





Generalizations of k-dimensional Weisfeiler–Leman stabilization

Anuj Dawar and Danny Vagnozzi

The family of Weisfeiler–Leman equivalences on graphs is a widely studied approximation of graph isomorphism with many different characterizations. We study these and other approximations of isomorphism defined in terms of refinement operators and Schurian polynomial approximation schemes (SPAS). The general framework of SPAS allows us to study a number of parameters of the refinement operators based on Weisfeiler–Leman refinement, logic with counting, lifts of Weisfeiler–Leman as defined by Evdokimov and Ponomarenko, the invertible map test introduced by Dawar and Holm, and variations of these, as well as to establish relationships between them.

1. Introduction

For convenience, we shall treat graphs as arc-coloured complete digraphs; that is to say, as labelled partitions of the set of ordered pairs of vertices (hereafter, we refer to the latter as arcs). For example, an undirected simple graph can be seen as a partition of its arcs into edges, nonedges, and loops. As such, the graph isomorphism problem is that of deciding whether there is a colour-preserving bijection between the sets of vertices of two given graphs. Computationally, this problem is polynomial-time equivalent to finding the orbits of the induced action of the automorphism group of a given graph G on a fixed power of its vertex set V [Mathon 1979]. For short, we refer to the partition of V^k (for any fixed k) obtained by this action as the orbit partition of V^k . The graph isomorphism problem (and likewise the problem of determining the orbit partition) is neither known to be solvable in polynomial time nor known to be NP-complete. The best known upper bound to their computational time is quasipolynomial. This follows from the well-known result by Babai [2016].

The classical Weisfeiler–Leman (WL) algorithm is a well-known method for approximating the orbits of the induced action of the automorphism group of a given graph on the set of pairs of arcs. It can be seen as a generalization of the so-called *naïve colour refinement*. Given a graph G, the WL algorithm produces a coherent configuration, which is a partition of the set of arcs of G satisfying certain stability conditions (see Section 2 for the definition). A natural generalization of this algorithm was given by Babai: for each $k \in \mathbb{N}$, the k-dimensional Weisfeiler–Leman (WL $_k$) algorithm outputs a labelled partition of k-tuples of vertices satisfying a similar stability condition and respecting local isomorphism. The running time of the WL $_k$ algorithm on a graph with n vertices is bounded by $n^{O(k)}$. The case k=1 coincides with the naïve colour refinement, and k=2 with the classical Weisfeiler–Leman algorithm.

It follows from a result by Cai, Fürer, and Immerman [Cai et al. 1992] that there is no fixed $k \in \mathbb{N}$ such that for all graphs the k-dimensional Weisfeiler–Leman algorithm outputs the partition of k-tuples of vertices into orbits of the induced action of the automorphism group of the input graph. Indeed, the

MSC2010: primary 05E15, 05E30, 05E99; secondary 03D15.

Keywords: Weisfeiler-Leman, combinatorics, invertible map, graph isomorphism, complexity theory, coherent configuration.

authors show how to construct a graph with O(k) vertices for which WL_k fails to produce the partition into such orbits. Thus, their result implies that a partition induced by this group action can be obtained for all graphs on n vertices only if one chooses k to be $\Omega(n)$. One can informally claim that the strength of the k-dimensional Weisfeiler-Leman algorithm increases with k. More precisely, for unlabelled partitions \mathcal{P} and \mathcal{Q} of some set A we write $\mathcal{P} \leq_A \mathcal{Q}$, and say \mathcal{Q} is a refinement of \mathcal{P} ; if, whenever $a, b \in A$ are in the same equivalence class of \mathcal{Q} , they are also in the same equivalence class of \mathcal{P} . By viewing a labelled partition of A as a function $\gamma: A \to L$ to a set of labels L (which we sometimes refer to as the *colour* set), the unlabelled partition induced by γ is $\{\gamma^{-1}(l) \mid l \in L\}$. We extend the partial order \leq_A to labelled partitions by writing $\gamma \leq_A \rho$ to mean that the unlabelled partition induced by ρ refines that induced by γ . Note that this does not require that the codomains of γ and ρ are the same. We omit the subscript A when the set is clear from the context. For a graph Γ , define $\overline{\operatorname{WL}}_1(\Gamma) = \Gamma$; and for $k \geq 2$, set $\overline{\operatorname{WL}}_k(\Gamma)$ to be the labelled partition of the set of arcs induced by the output of the k-dimensional Weisfeiler-Leman algorithm on input Γ . We can now state the following:

$$\overline{WL}_1(\Gamma) \leq \overline{WL}_2(\Gamma) \leq \cdots \leq \overline{WL}_n(\Gamma) = \overline{WL}_{n+1}(\Gamma) = \cdots = \overline{WL}_{\infty}(\Gamma),$$

where n is the number of vertices of Γ and $\overline{WL}_{\infty}(\Gamma)$ is the partition into the orbits of the induced action of the automorphism group on arcs. Also,

$$\overline{\mathrm{WL}}_l(\overline{\mathrm{WL}}_k(\Gamma)) = \overline{\mathrm{WL}}_k(\Gamma)$$

for all $l, k \in \mathbb{N}$ with $l \leq k$. This shows that the family of maps from the set of arc-coloured complete digraphs to itself $\{\overline{WL}_1, \overline{WL}_2, \dots\}$ forms a *Schurian polynomial approximation scheme* in the following sense, as defined in [Evdokimov and Ponomarenko 1999]:

