

involve

a journal of mathematics

Characterizing optimal point sets determining one distinct triangle

Hazel N. Brenner, James S. Depret-Guillaume,
Eyvindur A. Palsson and Robert W. Stuckey



Characterizing optimal point sets determining one distinct triangle

Hazel N. Brenner, James S. Depret-Guillaume,
Eyvindur A. Palsson and Robert W. Stuckey

(Communicated by Kenneth S. Berenhaut)

We determine the maximum number of points in \mathbb{R}^d which form exactly t distinct triangles, where we restrict ourselves to the case of $t = 1$. We denote this quantity by $F_d(t)$. It is known from the work of Epstein et al. (*Integers* **18** (2018), art. id. A16) that $F_2(1) = 4$. Here we show somewhat surprisingly that $F_3(1) = 4$ and $F_d(1) = d + 1$, whenever $d \geq 3$, and characterize the optimal point configurations. This is an extension of a variant of the distinct distance problem put forward by Erdős and Fishburn (*Discrete Math.* **160**:1-3 (1996), 115–125).

1. Introduction

Erdős [1946] proposed his distinct distance conjecture, which states that any set of n points in the plane will define at least $\Omega(n/\sqrt{\log n})$ distinct distances. Since that time, optimal points sets have been a heavily studied topic within the field of discrete geometry. Guth and Katz [2015] made significant progress towards proving this conjecture when they showed that a set of n points in the plane defined at least $\Omega(n/\log n)$ distinct distances. The analogous problems in dimensions 3 and higher remain open.

Erdős and Fishburn [1996] asked a question related to this: given a positive integer k , what is the maximum number of points which can be embedded in the plane such that precisely k distinct distances are defined, and can all such point configurations be characterized? In their paper, Erdős and Fishburn characterized the optimal configurations for $1 \leq k \leq 4$, and this was extended by Shinohara [2008] for $k = 5$ and by Wei [2012] for $k = 6$. Further, Erdős [1975] conjectured that for sufficiently large values of k , an optimal point configuration exists in the triangular lattice, which continues as an open conjecture. (Figure 1 shows the optimal configurations for k distinct distances in the plane when $2 \leq k \leq 6$.)

MSC2010: primary 52C10; secondary 52C35.

Keywords: one-triangle problem, Erdős problem, optimal configurations, finite point configurations. The work of Palsson was supported in part by Simons Foundation Grant #360560.

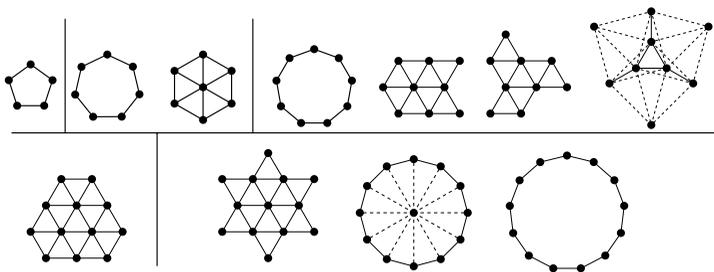


Figure 1. Optimal, or maximal, point set configurations determining exactly k distinct distances in the plane for $2 \leq k \leq 6$ [Brass et al. 2005]. For all k , $2 < k \leq 6$, there exists a configuration in the triangular lattice; Erdős conjectured that this is always true when k is sufficiently large.

