap B', a)$ is well defined. Note that $(B \cap A', b)$ is well defined by the same argument.

Next we will show that either $(A \cap B', a)$ or $(B \cap A', b)$ shortens $\alpha^{-1} \cdot [W]$. By Equation (3-6), we know that $0 > D(\alpha) + D(\beta)$. By Lemma 3.17, we know that

$$D(\alpha) = \langle A, A' \rangle_a - \langle a, L \rangle_a = \langle A \cap B', A' \rangle_a + \langle A \cap B, A' \rangle_a - \langle a, L \rangle_a$$

and that:

$$D(\beta) = \langle B, B' \rangle_b - \langle b, L \rangle_b = \langle B \cap A', B' \rangle_b + \langle B \cap A, B' \rangle_b - \langle b, L \rangle_b$$

Also from Lemma 3.17, we know that

$$\begin{aligned} D((A \cap B', a)) &= \langle A \cap B', A' \cup B \rangle_a - \langle a, L \rangle_a \\ &= \langle A \cap B', A' \rangle_a + \langle B \cap A, B' \cap A \rangle_a - \langle a, L \rangle_a \end{aligned}$$

and that:

$$\begin{aligned} D((B \cap A', b)) &= \langle B \cap A', B' \cup A \rangle_b - \langle b, L \rangle_b \\ &= \langle B \cap A', B' \rangle_b + \langle A \cap B, A' \cap B \rangle_b - \langle b, L \rangle_b \end{aligned}$$

We claim that $\langle A \cap B, A' \rangle_a \geq \langle A \cap B, A' \cap B \rangle_b$. Since $b \notin A$, Sublemma 3.20 says that $\text{lk}_L(b) \cap A = \emptyset$. If $(cud^{-1})^{\pm 1}$ is a subsegment of W' with $c \in A \cap B$, $d \in A' \cap B$, and u a word in $\text{lk}_L(b)$, then either u is a word in $\text{lk}_L(b) \cap \text{lk}_L(a)$, or $u = u'u_1u''$ where u' a word in $\text{lk}_L(b) \cap \text{lk}_L(a)$ and $u_1 \in \text{lk}_L(b) - \text{lk}_L(a)$. If the former is true,

then $cu d^{-1}$ is counted by $\langle A \cap B, A' \rangle_a$; if the latter is holds, then instead $cu' u_1$ is counted by $\langle A \cap B, A' \rangle_a$ (since $\text{lk}_L(b) \cap A = \emptyset$). Either way, each subsegment of W' counted by one counter is also counted by the other, showing the inequality. Similarly, we know $\langle B \cap A, B' \rangle_b \geq \langle B \cap A, B' \cap A \rangle_a$.

Putting this all together, we have that

$$0 > D(\alpha) + D(\beta) > D((A \cap B', a)) + D((B \cap A', b)).$$

So one of $(A \cap B', a)$ and $(B \cap A', b)$ shortens $[W']$. \square

Remark 3.22 The *pure automorphism group* $\text{Aut}^0 A_\Gamma$ of A_Γ is the subgroup of $\text{Aut } A_\Gamma$ generated by dominated transvections, partial conjugations and inversions. It contains those graphic automorphisms which can be expressed as products of transvections and inversions; depending on Γ , $\text{Aut}^0 A_\Gamma$ may or may not be all of $\text{Aut } A_\Gamma$. In any case, $\text{Aut}^0 A_\Gamma$ is a finite-index normal subgroup of $\text{Aut } A_\Gamma$.

Define the *pure long-range Whitehead automorphisms* Ω_ℓ^0 to be $\Omega_\ell \cap \text{Aut}^0 A_\Gamma$. If $\alpha \in \langle \Omega_\ell^0 \rangle$, then in fact, we can peak-reduce α with respect to any k -tuple of conjugacy classes W by elements of Ω_ℓ^0 . To see this, consider the proof of Lemma 3.18: when we lower peaks in factorizations of α , we move around type (1) Whitehead automorphisms in case 1, and we introduce a type (1) Whitehead automorphism in subcase 3c that is in Ω_ℓ^0 , and in no other case do we introduce a type (1) Whitehead automorphism. So if we start with a factorization of α by elements of Ω_ℓ^0 and peak-reduce it, we will end up with a peak-reduced factorization of α by elements of Ω_ℓ^0 . This technical detail is important for the application in Day [4].

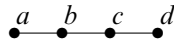
4 Attempting to extend peak reduction to $\text{Aut } A_\Gamma$

4.1 A failure of peak reduction

In this section we prove Proposition C.

Example 4.1 (Outer automorphisms of the four-vertex path) Let Γ be the four-vertex path, with labels as in Figure 1. Let P denote the subgroup of $\text{Out } A_\Gamma$ generated by the images of the inversions and the unique nontrivial graphic automorphism (which swaps a with d and b with c). Then $P \cong (\mathbb{Z}/2\mathbb{Z}) \times (\mathbb{Z}/2\mathbb{Z})^4$.