Definition 1.1 (Schurian polynomial approximation scheme). A family of mappings $\{X_1, X_2, ...\}$ is said to form a *Schurian polynomial approximation scheme* (SPAS) if for any graph Γ with vertex set V:

- (1) $X_k(\Gamma)$ is a graph with vertex set V for all $k \in \mathbb{N}$.
- (2) $X_1(\Gamma) \leq X_2(\Gamma) \leq \cdots \leq X_n(\Gamma) = X_{n+1}(\Gamma) = \cdots = X_{\infty}(\Gamma) = \text{Sch}(\Gamma)$, where n = |V| and $\text{Sch}(\Gamma)$ is the partition of arcs into orbits of the induced action of the automorphism group of Γ .
- (3) $X_l(X_m(\Gamma)) = X_m(\Gamma)$ for all $l, m \in \mathbb{N}$ with $l \leq m$.
- (4) $X_k(\Gamma)$ is computable in time $n^{O(k)}$.

More informally, one can think of SPAS as a collection of polynomially computable functions indexed by \mathbb{N} , each of which takes as input a partition and refines it. Moreover, the higher the index, the closer the output partition is to the orbit partition of the set of arcs.

Definition 1.2 (dominance and equivalence of SPAS). For any two SPAS, $S_X = \{X_1, X_2, ...\}$ and $S_Y = \{Y_1, Y_2, ...\}$:

- (1) we say S_X dominates S_Y , and write $S_Y \leq S_X$, if for each $k \in \mathbb{N}$ there is some $k' \in \mathbb{N}$ such that $Y_k(\Gamma) \leq X_{k'}(\Gamma)$ for all graphs Γ .
- (2) we say S_X is *equivalent* to S_Y , and write $S_X \simeq S_Y$, if $S_X \preceq S_Y$ and $S_Y \preceq S_X$.
- (3) we say S_X strictly dominates S_Y if $S_Y \leq S_X$, but $S_X \not\leq S_Y$.

In this paper, we deal with the following SPAS: S_{WL} , S_C , $S_{C,r}$, $S_{WL,r}$, and $S_{IM(\mathbb{F})}$. Such schemes arise from considering fixed points of *refinement operators*. The formal definition of *refinement operators* is given in Section 4. These are operators that take a partition γ , of V^k , to a refinement of itself. The SPAS S_{WL} and S_C arise from the following well-known concepts: Babai's generalization of the classical Weisfeiler–Leman algorithm for the former and first order logic with counting quantifiers for the latter. In each case, the label of a tuple $\vec{v} \in V^k$ in the refined partition is determined by its label in γ and the partition of V that is induced by considering $\gamma(\vec{u})$ for the tuples \vec{u} obtained by substituting elements of V in \vec{v} . For $v \in \mathbb{N}$, $v_{WL,r}$ and $v_{C,r}$ are further generalizations of v_{WL} and $v_{C,r}$ respectively. In these generalizations, the label associated to each $v_{C,r}$ are further generalization by $v_{C,r}$ and the partition of $v_{C,r}$ obtained by considering $v_{C,r}$ for tuples $v_{C,r}$ obtained by substituting $v_{C,r}$ and definitions of these are given later. Here we note that our first result shows that the parameter $v_{C,r}$ does not strengthen the SPAS $v_{C,r}$ and $v_{C,r}$

Theorem 1.3. For any $r \in \mathbb{N}$, $S_{WL} \simeq S_C \simeq S_{WL,r} \simeq S_{C,r}$.

The reasons for considering the additional parameter r is that it appears to be of interest in another scheme we consider. The scheme $\mathcal{S}_{\text{IM}(\mathbb{F})}$ arises from the *invertible map game* introduced in [Dawar and Holm 2017]. It has been shown to have a close relationship to logics with linear algebraic operators [Dawar et al. 2019] over a field \mathbb{F} . The associated refinement operator $\text{IM}_k^{\mathbb{F}}$ maps each k-tuple of vertices \vec{v} and partition γ to a tuple of matrices, and the colour associated to \vec{v} by the refinement of γ is determined by the equivalence class of this tuple of matrices under simultaneous similarity.

For this SPAS we prove the following results:

Theorem 1.4. For any field \mathbb{F} of characteristic 0, $\mathcal{S}_{\text{IM}(\mathbb{F})} \simeq \mathcal{S}_{\text{WL}}$.