Erdős' distance problem can be extended to consider triangles in place of distances. Since the set of distances generated by a point set can be thought of as being determined by the collection of two-point subsets, we may analogously consider the set of triangles formed by a point set as determined by the collection of three-point subsets. Erdős' distance problem then becomes: what is the minimum number of distinct triangles formed by a collection of n points in the plane? Hence, the analogue of Erdős and Fishburn's problem is: given a positive integer t , what is the maximum number of points, n , placed in the plane which define exactly t distinct triangles? Epstein et al. [2018] focused on the latter of these analogues and showed that $n = 4$ for $t = 1$, and $n = 5$ for $t = 2$. Finding maximal point sets in the plane remains an open question for higher values of t . As mentioned above, the higher-dimensional analogues of Erdős' distance problem are as yet open. Here we concern ourselves with the higher-dimensional analogue of Erdős and Fishburn's question, rather than with higher values of k . Our main result is the following:

Theorem 1.1. *Suppose $S \subset \mathbb{R}^d$ determines a single distinct triangle T :*

- (1) *If T is equilateral, then S is contained in the set of vertices of a regular d -simplex, and in particular $|S| \leq d + 1$.*
- (2) *If T is not equilateral, then $|S| \leq 4$.*

We can then state the following corollary in the language of optimal point configurations:

Corollary 1.2. *Let $F_d(t)$ denote the maximum number of points which can be placed in \mathbb{R}^d to determine exactly t distinct triangles. Then:*

- (1) *$F_3(1) = 4$ and the only configurations which achieve this are the vertices of a square, a rectangle, or a tetrahedron.*

- (2) $F_d(1) = d + 1$ when $d > 3$ and the only configuration which achieves this is the regular d -simplex.

We also make three observations. First, in \mathbb{R}^d , $d > 3$, the d -simplex is the unique optimal point configuration, yet in dimensions 2 and 3 this is not so. Second, in addition to the above, the 2-simplex fails to be optimal in \mathbb{R}^2 . Third, note that in \mathbb{R}^3 there are two optimal configurations, viz. the tetrahedron (3-simplex) and the rectangle (to include the square), while in \mathbb{R}^2 and \mathbb{R}^d , $d > 3$, there is a single family of solutions (by considering the square to be a special case of the rectangle), and hence, the d -simplex fails to be unique. This transition that happens in \mathbb{R}^3 is surprising and novel. In addition to the above, we have the following notable remark:

Remark 1.3. If $d = 3$, both (1) and (2) of [Theorem 1.1](#) yield optimal configurations. For (1) these configurations are specifically the vertices of the regular tetrahedron, and for (2) they are the vertices of the square, the vertices of a tetrahedron with isosceles faces and the vertices of a tetrahedron with scalene faces. These can be uniquely determined as distance graphs which can be realized in \mathbb{R}^3 in the above ways.

In [Section 5](#) we offer a proof of this remark by way of a construction of said distance graphs and point sets in \mathbb{R}^3 that satisfy them. This remark is particularly interesting as our framework for an upcoming paper in preprint arrives at constructions for optimal configurations determining few distinct triangles by considering the number of distinct distances that can be determined by such configurations [[Brenner et al. 2019](#)]. This remark shows that there can exist distinct optimal configurations determining a given number of distinct triangles that determine different numbers of distinct distances. Interestingly, this is not the case for optimal configurations determining two distinct triangles, which may only determine two distinct distances.

2. Definitions and lemmas

We formalize the concepts of a triangle and set out our notation with the following definitions:

Definition 2.1. Given a finite point set $P \subset \mathbb{R}^d$, $d \geq 3$, we say two triples (a, b, c) , $(a', b', c') \in P^3$ are equivalent if there is an isometry mapping one to the other, and we denote this as $(a, b, c) \sim (a', b', c')$.

Definition 2.2. Given a finite point set $P \subset \mathbb{R}^d$, $d \geq 3$, we denote by P_{nc}^3 the set of noncollinear triples $(a, b, c) \in P^3$.

Definition 2.3. Given a finite point set $P \subset \mathbb{R}^2$, we define the set of distinct triangles determined by P as

$$T(P) := P_{\text{nc}}^3 / \sim . \tag{2-1}$$

In this paper when we discuss and count the number of distinct triangles of a finite point set $P \in \mathbb{R}^d$ we are precisely working with the set $T(P)$. Note that this excludes degenerate triangles where all three points lie on a line.