Let N denote the subgroup of $\text{Out } A_\Gamma$ generated by the images of the dominated transvections. We have adjacent dominations $b \geq a$ and $c \geq d$, and nonadjacent dominations $c \geq a$ and $b \geq d$. These are the only examples of domination in Γ . This

Figure 1: A graph Γ such that peak reduction fails on A_Γ

gives us six infinite cyclic subgroups of $\text{Aut } A_\Gamma$ generated by dominated transvections: two for each example of nonadjacent domination (multiplying on the right and on the left) and one for each example of adjacent domination. Since $e(b, c)$, these transvections commute and generate a copy of $\mathbb{Z}^6 < \text{Aut } A_\Gamma$. Each of our pairs of nonadjacent transvections differ by an inner automorphism, so $N \cong \mathbb{Z}^4$. From Equation (R6), we know that P normalizes N . No vertex of Γ has a star that separates Γ , so each partial conjugation is a full conjugation. Then by Laurence's Theorem (Theorem 2.3), we have

$$\text{Out } A_\Gamma \cong P \rtimes N.$$

Let $\phi: \mathbb{Z}^4 \rightarrow N$ be given by $\phi(p, q, r, s)(a) = ab^p c^q$ and $\phi(p, q, r, s)(d) = b^r c^s d$ (note $\phi(p, q, r, s)$ fixes the conjugacy classes b and c). Let $k \geq 2$ and let w be the conjugacy class of the cyclic word ad^k . For any $(p, q, r, s) \in \mathbb{Z}^4$, we have

$$\phi(p, q, r, s)(w) = ab^{p+r} c^{q+ks} d(b^r d)^{k-1}.$$

Note that the word on the right side is a graphically reduced cyclic word. Then $\phi(p, q, r, s)$ fixes w if and only if $p = 0$, $r = 0$ and $q = -ks$. So the stabilizer N_w is $\langle \phi(0, -k, 0, 1) \rangle$. Further, the only classes in $(\text{Out } A_\Gamma) \cdot w$ with length less than or equal to $|w|$ are the 8 classes in $P \cdot w$. Also note that if $w' = \sigma(w)$ with $\sigma \in P$, then the stabilizer $N_{w'}$ is $\langle \sigma \phi(0, -k, 0, 1) \sigma^{-1} \rangle$.

Proof of Proposition C Let Γ be as in Example 4.1. For $\alpha \in \text{Aut } A_\Gamma$, let $|\alpha|$ denote the length of the class of α in $\text{Out } A_\Gamma$ with respect to Laurence's generators. Pick a natural number k such that

$$k > 1 + \sup_{\alpha \in S} |\alpha|.$$

Let w be the conjugacy class of the cyclic word ad^k . Let $\beta \in \text{Aut } A_\Gamma$ represent the class of $\phi(0, -k, 0, 1) \in \text{Out } A_\Gamma$, with ϕ as in Example 4.1. Note that β fixes w and β does not represent an element of P .

Suppose that β can be peak-reduced with respect to w by elements of S . Since w is a minimal-length element of $(\text{Out } A_\Gamma) \cdot w$, this means we can factor β as $\gamma_m \cdots \gamma_1$ for some $\gamma_1, \dots, \gamma_m \in S$, such that for each j , $1 \leq j \leq m$, we have $|\gamma_j \cdots \gamma_1(w)| = |w|$. Each $\gamma_j \cdots \gamma_1(w)$ is the same length as w and in the same orbit, so by Example 4.1, it is in $P \cdot w$.

Fix a j , $1 \leq j \leq m$. Let $w' = \gamma_{j-1} \cdots \gamma_1(w)$. There is σ representing an element of P such that $\sigma\gamma_j$ represents an element of N . Since $|\sigma\gamma_j(w')| = |w'|$, we know from Example 4.1 that $\sigma\gamma_j(w') = w'$ (since $w' \in P \cdot w$). Also by Example 4.1, $\sigma\gamma_j$ represents $\rho\phi(0, -ks, 0, s)\rho^{-1}$ for some $s \in \mathbb{Z}$ and some $\rho \in P$. Then $|\sigma\gamma_j| = s(k+1)$; but since $\gamma_j \in S$ and therefore $|\sigma\gamma_j| \leq k$, this s must be 0. Therefore each γ_j represents an element of P . Then β represents an element of P , a contradiction. \square

4.2 Automorphisms fixing a set of basis elements

In this section we prove Proposition D.

Lemma 4.2 *Suppose x is the conjugacy class of an element of X and $\alpha \in \langle \Omega_s \rangle$. Then $\alpha(x)$ cannot be shortened by a member of Ω_ℓ .*

Proof If an element of Ω_ℓ shortens $\alpha(x)$, then it must be a type (2) automorphism (A, a) . Further, we must have $a \in (\text{supp } \alpha(x))^{\pm 1}$ and $(A-a) \cap (\text{supp } \alpha(x))^{\pm 1} \neq \emptyset$ by Lemma 3.16. We know $\text{supp } \alpha(x)$ is a clique in Γ by Lemma 3.8, so this contradicts the fact that (A, a) is long-range. \square

The image of X in H_Γ is a basis for H_Γ . By declaring this basis to be orthonormal, we get an inner product $\langle -, - \rangle$ on H_Γ .