Theorem 1.5. For any field \mathbb{F} of positive characteristic, $S_{\text{IM}(\mathbb{F})}$ strictly dominates S_{WL} .

The paper is structured as follows: after a brief overview of the required notions on coherent configurations and algebras, we formally define and discuss the concepts of refinement operators and procedures. We then prove Theorem 1.3 and use a similar method to show the equivalence between \mathcal{S}_{WL} and \mathcal{S}_{EP} , a SPAS introduced by Evdokimov and Ponomarenko in [1999]. The final two sections contain the proofs of Theorems 1.4 and 1.5 and a short discussion on a variant of the SPAS $\mathcal{S}_{IM(\mathbb{F})}$, namely $\mathcal{S}_{IMt(\mathbb{F})}$. This variant is motivated by looking at the difference between the definitions of \mathcal{S}_{WL} and \mathcal{S}_{C} , and applying a similar variation to the definition of $\mathcal{S}_{IM(\mathbb{F})}$. The discussion leads to a proof of the following:

Theorem 1.6. For any field \mathbb{F} , $S_{IM(\mathbb{F})} \simeq S_{IMt(\mathbb{F})}$.

Throughout the text, all sets are finite. Given two sets V and I, a tuple in V^I is denoted by \vec{v} , and its i-th entry by v_i , for each $i \in I$. We use the notation $(v_i)_{i \in I}$ to denote the element of V^I with i-th element equal to v_i . We set $[k] = \{1, 2, \ldots, k\} \subset \mathbb{N}$ and $[k]^{(r)} = \{\vec{x} \in [k]^r \mid x_i \neq x_j \ \forall i, j \in [r], i \neq j\}$. Recall that a labelled partition of a set A is a function $\gamma : A \to \operatorname{Im} \gamma$. The class of all labelled partitions of A is denoted by $\mathcal{P}(A)$. Recall also that for $\gamma, \rho \in \mathcal{P}(A), \gamma \leq \rho$ denotes that the unlabelled partition induced by ρ is a refinement of that induced by γ . If $\gamma \leq \rho$ and $\gamma \succeq \rho$ are both satisfied, we write $\gamma \approx \rho$. Note that $\gamma \approx \rho$ does not imply $\gamma = \rho$, as they may have different codomains. The equivalence class of $\alpha \in A$ with respect to the partition γ is denoted by $[\alpha]_{\gamma}$. Fix some set V and $k, r \in \mathbb{N}$ with $r \leq k$. For

any $\vec{v} \in V^k$, $\vec{u} \in V^r$, and $\vec{i} \in [k]^{(r)}$ we define $\vec{v}(\vec{i}, \vec{u}) \in V^k$ to be the tuple with entries

$$(\vec{v}\langle \vec{i}, \vec{u}\rangle)_j = \begin{cases} u_{i_s} & \text{if } j = i_s \text{ for some } s \in [r], \\ v_j & \text{otherwise.} \end{cases}$$

Given two tuples $\vec{v} \in V^r$ and $\vec{w} \in V^s$, their *concatenation* is denoted by $\vec{v} \cdot \vec{w} \in V^{r+s}$. More precisely, $\vec{v} \cdot \vec{w}$ is the tuple with entries

$$(\vec{v} \cdot \vec{w})_i = \begin{cases} v_i & \text{if } i \in [r], \\ w_j & \text{if } i = j + r. \end{cases}$$

For a relation $R \subseteq V^2$, we define the adjacency matrix of R to be the $V \times V$ matrix whose (u, v)-entry is 1 if $(u, v) \in R$ and 0 otherwise. The set of multisets of elements of V is denoted by $\operatorname{Mult}(V)$, and the multiset of entries of a tuple $\vec{v} \in V^I$ is denoted by $\{\{v_i \mid i \in I\}\}$. For all $\gamma \in \mathcal{P}(V^k)$ and natural numbers $r \leq k$, set $\Phi^{\gamma,r} = \operatorname{Im} \gamma^{[k]^{(r)}}$.

2. Coherent configurations and coherent algebras

This section introduces notions on coherent configurations and algebras necessary throughout the paper. For a more in-depth account see [Chen and Ponomarenko 2019] or [Cameron 1999]. Our formulation is, in general, different from the more traditional treatment, as we deal with labelled partitions and extend the notion of coherent algebras to arbitrary fields. Also note that rainbows and coherent configurations were originally defined for *unlabelled partitions*. Thus, Definitions 2.1 and 2.2 define, strictly speaking, a *labelled* rainbow and a *labelled* coherent configuration, respectively.