Definition 2.4. Let p and q be points in \mathbb{R}^d for $d \geq 1$. We denote the Euclidean distance between p and q by $d(p, q)$.

Theorem 1.1(1) is a direct consequence of the following lemma:

Lemma 2.5. Let S be a set of points in \mathbb{R}^d , $d \geq 3$, which defines a single distinct equilateral triangle. Then S has at most $d + 1$ points.

3. Proof of Theorem 1.1

As stated previously, **Theorem 1.1(1)** follows directly from **Lemma 2.5**, so we will omit proof in this section in favor of proving **Lemma 2.5** in **Section 4**. Then, to prove **Theorem 1.1(2)**, we will consider the case where T is isosceles and the case where T is scalene separately. In both cases, we assume towards a contradiction that there exists a point set S containing five points determining such a triangle. Our argument will be made purely on distance graphs and will thus not depend on dimension.

Proof of Theorem 1.1(2). Assume towards a contradiction that there exists a point set S containing five points which determines one distinct nonequilateral triangle. For convenience, we then split into cases, dealing with scalene and isosceles triangles separately.

Scalene: Fix an arbitrary point \mathcal{O} in S and consider the distances from \mathcal{O} to the remaining four points. Note that clearly a point set determining exactly one distinct, scalene triangle determines only three distinct distances. So, by the pigeonhole principle, two of the distances from \mathcal{O} to the remaining points are equal. Without loss of generality, say that the repeated distance is d_1 and specifically $\mathcal{O}A = \mathcal{O}B = d_1$. Then, clearly $\triangle \mathcal{O}AB$ is an isosceles triangle, but we assumed that the only distinct triangle determined by this point set is scalene, so we have the desired contradiction.

Isosceles: Let d_1 denote the repeated edge length of T and d_2 the remaining edge length. Similarly to the above, fix an arbitrary point \mathcal{O} and consider the distances from \mathcal{O} to the remaining points. Clearly, by the same argument as the above, if any two of these distances were d_2 , an isosceles triangle with repeated edge length d_2 would be determined, which would be a contradiction. So, assume that at least three of the four distances are d_1 and label the three points determining them A , B and C . Then consider each of the triangles consisting of two of the points and \mathcal{O} , e.g., $\triangle \mathcal{O}AB$. Clearly this must be congruent to T , and $\mathcal{O}A = \mathcal{O}B = d_1$, so we must have $AB = d_2$. The same holds for all such triangles, so we have $AB = BC = AC = d_2$. But, again, we assumed that the only triangle determined by S was isosceles, so this equilateral triangle of edge length d_2 yields a contradiction. \square

4. Proofs of lemmas

It is well known in the literature that [Lemma 2.5](#) holds, and that in \mathbb{R}^d , $d \geq 2$, a set of points determining a single distinct distance has at most $d + 1$ points. However, in the interest of completeness, we include here a proof of the lemma by induction on the dimension d :

Proof of Lemma 2.5. Base case ($d = 3$): Let $S = \{A_1, A_2, A_3, A_4, A_5\} \subset \mathbb{R}^d = \mathbb{R}^3$ be a point set containing $d + 2 = 5$ points, which defines a single distinct equilateral triangle, call it T . Thus, the triangle $\triangle A_1 A_2 A_3$ must form the triangle T . Define e to be the edge length of T , and let P denote the plane defined by $\{A_1, A_2, A_3\}$.

Since S defines an equilateral triangle, it follows that A_4 and A_5 must be equidistant from $\{A_1, A_2, A_3\}$ and lie upon a line normal to P , which goes through a point $p \in P$, where p is equidistant to $\{A_1, A_2, A_3\}$; p is called the circumcenter of the equilateral triangle $\triangle A_1 A_2 A_3$.