Lemma 4.3 *If $\alpha \in \langle \Omega_s \rangle$, then for any $a, b \in X$, we have $\langle \alpha_* b, a \rangle \neq 0$ implies $a = b$, or $a \geq b$ with $e(a, b)$.*

Proof We induct on the length of α with respect to the subset of transvections $\tau_{a,b} \in \Omega_s$. The assertion is obvious if the length is zero. Assume it is true for β and that $\alpha = \beta\gamma$ where $\gamma = \tau_{c,d} \in \Omega_s$. Suppose $\gamma = \tau_{c,d}$ for some $c, d \in L$ with $c \geq d$. Then $\langle \alpha_* b, a \rangle \neq 0$ implies either that $\langle \beta_* b, a \rangle \neq 0$ or that $v(a) = v(c)$ and $\langle \beta_* b, d \rangle \neq 0$. In the first case, the lemma follows. In the second case, we have $a \geq d$ and also $d \geq b$ with $e(d, b)$ by inductive assumption. \square

Lemma 4.4 *Let $S \subset X$. The pointwise stabilizer of S in $\langle \Omega_s \rangle$ is generated by transvections $\tau_{a,b} \in \Omega_s$ with $v(b) \notin S$.*

Proof Suppose $\alpha \in \langle \Omega_s \rangle$ and α fixes S pointwise. Let $A = \alpha_* \in \text{Aut } H_\Gamma$. Since α fixes S pointwise, for any $x \in S$, $\langle Ax, y \rangle = 0$ for all $y \in X - x$ and $\langle Ax, x \rangle = 1$. Let $X = C_1 \cup \cdots \cup C_m$ be the decomposition of X into adjacent domination equivalence classes. First of all, for each i , we can row-reduce A such that for any $a, b \in C_i$, $\langle Ab, a \rangle$ is 0 if $a \neq b$ and 1 if $a = b$. In fact, we can do this by multiplying α by

transvections $\tau_{a,b}$ for various $a, b \in C_i$ with $v(b) \notin S$ (each $\tau_{a,b}$ corresponds to a row operation). Now suppose some $\langle Ab, a \rangle \neq 0$ with $a \not\sim b$. Then $a \geq b$ and $e(a, b)$ by Lemma 4.3, and $v(b) \notin S$. Since we have already reduced the diagonal, applying some power of $\tau_{a,b}$ to α will change this entry to zero. Of course, by doing this in appropriate order to the nonzero entries with $a \not\sim b$, we can row-reduce the rest of A . So we can reduce α to the identity by applying elements $\tau_{a,b} \in \Omega_S$ with $v(b) \notin S$. \square

Proof of Proposition D Suppose $W = (w_1, \dots, w_k)$ is a k -tuple of conjugacy classes with each $|w_i| = 1$ and suppose $\alpha \in \text{Aut } A_\Gamma$ with $|\alpha \cdot W| = |W|$. By Theorem B, we write $\alpha = \beta\gamma$ where $\beta \in \langle \Omega_\ell \rangle$ and $\gamma \in \langle \Omega_S \rangle$. Also by Theorem B, we have a factorization $\beta = \delta_m \cdots \delta_1$ by elements of Ω_ℓ that is peak-reduced with respect to $\gamma \cdot W$. By Lemma 4.2, this δ_1 cannot shorten $\gamma \cdot W$. So since $\delta_m \cdots \delta_1$ is peak-reduced, we have $|\gamma \cdot W| = |\alpha \cdot W| = |W|$.

Since each w_i is a minimal-length representative of its $\text{Aut } A_\Gamma$ -orbit, it follows that each $\gamma(w_i)$ is the conjugacy class of an element of L . If $x, y \in L$ with y conjugate to $\gamma(x)$, then $x \sim y$ by Corollary 3.11 and $e(x, y)$ by Lemma 3.8. In general, if S is a basis for \mathbb{Z}^j for some j , $S' \subset S$ and $A \in \text{SL}(j, \mathbb{Z})$ sends S' to a subset of $S^{\pm 1}$, then there is $B \in \text{SL}(j, \mathbb{Z})$ such that $B|_{S'} = A|_{S'}$ and B restricts to a permutation on $S \cup (-S)$ (this can be proven by a row reduction argument). Then from Corollary 3.11, we deduce that there is a type (1) Whitehead automorphism σ such that $\sigma\gamma \cdot W = W$ and $\sigma \in \langle \Omega_S \rangle$.

Then by Lemma 4.4, we can write $\sigma\gamma$ as a product $\phi_r \cdots \phi_1$ of elements ϕ_1, \dots, ϕ_r of Ω_S that fix $\text{supp } W$ pointwise. Then the following is a peak-reduced factorization of α by elements of $\Omega_\ell \cup \Omega_S$:

$$\alpha = \delta_m \cdots \delta_1 \sigma^{-1} \phi_r \cdots \phi_1. \quad \square$$

Corollary 4.5 Suppose $W = (w_1, \dots, w_k)$ is a k -tuple of conjugacy classes of A_Γ with each $|w_i| = 1$. Then the stabilizer $(\text{Aut } A_\Gamma)_W$ of W in $\text{Aut } A_\Gamma$ is generated by $(\Omega_\ell \cup \Omega_S) \cap ((\text{Aut } A_\Gamma)_W)$.