Definition 2.1 (rainbow). A labelled partition ρ of V^2 is said to be a *rainbow* on V if:

(1) There is a set $\mathcal{I} \subseteq \operatorname{Im} \rho$ such that

$$\bigcup_{\sigma \in \mathcal{T}} \{ \vec{x} \in V^2 \mid \rho(\vec{x}) = \sigma \} = \{ (v, v) \in V^2 \mid v \in V \}.$$
 (1)

(2) For all $(u, v), (u', v') \in V^2, \ \rho(u, v) = \rho(u', v') \iff \rho(v, u) = \rho(v', u').$

We set $Cel(\rho) = \{U \subset V \mid \exists \sigma \in Im \ \rho, \ \rho(u, u) = \sigma, \ \forall u \in U\}$ and call its elements the *cells* of ρ .

It was stated in the introduction that, in this paper, graphs are viewed as partitions of the set of their arcs, hence as arc-coloured complete digraphs. For example, an uncoloured loop-free undirected graph can be seen as a complete digraph with its arcs partitioned into three colour classes: edges, nonedges, and loops. Hence, we can always see a graph as a rainbow in the above sense. This view is natural, since our interest is in the partition into orbits of the induced action of the automorphism group on arcs, and this partition is, necessarily, a rainbow. Furthermore, given any group action on V, the partition into the orbits of the induced action on V^2 forms a *coherent configuration* [Cameron 1999]¹:

Definition 2.2 (coherent configuration). A rainbow ρ on V is said to be a *coherent configuration* on V if for each σ , τ , $\kappa \in \text{Im } \rho$, there is a constant $p_{\sigma\tau}^{\kappa}$ such that for any $(u, v) \in V^2$ with $\rho(u, v) = \kappa$,

$$\big|\{x\in V\mid \rho(u,x)=\sigma, \rho(x,v)=\tau\}\big|=p_{\sigma\tau}^{\kappa}.$$

¹Although all group actions give rise to coherent configurations, not all coherent configurations arise from group actions; see [Cameron 1999].

Observe that if X is a union of cells of a coherent configuration, the restriction $\rho|_{X^2}$ is a coherent configuration on X.

The constants $p_{\sigma\tau}^{\kappa}$ are called the *intersection numbers* of ρ and may be interpreted algebraically as follows: For every $\sigma \in \operatorname{Im} \rho$, let A_{σ} be the adjacency matrix of the relation $\rho^{-1}(\sigma)$. Then for all $\sigma, \tau \in \operatorname{Im} \rho$,

$$A_{\sigma}A_{\tau} = \sum_{\kappa \in \operatorname{Im} \rho} p_{\sigma\tau}^{\kappa} A_{\kappa}.$$

Thus, taking $p_{\sigma\tau}^{\kappa}$ as rational numbers in a field \mathbb{F} of characteristic zero, we see that the \mathbb{F} -span of the set $\mathcal{A}_{\rho} = \{A_{\sigma} \mid \sigma \in \text{Im } \rho\}$ is an \mathbb{F} -algebra. The same is true if we take \mathbb{F} to be a field of characteristic q and consider the constants $p_{\sigma\tau}^{\kappa}$ modulo q. We refer to this algebra as the \mathbb{F} -adjacency algebra of ρ and denote it by $\mathbb{F}\mathcal{A}_{\rho}$. Such an algebra is a *coherent algebra* in the following sense:

Definition 2.3 (coherent algebra). A subalgebra of $Mat_V(\mathbb{F})$ is said to be a *coherent algebra* on V if it is a unital algebra with respect to matrix multiplication and Schur–Hadamard (componentwise) multiplication , and it is closed under transposition.

We indicate the Schur–Hadamard multiplication by \star ². At this point, it needs to be pointed out that in most literature, when $\mathbb{F} = \mathbb{C}$, closure under transposition is usually replaced by *closure under Hermitian conjugation* for the definition of a coherent algebra. However, we show in Proposition 2.6, that an algebra satisfies Definition 2.3 if and only if it has a basis of 0-1-matrices satisfying the *coherence conditions* (Definition 2.4). In Section 2.3 of [Chen and Ponomarenko 2019] it is shown that an algebra over \mathbb{C} satisfies Definition 2.3 with closure under transposition replaced by closure under Hermitian conjugation if and only if it has a basis of 0-1-matrices satisfying the coherence conditions. Hence, over \mathbb{C} , Definition 2.3 is equivalent to the original one by D. Higman in [1987], but has the advantage that it can be extended to any field.

It is clear from the definition, that for any coherent configuration ρ , the \mathbb{F} -adjacency algebra $\mathbb{F}\mathcal{A}_{\rho}$ is a coherent algebra for any field \mathbb{F} . Indeed, the set \mathcal{A}_{ρ} is the unique basis of 0-1-matrices for $\mathbb{F}\mathcal{A}_{\rho}$ satisfying the *coherence conditions*:

Definition 2.4. A set of 0-1-matrices \mathcal{M} is said to satisfy the *coherence conditions* if:

- (1) $\sum_{A \in \mathcal{M}} A = \mathbb{J}$, where \mathbb{J} is the all 1s matrix.
- (2) For some $\mathcal{I} \subseteq \mathcal{M}$, $\sum_{A \in \mathcal{I}} A = \mathbb{I}$, where \mathbb{I} is the identity matrix.
- (3) $A^t \in \mathcal{M}$ for all $A \in \mathcal{M}$.