Since p is the circumcenter of the equilateral triangle $\triangle A_1 A_2 A_3$, it follows that

$$d(A_1, p) = d(A_2, p) = d(A_3, p) = \frac{\sqrt{3}}{3}e.$$

Since S defines only the triangle T , it follows that $d(A_4, A_5) = e$. Since A_4, A_5 , and p lie upon the same line, it follows that

$$d(A_4, p) + d(p, A_5) = d(A_4, A_5).$$

Since A_4 and A_5 are equidistant from $\{A_1, A_2, A_3\}$, they are also equidistant from the plane P , and hence, $d(A_4, p) = d(p, A_5) = \frac{1}{2}e$. Applying the Pythagorean theorem we obtain

$$d(A_4, p)^2 + d(A_1, p)^2 = d(A_4, A_1)^2,$$

which yields

$$\left(\frac{1}{2}e\right)^2 + \left(\frac{\sqrt{3}}{3}e\right)^2 = (e)^2.$$

Thus, we obtain

$$\frac{1}{4}e^2 + \frac{1}{3}e^2 = e^2,$$

which implies that $\frac{7}{12} = 1$, a clear contradiction. We note that the vertices of a 3-simplex, the regular tetrahedron, give a configuration in \mathbb{R}^3 which defines a single distinct equilateral triangle and has $3 + 1 = 4$ points. Therefore, a set S defining a single distinct equilateral triangle in \mathbb{R}^3 can have at most $3 + 1 = 4$ points.

Inductive assumption: Suppose that for dimensions $n < d$, a point set S defining a single distinct equilateral triangle can have at most $n + 1$ points.

Inductive step: Now let $S = \{A_1, \dots, A_d, A_{d+1}, A_{d+2}\} \subset \mathbb{R}^d$ be a point set containing $d + 2$ points, which defines a single distinct equilateral triangle, call it T . Thus, the points $\{A_1, \dots, A_d\}$ must be such that any triplet forms the triangle T ,

and so they must form a $(d-1)$ -simplex (this by our inductive assumption). Call this $(d-1)$ -simplex $\Delta A_1 \cdots A_d$. Let e be the edge length of T , and let P be the $(d-1)$ -dimensional hyperplane defined by $\{A_1, \dots, A_d\}$.

Since S defines an equilateral triangle, it follows that A_{d+1} and A_{d+2} must be equidistant to $\{A_1, \dots, A_d\}$ and lie upon a line normal to P , which passes through a point $p \in P$ (viz. the circumcenter of $\Delta A_1 \cdots A_d$), where p is equidistant from $\{A_1, \dots, A_d\}$.

From this point, we can follow the same construction as in the base case, using the points A_1, p, A_{d+1} , and A_{d+2} to arrive at a contradiction. We note here that the d -simplex gives a configuration in \mathbb{R}^d which defines a single distinct equilateral triangle and has $d+1$ points. Thus, a set S in \mathbb{R}^d defining a single distinct equilateral triangle can have at most $d+1$ points, as desired. \square

5. Constructions for Remark 1.3

Clearly the regular tetrahedron is an optimal configuration (of four points) determining one distinct triangle in three dimensions, and it is the unique configuration satisfying Theorem 1.1(1). Then, in the following, we will characterize the distance graphs of the four-point sets satisfying Theorem 1.1(2) (and thus optimal in three dimensions). We then provide constructions of point sets in \mathbb{R}^3 satisfying these distance graphs. For convenience, we will again consider the isosceles and scalene cases separately. Assume in all that follows that S is a set of four points determining one distinct triangle of the respective geometry.

Isosceles: Let d_1 be the repeated edge length of T . Following the framework used in the main proof, fix a point \mathcal{O} in S such that $\mathcal{O}A = d_1$ and $\mathcal{O}B = d_2$ for some A and B in S . Then, notice that given the remaining point C in S , we have $\Delta \mathcal{O}AB \simeq \Delta \mathcal{O}CB \simeq T$, so specifically $\mathcal{O}C = BC = AB = d_1$, and similarly, we must have $AC = d_2$.