Proof Let $\bar{\Delta}$ be the directed multigraph whose vertices are k -tuples of conjugacy classes W' with $|W'| = |W|$, and an edge from W_1 to W_2 , labeled by α , if $\alpha \in \Omega_\ell \cup \Omega_S$ with $\alpha(W_1) = W_2$. Let Δ be the (undirected) connected component of W in $\bar{\Delta}$. This is called the *Whitehead graph* of W . We map the paths of Δ to $\text{Aut } A_\Gamma$ by composing their edge labels; a path from a vertex W_1 to a vertex W_2 will map to an automorphism α with $\alpha(W_1) = W_2$ (this is true for paths of length 1 and remains true under concatenations). In particular, $\pi_1(\Delta, W) \rightarrow (\text{Aut } A_\Gamma)_W$. By Proposition D, if $\alpha \in (\text{Aut } A_\Gamma)_W$, we can write $\alpha = \beta_m \cdots \beta_1$ where each $\beta_i \in \Omega_\ell \cup \Omega_S$ and for each i ,

$0 \leq i \leq m$, we have $|\beta_i \cdots \beta_1 \cdot W| = |W|$. Then $\beta_m \cdots \beta_1$ describes a path in Δ mapping to α , and the map $\pi_1(\Delta, W) \rightarrow (\text{Aut } A_\Gamma)_W$ is surjective.

To get generators for $\pi_1(\Delta, W)$, we pick a maximal tree for Δ . Since each $|w_i| = 1$, we know each vertex of Δ is the image of W under some permutation of L . Then we can pick our maximal tree T to be a union of edges labeled by type (1) Whitehead automorphisms originating at W . There is a unique loop in $\pi_1(\Delta, W)$ for each (directed) edge in $\Delta - T$ (the loop leaving T only to cross this edge once); these loops generate $\pi_1(\Delta, W)$, and the images of these loops in $(\text{Aut } A_\Gamma)_W$ generate.

If α is a type (2) Whitehead automorphism labeling an edge in $\Delta - T$, then α labels a loop from a vertex W' to itself (if a type (2) Whitehead automorphism changes a vertex W' of Δ , then it lengthens it). So if $W \neq W'$ there is a type (1) Whitehead automorphism σ labeling the edge in T from W to W' , and by Equation (R6), the automorphism $\sigma\alpha\sigma^{-1}$ is a Whitehead automorphism labeling an edge from W to itself. If α is a type (1) Whitehead automorphism labeling an edge in $\Delta - T$ that is not a loop at W , then by relations of type (R7), the loop based at W through edges in T and α is redundant with a type (1) Whitehead automorphism labeling an edge from W to itself. So in fact, the loops in $\pi_1(\Delta, W)$ of length 1 map to a generating set for $(\text{Aut } A_\Gamma)_W$. By definition, they map to $(\Omega_\ell \cup \Omega_s) \cap ((\text{Aut } A_\Gamma)_W)$. \square

Remark 4.6 There is another case where a peak reduction theorem holds for $\text{Aut } A_\Gamma$: the author has shown in [4] that if $w = [a_1, b_1] \cdots [a_k, b_k]$ for distinct $a_1, \dots, a_k, b_1, \dots, b_k \in X$ and $\alpha \in \text{Aut } A_\Gamma$ with $\alpha(w) = w$, then α can be peak-reduced with respect to w by elements of Ω .

5 A presentation for $\text{Aut } A_\Gamma$

The goal of this section is to prove Theorem A. Recall that Φ is the free group on Ω . Let $\Phi_\ell < \Phi$ be the subgroup generated by Ω_ℓ . Let $R_\ell = R \cap \Phi_\ell$. Denote the normal closure of $\langle R_\ell \rangle$ in Φ_ℓ by $\overline{\langle R_\ell \rangle}$. Say that w_1 and w_2 in Φ_ℓ are congruent modulo R_ℓ if $w_1 w_2^{-1} \in \overline{\langle R_\ell \rangle}$. Similarly, we denote the normal closure of $\langle R \rangle$ in Φ by $\overline{\langle R \rangle}$ and say that w_1 and w_2 in Φ are congruent modulo R if $w_1 w_2^{-1} \in \overline{\langle R \rangle}$.

Lemma 5.1 *Suppose $\alpha, \beta \in \Omega_\ell$ and $[W]$ is a k -tuple of conjugacy classes of A_Γ . Suppose $\beta\alpha^{-1}$ forms a peak with respect to $[W]$. Then there exist $\delta_1, \dots, \delta_k \in \Omega_\ell$ such that, when multiplied in Φ_ℓ , $\beta\alpha^{-1}$ is congruent to $\delta_k \cdots \delta_1$ modulo R_ℓ and for each $i, 1 \leq i < k$, we have*