We now show that any coherent algebra over any field has a unique basis of 0-1-matrices satisfying the coherence conditions. We refer to this basis as the *standard basis* of a coherent algebra. The argument that follows is analogous to that used to prove Theorem 2.3.7 in [Chen and Ponomarenko 2019].

Let W be a coherent algebra over V. As explained in Section 2.3 in [Chen and Ponomarenko 2019], one may write

$$\mathbb{J} = \sum_{i \in [r]} E_i,\tag{2}$$

²In the literature, the Schur–Hadamard multiplication is often denoted by ∘. However, we reserve the latter for function composition.

where $\{E_i \mid i \in [r]\}$ is the full set of primitive idempotents of W with respect to the Schur–Hadamard product. In order to be an idempotent, E_i must be a 0-1 matrix; and hence, the adjacency matrix of some relation $R_i \subseteq V^2$. Since for $i \neq j$, E_i and E_j are orthogonal, for all $u, v \in V$, $(E_i)_{uv} = 1 \Longrightarrow (E_j)_{uv} = 0$. Thus, from (2), it follows that $\{R_i \mid i \in [r]\}$ forms a partition of V^2 .

Lemma 2.5. $\{E_i \mid i \in [r]\}$, as above, satisfies the coherence conditions.

Proof. Condition (1) of Definition 2.4 is satisfied because of (2). Because $\mathbb{I} \in W$ is an idempotent, it can be written as a sum of primitive idempotents. Thus, $\{E_i \mid i \in [r]\}$ satisfies condition (2) in Definition 2.4. Finally, E_i^t is also a primitive idempotent, since W is closed under transposition.

Proposition 2.6. For any field \mathbb{F} , a coherent algebra on V over \mathbb{F} has a unique basis of 0-1-matrices satisfying the coherence conditions.

Proof. The set $B = \{E_i \mid i \in [r]\}$ satisfies the coherence conditions by Lemma 2.5.

Suppose \mathbb{F} is algebraically closed. Then B is a basis for W, since W is commutative with respect to the Schur–Hadamard product and a basis of a semisimple commutative algebra over an algebraically closed field is given by the set of its primitive idempotents.

Suppose \mathbb{F} is not algebraically closed. Since B is a linearly independent set, there is some $B' \subseteq W$ such that $B \cup B'$ is a basis for W. Let \mathbb{G} be the algebraic closure of \mathbb{F} , and consider the linear space $\mathbb{G}(B \cup B') \subseteq \operatorname{Mat}_V(\mathbb{G})$. By construction, $\mathbb{G}(B \cup B')$ is closed under transposition, matrix multiplication, and Schur–Hadamard multiplication and is thus a coherent algebra over \mathbb{G} . From the above, $\mathbb{G}(B \cup B')$ must then have a basis B'' of 0-1-matrices satisfying the coherence conditions. Since all entries of the elements of B'' are 0 and 1, $B'' \subset \mathbb{F}(B \cup B')$ is a basis for W as well. As B'' is a set of primitive orthogonal idempotents of W, it holds that $B'' \subseteq B$. But B is a linearly independent set, and B'' is a basis for W. Whence $B'' = B = \{E_i \mid i \in [r]\}$.

The uniqueness of B follows from (2), which implies that any basis of 0-1 matrices satisfying the coherence conditions must be the set of primitive idempotents of W.

We can denote a coherent algebra on V over \mathbb{F} as $\mathbb{F}\mathcal{A}$, where \mathcal{A} is some set of 0-1-matrices satisfying the coherence conditions. It is easily seen that for any $(u, v) \in V^2$ and binary relations S and T on V with adjacency matrices A_S , $A_T \in \mathrm{Mat}_V(\mathbb{F})$, respectively,

$$\left| \{ x \in V \mid (u, x) \in S, (x, v) \in T \} \right| = (A_S A_T)_{uv} \tag{3}$$

if $\operatorname{char}(\mathbb{F})=0$. Set $\rho_W:V^2\to [r]$ to be $\rho_W(u,v)=i$ if $(u,v)\in R_i$. Since there are constants $p_{ij}^k\in\mathbb{F}$ such that

$$E_i E_j = \sum_{k \in [r]} p_{ij}^k E_k,$$

it then follows that

$$p_{ij}^{k} = \left| \{ x \in V \mid \rho_{W}(u, x) = i, \rho_{W}(x, v) = j \} \right|$$
 (4)

is the same for all (u, v) such that $\rho_W(u, v) = k$; and hence, ρ_W is a coherent configuration. Otherwise, if \mathbb{F} has characteristic q > 0, (3) holds modulo q. In particular if q > |V|, ρ_W is a coherent configuration.