Notice that any four noncoplanar points form the vertices of a tetrahedron (if the points are coplanar, this distance graph is clearly uniquely realized by the vertices of the square). And, by the above, a tetrahedron $ABCD$ satisfying the above distance graph must have $AB = d_2$ and $CD = d_2$, with the remaining edges of length d_1 . To construct such a tetrahedron, consider taking points $P = (\frac{1}{2}d_2, 0, 0)$, $Q = (-\frac{1}{2}d_2, 0, 0)$, $R = (0, \frac{1}{2}d_2, 0)$ and $S = (0, -\frac{1}{2}d_2, 0)$. Note that for $d_1 = \frac{\sqrt{2}}{2}d_2$, $PQRS$ satisfies the above distance graph, although it is planar. Then, by arbitrarily translating R and S by the same distance along the z -axis, we can construct any tetrahedron satisfying this distance graph (clearly no tetrahedron with d_1 not satisfying this inequality may exist).

Scalene: Let A, B and C determine $\Delta ABC \simeq T$ in S . Without loss of generality let $AB = d_1$, $BC = d_2$ and $AC = d_3$. Then, notice that each point may determine

each distance exactly once (otherwise they would determine an isosceles triangle, producing a contradiction similar to that used in the main proof). Then, since each point A , B and C determines exactly two of the distances, we may simply fill in that $AD = d_2$, $BD = d_3$ and $CD = d_1$. It is easy to verify that this distance graph determines only one distinct triangle.

Similarly to the isosceles case, we observe that if all four points are coplanar, this distance graph uniquely determines a rectangle. Supposing instead that four points A , B , C and D satisfy the above distance graph and are noncoplanar (thus form a tetrahedron), we clearly must then have that, up to relabeling, $AB = CD = d_1$, $AC = BD = d_2$ and $AD = BC = d_3$. Namely, “opposite” pairs of edges of the tetrahedron are congruent. To construct all such tetrahedra, consider fixing points $P = (\frac{1}{2}d_1, 0, 0)$ and $Q = (-\frac{1}{2}d_1, 0, 0)$. Then, fix $R' = (0, \frac{1}{2}d_1, z')$ and $S' = (0, -\frac{1}{2}d_1, z')$ for arbitrary z' . Consider the unique circle lying in the plane $z = z'$ passing through R' and S' . Then choose R and S as the endpoints of any diameter of this circle such that $R, S \neq R', S'$ and P, Q, R and S are not all coplanar. Clearly, then, $PQRS$ satisfies the above distance graph.

Acknowledgements

We would like to thank an anonymous referee for suggesting the inclusion of [Theorem 1.1](#) in its current form. Their proposed restructuring around this result has significantly improved the paper.

References

- [Brass et al. 2005] P. Brass, W. Moser, and J. Pach, *Research problems in discrete geometry*, Springer, 2005. [MR](#) [Zbl](#)
- [Brenner et al. 2019] H. N. Brenner, J. S. Depret-Guillaume, E. A. Palsson, and S. Senger, “Uniqueness of optimal point sets determining two distinct triangles”, preprint, 2019. [arXiv](#)
- [Epstein et al. 2018] A. Epstein, A. Lott, S. J. Miller, and E. A. Palsson, “Optimal point sets determining few distinct triangles”, *Integers* **18** (2018), art. id. A16. [MR](#) [Zbl](#)
- [Erdős 1946] P. Erdős, “On sets of distances of n points”, *Amer. Math. Monthly* **53** (1946), 248–250. [MR](#) [Zbl](#)
- [Erdős 1975] P. Erdős, “On some problems of elementary and combinatorial geometry”, *Ann. Mat. Pura Appl.* (4) **103** (1975), 99–108. [MR](#) [Zbl](#)
- [Erdős and Fishburn 1996] P. Erdős and P. Fishburn, “Maximum planar sets that determine k distances”, *Discrete Math.* **160**:1-3 (1996), 115–125. [MR](#) [Zbl](#)
- [Guth and Katz 2015] L. Guth and N. H. Katz, “On the Erdős distinct distances problem in the plane”, *Ann. of Math.* (2) **181**:1 (2015), 155–190. [MR](#) [Zbl](#)
- [Shinohara 2008] M. Shinohara, “Uniqueness of maximum planar five-distance sets”, *Discrete Math.* **308**:14 (2008), 3048–3055. [MR](#) [Zbl](#)
- [Wei 2012] X. Wei, “A proof of Erdős–Fishburn’s conjecture for $g(6) = 13$ ”, *Electron. J. Combin.* **19**:4 (2012), art. id. P38. [MR](#) [Zbl](#)