$$|(\delta_i \cdots \delta_1) \cdot [W]| < |\alpha^{-1} \cdot [W]|.$$

Proof This lemma is a refinement of Lemma 3.18, so to prove it, it is enough to review the proof of Lemma 3.18, noting in each case that the peak-lowering factorization $\delta_k \cdots \delta_1$ is congruent to $\beta\alpha^{-1}$ modulo R_ℓ . This will be true if in each case, the only manipulations we apply to elements of $\text{Aut } A_\Gamma$ are applications of relations in R . At the start of the proof, we established that if $\alpha = (A, a)$ and $\beta = (B, b)$, we may switch α and β or swap β with $(L - B - \text{lk}_L(b), b^{-1})$. By the symmetry in the statement, it is again apparent that we may still switch α and β if necessary. In showing we could swap β with $(L - B - \text{lk}_L(b), b^{-1})$, we used Relations (R8)–(R10). In case 1, we used Relation (R6). We used Relation (R1) in cases 2, 3 and 4. In case 2, we used Relation (R3)(b). In case 3, we used Relation (R2) in subcase 3a; Relations (R2), (R3)(a) and (R4)(a) in subcase 3b; and Relations (R2) and (R5) in subcase 3c. In case 4, we used Relation (R2) and invoked case 3. These were the only manipulations done to elements of $\text{Aut } A_\Gamma$ in that proof, so we are done. \square

The following lemma is similar to Proposition 6.2.5 of Culler–Vogtmann.

Lemma 5.2 *Let V be a k -tuple of conjugacy classes whose elements are all the conjugacy classes in A_Γ of length 2, each appearing once. If $(A, a) \in \Omega_\ell$ and $|(A, a) \cdot V| \leq |V|$, then (A, a) is trivial or is the conjugation $(L - a^{-1}, a)$.*

Proof We partition L into the following seven sets:

$$L = (A \cap A^{-1}) + (A - A^{-1} - a) + (A^{-1} - A - a^{-1}) \\ + (L - \text{lk}_L(a) - A \cup A^{-1}) + \text{lk}_L(a) + \{a\} + \{a^{-1}\}$$

If bc is a cyclic word of length 2 (not necessarily with $b \neq c$), then we can use Lemma 3.16 to compute $D_{[bc]}((A, a))$ according to the sets in the partition of L that b and c are members of. Note that since bc is a cyclic word, we may switch b with c in our enumeration of cases. Also, if both $b, c \in (L - \text{lk}_L(a) - A \cup A^{-1}) + \text{lk}_L(a) + \{a\} + \{a^{-1}\}$, then $D_{[bc]}((A, a)) = 0$. We list the remaining cases in Table 1.

As usual, $n = |X|$. Let $m = |X - \text{lk}(v(a))|$, let $x = \frac{1}{2}|A \cap A^{-1}|$, and let $y = |A - A^{-1} - a| = |A^{-1} - A - a^{-1}|$. Then $|L - \text{lk}_L(a) - A \cup A^{-1}| = 2(m - x - y)$. We list the number of conjugacy classes appearing in V of the form $[bc]$ as b and c are in the different subsets of L in Table 2, leaving out the cases in which $D_{[bc]}((A, a)) = 0$.

We compute $D_V((A, a))$ from the two tables by taking products and summing:

$$D_V((A, a)) = 4xy + 8x(m - x - y) + 2y(y + 1) + 4y(m - x - y) + 4y(n - m)$$

Note that the contribution to $D_V((A, a))$ from the entries in V containing a copy of a or a^{-1} cancel each other out. Since the numbers x , y , $(m - x - y)$ and $(n - m)$ are

	c	$A \cap A^{-1}$	$A - A^{-1} - a$	$A^{-1} - A - a^{-1}$
b				
$A \cap A^{-1}$		0		
$A - A^{-1} - a$		1	2	
$A^{-1} - A - a^{-1}$		1	0	2
$L - \text{lk}_L(a) - A \cup A^{-1}$		2	1	1
$\text{lk}_L(a)$		0	1	1
$\{a\}$		0	1	-1
$\{a^{-1}\}$		0	-1	1

Table 1: The value of $D_{[bc]}((A, a))$ as b and c are in different subsets of L

	c	$A \cap A^{-1}$	$A - A^{-1} - a$	$A^{-1} - A - a^{-1}$
b				
$A \cap A^{-1}$		-		
$A - A^{-1} - a$		$2xy$	$\frac{y(y+1)}{2}$	
$A^{-1} - A - a^{-1}$		$2xy$	-	$\frac{y(y+1)}{2}$
$L - \text{lk}_L(a) - A \cup A^{-1}$		$4x(m-x-y)$	$2y(m-x-y)$	$2y(m-x-y)$
$\text{lk}_L(a)$		-	$y(n-m)$	$y(n-m)$
$\{a\}$		-	y	y
$\{a^{-1}\}$		-	y	y

Table 2: The number of conjugacy classes in V of the form $[bc]$, as b and c are in different subsets of L

all nonnegative (they count the cardinalities of sets), we know that $D_V((A, a))$ cannot be negative; further, for $D_V((A, a))$ to be zero, we must have each of the terms equal to zero. This implies that $y = 0$ and that $x(m - x) = 0$, which means that (A, a) is either the trivial automorphism $(\{a\}, a)$ or the conjugation $(L - \text{lk}_L(a) - a^{-1}, a)$. \square