Remark 2.7. In the literature, coherent algebras over a field \mathbb{F} of positive characteristic have usually been defined to be the \mathbb{F} -span of the adjacency matrices of the relations $\{\rho^{-1}(\sigma)|\sigma\in\operatorname{Im}\rho\}$ for some

coherent configuration ρ . The latter discussion thus shows that, over fields of positive characteristic, Definition 2.2 defines a potentially larger class of algebras.

The most intuitive morphisms between coherent configurations arise from the algebraic setting. Let $\mathcal{A}, \mathcal{A}' \subseteq \operatorname{Mat}_V(\mathbb{F})$ satisfy the coherence conditions.

Definition 2.8 (isomorphism of coherent algebras). An \mathbb{F} -linear bijection $\psi : \mathbb{F} \mathcal{A} \to \mathbb{F} \mathcal{A}'$ is said to be an *isomorphism* of coherent algebras if:

- (1) $\psi(\mathbb{I}) = \mathbb{I}$.
- (2) $\psi(\mathbb{J}) = \mathbb{J}$.
- (3) $\psi(AB) = \psi(A)\psi(B)$ and $\psi(A \star B) = \psi(A) \star \psi(B)$, for all $A, B \in \mathbb{F} \mathcal{A}_{\rho}$.

That is, ψ preserves the structure of $\mathbb{F}\mathcal{A}$, both as a matrix algebra and as an algebra, with respect to \star . As a consequence, the image under ψ of an element of the standard basis of $\mathbb{F}\mathcal{A}$ must be an element of the standard basis of $\mathbb{F}\mathcal{A}$, since the standard basis of a coherent algebra is the set of its primitive idempotents with respect to \star . Conversely, the \mathbb{F} -linear extension of any bijection between the standard bases of $\mathbb{F}\mathcal{A}$ and $\mathbb{F}\mathcal{A}'$ is a coherent algebra isomorphism provided it is also a matrix algebra isomorphism.

Let ρ and ρ' be coherent configurations and denote their intersection numbers by $p_{\sigma\tau}^{\kappa}$ and $q_{\sigma'\tau'}^{\kappa'}$, respectively.

Definition 2.9 (algebraic isomorphism). A bijection $\phi : \operatorname{Im} \rho \to \operatorname{Im} \rho'$ is said to be an *algebraic isomorphism* if for all $\sigma, \tau, \kappa \in \operatorname{Im} \rho$,

$$p_{\sigma\tau}^{\kappa} = q_{\phi(\sigma)\phi(\tau)}^{\phi(\kappa)}.$$

Thus, an algebraic isomorphism between coherent configurations induces a bijection between the standard bases of their respective adjacency algebras. Such a bijection linearly extends to a coherent algebra isomorphism.

Crucial to this paper is the fact that when coherent algebras are semisimple (with respect to matrix product), isomorphisms between them assume a very simple form. Indeed, it is an easy consequence of the Skolem–Nöther Theorem that if $\psi: W_1 \to W_2$ is an algebra isomorphism, where W_1 and W_2 are semisimple subalgebras of $\mathrm{Mat}_n(\mathbb{F})$, then there is some $S \in \mathrm{GL}_n(\mathbb{F})$ such that $\psi(A) = SAS^{-1}$ for all $A \in W_1$. The following is a direct consequence of Theorem 4.1.3 in [Zieschang 1996]³:

Theorem 2.10. The Jacobson radical of a coherent algebra $\mathbb{F}A$ is a subspace of the span of the elements of the standard basis whose number of nonzero entries is divisible by $char(\mathbb{F})$.

For a coherent configuration ρ on V, choose $u, v \in V$ such that $\rho(u, v) = \sigma$ and $\rho(u, u) = \tau$. It follows from formula (2.1.5) in [Chen and Ponomarenko 2019] that

$$|\rho^{-1}(\sigma)| = |\rho^{-1}(\tau)| |\{x \in V \mid \rho(u, x) = \sigma\}|.$$

Since both factors on the right-hand side are no larger than |V|, it is clear that the prime factors of the size of any equivalence class of ρ is no larger than |V|.

³The author actually proves this statement for coherent algebras whose diagonal matrices are multiples of the identity matrix. However, the same argument applies to the more general case; and, in particular, to our more general notion of coherent algebras in the sense of Definition 2.2.

Corollary 2.11. A coherent algebra on V over \mathbb{F} is semisimple with respect to matrix product if $\operatorname{char}(\mathbb{F}) = 0$ or $\operatorname{char}(\mathbb{F}) > |V|$.

3. Graph-like partitions

In this section, we describe some restrictions to be imposed on the partitions dealt with in the paper. Such restrictions are natural in the sense that they are necessary conditions to be satisfied by a partition of k-tuples into the orbits of an induced action of a group on 1-tuples.