Received: 2019-02-12

Revised: 2019-09-29

Accepted: 2019-11-11

hazlebrenner@vt.edu

*Department of Mathematics, Virginia Tech, Blacksburg, VA,
United States*

jdg@vt.edu

*Department of Mathematics, Virginia Tech, Blacksburg, VA,
United States*

palsson@vt.edu

*Department of Mathematics, Virginia Tech, Blacksburg, VA,
United States*

rstuckel@kent.edu

*Department of Mathematics, Virginia Tech, Blacksburg, VA,
United States*

Current address:

*Department of Mathematical Sciences, Kent State University,
Kent, OH, United States*

INVOLVE YOUR STUDENTS IN RESEARCH

Involve showcases and encourages high-quality mathematical research involving students from all academic levels. The editorial board consists of mathematical scientists committed to nurturing student participation in research. Bridging the gap between the extremes of purely undergraduate research journals and mainstream research journals, *Involve* provides a venue to mathematicians wishing to encourage the creative involvement of students.

MANAGING EDITOR

Kenneth S. Berenhaut Wake Forest University, USA

BOARD OF EDITORS

Colin Adams	Williams College, USA	Robert B. Lund	Clemson University, USA
Arthur T. Benjamin	Harvey Mudd College, USA	Gaven J. Martin	Massey University, New Zealand
Martin Bohner	Missouri U of Science and Technology, USA	Mary Meyer	Colorado State University, USA
Amarjit S. Budhiraja	U of N Carolina, Chapel Hill, USA	Frank Morgan	Williams College, USA
Pietro Cerone	La Trobe University, Australia	Mohammad Sal Moslehian	Ferdowsi University of Mashhad, Iran
Scott Chapman	Sam Houston State University, USA	Zuhair Nashed	University of Central Florida, USA
Joshua N. Cooper	University of South Carolina, USA	Ken Ono	Univ. of Virginia, Charlottesville
Jem N. Corcoran	University of Colorado, USA	Yuval Peres	Microsoft Research, USA
Toka Diagana	University of Alabama in Huntsville, USA	Y.-F. S. Pétermann	Université de Genève, Switzerland
Michael Dorff	Brigham Young University, USA	Jonathon Peterson	Purdue University, USA
Sever S. Dragomir	Victoria University, Australia	Robert J. Plemmons	Wake Forest University, USA
Joel Foisy	SUNY Potsdam, USA	Carl B. Pomerance	Dartmouth College, USA
Erin W. Fulp	Wake Forest University, USA	Vadim Ponomarenko	San Diego State University, USA
Joseph Gallian	University of Minnesota Duluth, USA	Bjorn Poonen	UC Berkeley, USA
Stephan R. Garcia	Pomona College, USA	József H. Przytycki	George Washington University, USA
Anant Godbole	East Tennessee State University, USA	Richard Rebarber	University of Nebraska, USA
Ron Gould	Emory University, USA	Robert W. Robinson	University of Georgia, USA
Sat Gupta	U of North Carolina, Greensboro, USA	Javier Rojo	Oregon State University, USA
Jim Haglund	University of Pennsylvania, USA	Filip Saidak	U of North Carolina, Greensboro, USA
Johnny Henderson	Baylor University, USA	Hari Mohan Srivastava	University of Victoria, Canada
Glenn H. Hurlbert	Virginia Commonwealth University, USA	Andrew J. Sterge	Honorary Editor
Charles R. Johnson	College of William and Mary, USA	Ann Trenk	Wellesley College, USA
K. B. Kulasekera	Clemson University, USA	Ravi Vakil	Stanford University, USA
Gerry Ladas	University of Rhode Island, USA	Antonia Vecchio	Consiglio Nazionale delle Ricerche, Italy
David Larson	Texas A&M University, USA	John C. Wierman	Johns Hopkins University, USA
Suzanne Lenhart	University of Tennessee, USA	Michael E. Zieve	University of Michigan, USA
Chi-Kwong Li	College of William and Mary, USA		