Lemma 5.3 *The group of inner automorphisms $\text{Inn } A_\Gamma$ is a right-angled Artin group. Specifically, if Z is the intersection of the stars in Γ of the elements of X , and Γ' is the full subgraph of Γ on the vertices $X - Z$, then the map sending $x \in X - Z$ to conjugation by x in A_Γ is an isomorphism $A_{\Gamma'} \xrightarrow{\cong} \text{Inn } A_\Gamma$.*

Proof By the Servatius centralizer theorem (Theorem 2.1), we know that the center $Z(A_\Gamma)$ is $\langle Z \rangle < A_\Gamma$. The obvious inclusion $A_{\Gamma'} \hookrightarrow A_\Gamma$ induces an isomorphism $A_{\Gamma'} \xrightarrow{\cong} A_\Gamma / Z(A_\Gamma)$. By composing this map with the usual isomorphism $A_\Gamma / Z(A_\Gamma) \xrightarrow{\cong} \text{Inn } A_\Gamma$, we obtain the isomorphism in the statement. \square

The proof of the following proposition is based on McCool's argument from [8].

Proposition 5.4 *The group $\langle \Omega_\ell \rangle < \text{Aut } A_\Gamma$ has the presentation $\langle \Omega_\ell \mid R_\ell \rangle$.*

Proof We already know that every relation in R_ℓ is an identity of $\langle \Omega_\ell \rangle$, so it is enough to show that every word representing the trivial element in $\langle \Omega_\ell \rangle$ is a product of conjugates of elements of R_ℓ . Suppose $w \in \Phi_\ell$ represents the trivial element in $\langle \Omega_\ell \rangle < \text{Aut } A_\Gamma$. We claim that there is $w' \in \overline{\langle R_\ell \rangle}$ such that ww' is a product of type (1) Whitehead automorphisms and conjugations. Let V_0 be a k -tuple containing each conjugacy class of A_Γ of length 2 once.

We will prove the claim by induction on the peaks of w with respect to V_0 ; specifically, inducting on the number of points between peaks of maximal height and also on the maximum height of peaks. Write $w = \alpha_j \cdots \alpha_1$ for $\alpha_j, \dots, \alpha_1 \in \Omega_\ell$. In our base case, we assume that $\alpha_j \cdots \alpha_1$ is a factorization of w that is peak-reduced with respect to V_0 . By Lemma 5.2, we know that V_0 is a minimal-length representative of its $\text{Aut } A_\Gamma$ orbit. So since our factorization of w is peak-reduced, for each i , we have $|(\alpha_i \cdots \alpha_1) \cdot V_0| = |V_0|$. We claim that for each i , $(\alpha_i \cdots \alpha_1) \cdot V_0$ is a k -tuple containing each conjugacy class of length 2 once. This is true if $i = 0$ by assumption. Now assume it for $i - 1$; since $|(\alpha_i \cdots \alpha_1) \cdot V_0| = |(\alpha_{i-1} \cdots \alpha_1) \cdot V_0|$, we know by Lemma 5.2 that α_i is then either trivial, a conjugation or a type (1) Whitehead automorphism and the statement is then true for i . So in our base case, w is already a product of type (1) Whitehead automorphisms and conjugations.

For the inductive step, suppose that $\alpha_k \cdots \alpha_1$ has peaks with respect to V_0 . Let α_i be a peak of maximal height. Then by Lemma 5.1, there are $\delta_1, \dots, \delta_m \in \Omega_\ell$ such that $(\alpha_{i+1}\alpha_i)^{-1}\delta_m \cdots \delta_1 \in \overline{\langle R_\ell \rangle}$ and such that we can lower the peak at α_i in $\alpha_k \cdots \alpha_1$ by substituting in $\delta_m \cdots \delta_1$ for $\alpha_{i+1}\alpha_i$. So we define

$$w_1 = (\alpha_{i-1} \cdots \alpha_1)^{-1} (\alpha_{i+1}\alpha_i)^{-1} \delta_m \cdots \delta_1 (\alpha_{i-1} \cdots \alpha_1) \in \overline{\langle R_\ell \rangle}.$$

Then $ww_1 = \alpha_k \cdots \alpha_{i+2}\delta_m \cdots \delta_1\alpha_{i-1} \cdots \alpha_1$ has a smaller number of points between maximal-height peaks that w with respect to V_0 , or its maximal-height peak is shorter. So we have reduced the peaks of ww_1 , and we invoke the inductive hypothesis for ww_1 : we have a $w_2 \in \overline{\langle R_\ell \rangle}$ such that ww_1w_2 is a product of type (1) Whitehead automorphisms and conjugations. So $w_1w_2 \in \overline{\langle R_\ell \rangle}$ satisfies the conclusions of our inductive claim.

So we have that w is congruent modulo R_ℓ to a product of type (1) Whitehead automorphisms and conjugations. Then by applying instances of Relation (R6), we know that w is congruent to a product $\beta\gamma$ where β is a product of type (1) Whitehead automorphisms and γ is a product of conjugation automorphisms in Ω_ℓ . The subgroup

of $\text{Aut } A_\Gamma$ generated by type (1) Whitehead automorphisms acts faithfully on $\text{Aut } H_\Gamma$, so since α maps to the trivial element of $\text{Aut } A_\Gamma$ and γ is in the kernel of the homology representation, we deduce that β represents the trivial automorphism. So by some instances of Relation (R7), we know that w is congruent modulo R_ℓ to γ , which represents the trivial automorphism in $\text{Inn } A_\Gamma$.