Fix some $k \in \mathbb{N}$, a set V, and let $\gamma \in \mathcal{P}(V^k)$. Define an action of $\mathrm{Sym}(k)$ on V^k by letting, for each $\tau \in \mathrm{Sym}(k)$, \vec{v}^{τ} be the element of V^k with i-th entry $v_{\tau^{-1}(i)}$.

Definition 3.1 (invariance). γ is *invariant* if $\gamma(\vec{u}) = \gamma(\vec{v}) \Longrightarrow \gamma(\vec{u}^{\tau}) = \gamma(\vec{v}^{\tau})$ for all $\vec{u}, \vec{v} \in V^k$ and all $\tau \in \operatorname{Sym}(k)$.

Fix some $r \in [k]$ and $\vec{i} \in [k]^r$. For a tuple $\vec{v} \in V^k$, we define its *projection* on \vec{i} , denoted $\operatorname{pr}_{\vec{i}}\vec{v}$, to be the tuple in V^r with j-th entry v_{i_j} . Without ambiguity, we write $\operatorname{pr}_r\vec{v}$ for the tuple $\operatorname{pr}_{(1,\ldots,r)}\vec{v}$. We denote by $\operatorname{pr}_r\gamma$ the partition of V^r given by

$$\operatorname{pr}_r \gamma(\vec{v}) = \gamma(v_1, v_2, \dots, v_r, v_r, \dots, v_r),$$

and call it the *r-projection* of γ . Note that if γ is invariant, then for any $r \leq k$, $\operatorname{pr}_r \gamma$ is invariant.

Definition 3.2 (r-consistency). γ is said to be \vec{i} -consistent for some $\vec{i} \in [k]^r$ if for all $\vec{u}, \vec{v} \in V^k$,

$$\gamma(\vec{u}) = \gamma(\vec{v}) \Longrightarrow \mathrm{pr}_r \gamma(\mathrm{pr}_{\vec{i}}\vec{u}) = \mathrm{pr}_r \gamma(\mathrm{pr}_{\vec{i}}\vec{v}).$$

If, in addition, γ is \vec{i} -consistent for all $\vec{i} \in [k]^r$, we say that γ is r-consistent.

Observe that if for some $r \le k$, γ is r-consistent, then it is t-consistent for all $t \le r$. One may also verify that if γ is invariant, then it is k-consistent if and only if for all \vec{u} , $\vec{v} \in V^k$,

$$\gamma(\text{pr}_{k-1}\vec{u} \cdot u_{k-1}) = \gamma(\text{pr}_{k-1}\vec{v} \cdot v_{k-1}). \tag{5}$$

Definition 3.3 (graph-like partition). γ is said to be a *graph-like* partition of V^k if it is invariant, r-consistent for all $r \leq k$, and for all \vec{u} , $\vec{v} \in V^k$,

$$\gamma(\vec{u}) = \gamma(\vec{v}) \Longrightarrow (u_i = u_j \Longrightarrow v_i = v_j, \, \forall i, j \in [k]).$$
(6)

Note that if γ is graph-like, then $\operatorname{pr}_t \gamma$ is graph-like for all $t \in [k]$. An example of a graph-like partition is that of a coherent configuration, as introduced in Section 2.

Proposition 3.4. A coherent configuration is a graph-like partition.

Proof. Let ρ be a coherent configuration on V. Then ρ is a rainbow, and hence satisfies conditions (1) and (2) of Definition 2.1, from which we deduce that it is invariant and satisfies (6). To show that it is 1-consistent, let $u, v, u', v' \in V$ be such that $\rho(u, v) = \rho(u', v')$. For any $\sigma \in \text{Im } \rho$,

$$\{x \in V \mid \rho(u, x) = \sigma\} = \bigcup_{\tau \in \operatorname{Im} \rho} \{x \in V \mid \rho(u, x) = \sigma, \rho(x, v) = \tau\}.$$

⁴The concept of an invariant partition has already been introduced in Theorem 6.1 in [Evdokimov and Ponomarenko 1999].

The size of the right-hand side of the above equation is independent of the choice of (u, v) from the equivalence class $[(u, v)]_{\gamma}$. In particular, if $\sigma = \rho(u, u)$, then because $\rho(u, v) = \rho(u', v')$ there is exactly one $x \in V$ such that $\rho(u', x) = \sigma$, namely x = u'. Hence, $\rho(u, u) = \rho(u', u')$, and 1-consistency of ρ follows.

Arguments of this kind appear repeatedly in our proofs of Theorems 1.3, 1.4, and 1.5.

Another graph-like partition which will be useful throughout the paper is that of *atomic types* of k-tuples of vertices of a graph Γ , which we indicate by $\alpha_{k,\Gamma}$. To be precise, we define

$$\alpha_{k,\Gamma}: V^k \to (\operatorname{Im} \Gamma)^{[k]^{(2)}}, \quad \vec{v} \mapsto (\Gamma(v_i, v_j))_{(i,j) \in [k]^{(2)}},$$

where, as the reader may recall, we view graphs as labelled partitions.