PRODUCTION

Silvio Levy, Scientific Editor

Cover: Alex Scorpan

See inside back cover or msp.org/involve for submission instructions. The subscription price for 2020 is US \$205/year for the electronic version, and \$275/year (+\$35, if shipping outside the US) for print and electronic. Subscriptions, requests for back issues and changes of subscriber address should be sent to MSP.

Involve (ISSN 1944-4184 electronic, 1944-4176 printed) at Mathematical Sciences Publishers, 798 Evans Hall #3840, c/o University of California, Berkeley, CA 94720-3840, is published continuously online. Periodical rate postage paid at Berkeley, CA 94704, and additional mailing offices.

Involve peer review and production are managed by EditFlow® from Mathematical Sciences Publishers.

PUBLISHED BY

 **mathematical sciences publishers**

nonprofit scientific publishing

<http://msp.org/>

© 2020 Mathematical Sciences Publishers

involve

2020

vol. 13

no. 1

Structured sequences and matrix ranks	1
CHARLES JOHNSON, YAOXIAN QU, DUO WANG AND JOHN WILKES	
Analysis of steady states for classes of reaction-diffusion equations with hump-shaped density-dependent dispersal on the boundary	9
QUINN MORRIS, JESSICA NASH AND CATHERINE PAYNE	
The L-move and Markov theorems for trivalent braids	21
CARMEN CAPRAU, GABRIEL COLOMA AND MARGUERITE DAVIS	
Low stages of the Taylor tower for r-immersions	51
BRIDGET SCHREINER, FRANJO ŠARČEVIĆ AND ISMAR VOLIĆ	
A new go-to sampler for Bayesian probit regression	77
SCOTT SIMMONS, ELIZABETH J. MCGUFFEY AND DOUGLAS VANDERWERKEN	
Characterizing optimal point sets determining one distinct triangle	91
HAZEL N. BRENNER, JAMES S. DEPRET-GUILLAUME, EYVINDUR A. PALSSON AND ROBERT W. STUCKEY	
Solutions of periodic boundary value problems	99
R. AADITH, PARAS GUPTA AND JAGAN MOHAN JONNALAGADDA	
A few more trees the chromatic symmetric function can distinguish	109
JAKE HURYN AND SERGEI CHMUTOV	
One-point hyperbolic-type metrics	117
MARINA BOROVIKOVA, ZAIR IBRAGIMOV, MIGUEL JIMENEZ BRAVO AND ALEXANDRO LUNA	
Some generalizations of the ASR search algorithm for quasitwisted codes	137
NUH AYDIN, THOMAS H. GUIDOTTI, PEIHAN LIU, ARMIYA S. SHAIKH AND ROBERT O. VANDENBERG	
Continuous factorization of the identity matrix	149
YUYING DAI, ANKUSH HORE, SIQI JIAO, TIANXU LAN AND PAVLOS MOTAKIS	
Almost excellent unique factorization domains	165
SARAH M. FLEMING AND SUSAN LOEPP	