Let Z and Γ' be as in Lemma 5.3. Map the free group on $X - Z$ to Φ_ℓ by sending $a \in X - Z$ to $(L - \text{lk}_L(a) - a^{-1}, a)$. This sends the relations from the right-angled Artin group presentation of $A_{\Gamma'}$ to instances of Relation (R3)(b). Of course, this map descends to the isomorphism $A_{\Gamma'} \xrightarrow{\cong} \text{Inn } A_\Gamma$ in Lemma 5.3. Then since γ represents the trivial element of $\text{Inn } A_\Gamma$, it corresponds to an element w_γ of the free group on $X - Z$ that maps to the trivial element of $A_{\Gamma'}$. This w_γ is a product of conjugates of the relations from the presentation of A_Γ , so γ is a product of conjugates of instances of Relation (R3)(b). So γ is in $\overline{\langle R_\ell \rangle}$, and therefore w is in $\overline{\langle R_\ell \rangle}$. \square

If $a, b \in X$ with $a \in \text{lk}_L(b)$ and $a \sim b$, then the type (1) Whitehead automorphism $\sigma_{a,b}$ of Relation (R5) exists. According to that relation, we have $\sigma_{a,b} \in \langle \Omega_s \rangle$. Let $P_s \subset \Omega$ be the finite subgroup of $\text{Aut } A_\Gamma$ generated by such $\sigma_{a,b}$ as a and b range over all adjacent domination-equivalent pairs in X . Let Φ_s be the subgroup of Φ generated by $\Omega_s \cup P_s$. Let $R_s = R \cap \Phi_s$.

Proposition 5.5 *The group $\langle \Omega_s \rangle$ has the presentation $\langle \Omega_s \cup P_s \mid R_s \rangle$.*

Proof Let $G = \langle \Omega_s \rangle < \text{Aut } A_\Gamma$ and let $\tilde{G} = \langle \Omega_s \cup P_s \mid R_s \rangle$. We know that each of the relations in R_s is an identity in \tilde{G} , so we have homomorphism $G \rightarrow \tilde{G}$ by sending each element of Ω_s to its own coset. We will show this map is an isomorphism by constructing an inverse.

By Corollary 3.11, the group G has a presentation where the generators are

$$\{E_{a,b} \mid a, b \in X, a \in \text{lk}_L(b) \text{ and } a \geq b\}$$

and the relations are all the relations of the forms (1)–(4) from Proposition 3.10. This presentation identifies each $E_{a,b}$ with the corresponding $\tau_{a,b}$.

By Relations (R1), (R2), (R5) and (R7), we know \tilde{G} is generated by the transvections $(\{a, b\}, a)$ with $a \in \text{lk}_L(b)$ and $a \geq b$. We map \tilde{G} to G by sending each $(\{a, b\}, a)$ to the corresponding $\tau_{a,b}$. We will show that this is a homomorphism by checking that the relations of our presentation for G from Proposition 3.10 already hold in \tilde{G} . Relation (1) follows from Relations (R2) and (R3)(b). Relation (2) follows from Relation (R4)(b). For any $a, b \in X$ with $a \in \text{lk}_L(b)$ and $a \sim b$, we know from Relation (R5) that $\tau_{a,b} \tau_{b,a}^{-1} \tau_{a,b}$ is $\sigma_{a,b}$, which has order 4 by Relation (R7) (here we are

using that $\tau_{a^{-1},b^{-1}} = \tau_{a,b}$, which holds because $a \in \text{lk}_L(b)$, and that $\tau_{a^{-1},b} = \tau_{a,b}^{-1}$. This means that Relation (3) already holds in \tilde{G} . By Relation (R5), $(\tau_{a,b}\tau_{b,a}^{-1}\tau_{a,b}\tau_{b,a})^3$ is $(\sigma_{a,b}\tau_{b,a})^3$, which is $\tau_{a,b^{-1}}\tau_{b^{-1},a^{-1}}\tau_{a^{-1},b}\sigma_{a,b}^3$ by Relation (R6), which is $\sigma_{a,b}^{-1}\sigma_{a,b}^3$ by Relation (R5) and using the facts that

$$\tau_{a,b^{-1}} = \tau_{a,b}^{-1}, \quad \tau_{b^{-1},a^{-1}} = \tau_{a^{-1},b}^{-1} \quad \text{and} \quad \tau_{a^{-1},b} = \tau_{a^{-1},b^{-1}}^{-1}.$$

Then Relation (4) already holds in \tilde{G} .

So we map \tilde{G} to G homomorphically by sending $(\{a,b\}, a) \in \Omega_s$ to $E_{a,b}$. It is apparent (from looking at the action on generating sets) that this homomorphism is the inverse to the homomorphism G to \tilde{G} that sends each element of $\Omega_s \cup P_s$ to its own coset. So $\langle \Omega_s \rangle < \text{Aut } A_\Gamma$ has the presentation $\langle \Omega_s \cup P_s \mid R_s \rangle$. \square

Proposition 5.6 *Every $w \in \Phi$ is congruent modulo R to a product uv for some $u \in \Phi_\ell$ and $v \in \Phi_s$.*

Proof This proposition is a refinement of part (1) of Theorem B. The only manipulations of elements of $\text{Aut } A_\Gamma$ done in that proof are through the sorting substitutions in Definition 3.2. Each of the sorting substitutions comes from applications of relations from R , as in Lemma 3.4. So the entire argument goes through for $\langle \Omega \mid R \rangle$. \square

Proof of Theorem A We have already shown that all the relations in R are identities of $\text{Aut } A_\Gamma$ (Proposition 2.10), so it is enough to show that any element of Φ representing the trivial element of $\text{Aut } A_\Gamma$ is in $\overline{\langle R \rangle}$. Let $w \in \Phi$ represent the trivial element of $\text{Aut } A_\Gamma$. By Proposition 5.6, w is congruent modulo R to a product uv for $u \in \Phi_\ell$ and $v \in \Phi_x$. Let $[u] \in \langle \Omega_\ell \rangle$ and $[v] \in \langle \Omega_s \rangle$ denote the elements of $\text{Aut } A_\Gamma$ they represent.

Let W_0 be the elements of X as an n -tuple of conjugacy classes. Suppose that $[v]$ is not a type (1) Whitehead automorphism; then $[v]$ sends W_0 to a strictly longer n -tuple. By part (3) of Theorem B, we peak-reduce $[u]$ with respect to $[v] \cdot W_0$. Since $[u][v]$ is trivial, $[u]$ sends $[v] \cdot W_0$ to W_0 ; since we have peak-reduced $[u]$, the first automorphism $\alpha \in \Omega_\ell$ in our peak-reduced factorization of $[u]$ shortens $[v] \cdot W_0$. However, this contradicts Lemma 4.2.

So $[v]$ is a type (1) Whitehead automorphism, which we write as σ . Then w is congruent to $u\sigma^{-1}\sigma v$ modulo R . From Proposition 5.5, we know that $u\sigma^{-1}$ is a product of conjugates of members of $R_s \subset R$, and from Proposition 5.4, we know that σv is a product of conjugates of members of $R_\ell \subset R$. So $w \in \overline{\langle R \rangle}$. \square

6 Closing remarks

The applications of peak reduction on F_n mentioned in the introduction all suggest further applications of Theorem B. Firstly, peak reduction can be used to get finite generation and finite presentation results for stabilizers of k -tuples of conjugacy classes in $\text{Aut } F_n$, as in McCool [8]. Along these lines, the author has used Theorem B in [4] to show that an analog of the mapping class group of a surface inside $\text{Aut } A_\Gamma$ is finitely generated. Generally, one could obtain further results similar to Corollary 4.5 by proving propositions similar to Proposition D, *ie* by finding additional special cases where peak reduction works on all of $\text{Aut } A_\Gamma$.

Peak reduction on the free group F_n makes an algorithm possible that determines whether two k -tuples of conjugacy classes in F_n are in the same $\text{Aut } F_n$ orbit (and makes it possible to find an automorphism taking one to the other, if it exists). Please see Lyndon–Schupp [7, Proposition I.4.19], for a description of this algorithm. As for free abelian groups, row reduction lets us transform k -tuples of elements of \mathbb{Z}^n standard representatives of their $GL(n, \mathbb{Z})$ -orbits (and more carefully, to find an automorphism taking one to another if it exists). So it seems natural to conjecture the existence of a similar algorithm for $\text{Aut } A_\Gamma$:

Conjecture 6.1 *There is an algorithm which, given $u, v \in A_\Gamma$, produces $\alpha \in \text{Aut } A_\Gamma$ with $\alpha(u) = v$ or determines in finite time that no such automorphism exists.*

Part (3) of Theorem B easily implies such an algorithm if we are only considering $\alpha \in \langle \Omega_\ell \rangle$, and part (2) of Theorem B suggests a row reduction approach if we are only considering $\alpha \in \langle \Omega_s \rangle$. However, it is not clear how these methods could be extended to apply to all of $\text{Aut } A_\Gamma$. Proposition C indicates that it will not be possible to produce the algorithm in Conjecture 6.1 by a direct generalization of the approach for free groups.

Finally, it may be possible to use these algorithmic techniques to improve our understanding of spaces that $\text{Aut } A_\Gamma$ acts on. As in Culler–Vogtmann [3], it should be possible to use peak-reduction techniques to find paths in $\text{Aut } A_\Gamma$ -spaces that behave nicely with respect to combinatorial Morse functions. In particular, this should help us to better understand outer space of right-angled Artin groups, as defined in Charney–Crisp–Vogtmann [2, Section 4] for triangle-free Γ . For general Γ , certain spaces of isometric actions of A_Γ on CAT(0) cubical complexes are $\text{Aut } A_\Gamma$ -spaces. Hopefully our techniques could lead to a better understanding of these spaces as well.

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*Department of Mathematics, California Institute of Technology
1200 E California Blvd, Pasadena, CA 91101, USA*

`mattday@caltech.edu`

<http://www.its.caltech.edu/~mattday>

Proposed: Walter Neumann
Seconded: Joan Birman, Leonid Polterovich

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