Definition 3.5 (Γ -partition). We say that γ is a Γ -partition if $\alpha_{k,\Gamma} \leq \gamma$ for some graph Γ .

The following result is a useful property of graph-like partitions.

Lemma 3.6. For any graph-like $\gamma \in \mathcal{P}(V^k)$ and any $\vec{u}, \vec{v} \in V^k$ such that $\gamma(\vec{u}) = \gamma(\vec{v})$,

$$\gamma(\vec{u}\langle\vec{\imath}, \operatorname{pr}_{\vec{\imath}}\vec{u}\rangle) = \gamma(\vec{v}\langle\vec{\imath}, \operatorname{pr}_{\vec{\imath}}\vec{v}\rangle)$$

for all $\vec{i} \in [k]^{(r)}$ and $\vec{j} \in [k]^r$.

Proof. Let \vec{u} , \vec{v} be as in the statement. For any s, $t \in [k]$, $\gamma(\vec{u}\langle s, v_t \rangle) = \gamma(\vec{v}\langle s, u_t \rangle)$ holds true by invariance and (k-1)-consistency of γ .

Fix some r < k and set $\vec{u}' = \operatorname{pr}_{\vec{\jmath}} \vec{u}$ and $\vec{v}' = \operatorname{pr}_{\vec{\jmath}} \vec{v}$ for some $\vec{\jmath} \in [k]^r$. Suppose that $\gamma(\vec{v}\langle\vec{\imath},\vec{v}'\rangle) = \gamma(\vec{u}\langle\vec{\imath},\vec{u}'\rangle)$ holds for all $\vec{\imath} \in [k]^{(r)}$. Because r < k, there is some $i' \in [k]$ which is not an entry of $\vec{\imath}$. For some $l \in [k]$, let $x = v_l$ and $y = u_l$. Then,

$$\gamma((\vec{v}\langle\vec{i},\vec{v}'\rangle)\langle i',x\rangle) = \gamma((\vec{u}\langle\vec{i},\vec{u}'\rangle)\langle i',y\rangle)$$

by invariance and (k-1)-consistency of γ . From this, one concludes that

$$\gamma(\vec{v}\langle\vec{i}\cdot i', \vec{v}'\cdot x\rangle) = \gamma(\vec{u}\langle\vec{i}\cdot i', \vec{u}'\cdot y\rangle);$$

and hence, by setting $\vec{q} = \vec{l} \cdot l$,

$$\gamma(\vec{v}\langle\vec{i}\cdot i', \operatorname{pr}_{\vec{q}}\vec{v}\rangle) = \gamma(\vec{u}\langle\vec{i}\cdot i', \operatorname{pr}_{\vec{q}}\vec{u}\rangle).$$

We deduce the desired statement by induction.

4. Refinement operators and procedures

As previously stated, all the SPAS in this paper arise from *refinement operators*, which we now define. Fix some $k \in \mathbb{N}$.

Definition 4.1 (refinement operator). A *k*-refinement operator R is a mapping which, for each set V, assigns to each $\gamma \in \mathcal{P}(V^k)$ a partition $R \circ \gamma \in \mathcal{P}(V^k)$ such that $\gamma \leq R \circ \gamma$.

Since labelled partitions are seen as mappings, the symbol o really indicates composition thereof.

Definition 4.2 (fixed point). $\gamma \in \mathcal{P}(V^k)$ is said to be a *fixed point* of a *k*-refinement operator *R* if $\gamma \approx R \circ \gamma$. In such cases, we also say that γ is *R*-stable.

Fix some $\gamma \in \mathcal{P}(V^k)$, set $X^0 = \gamma$ and $X^i = R \circ X^{i-1}$. We then have an increasing sequence

$$X^0 \leq X^1 \leq \cdots \leq X^i \leq \cdots$$
.

Because all elements of this sequence are bounded by a labelled partition of V^k with exactly one element per equivalence class, there must be some $s \in \mathbb{N}$ such that for all $i \geq s$, X^i is fixed point of R. For the smallest such s, denote X^s by $[\gamma]^R$.

Definition 4.3 (graph-like operator). A *k*-refinement operator *R* is *graph-like* if $R \circ \gamma$ is graph-like for all graph-like $\gamma \in \mathcal{P}(V^k)$ and sets *V*.

Definition 4.4 (refinement procedure). The family of mappings $\{R_1, R_2, \dots\}$ is said to be a *refinement procedure* if, for each $k \in \mathbb{N}$:

- (1) R_k is a graph-like k-refinement operator;
- (2) if γ is a graph-like fixed point of R_k then $\operatorname{pr}_{k-1}\gamma$ is a fixed point of R_{k-1} ; and
- (3) for all sets V and $\gamma \in \mathcal{P}(V^k)$, $[\gamma]^{R_k}$ is computable in time $|V|^{O(k)}$.

For each $k, r \in \mathbb{N}$, the k-refinement operators of interest in this paper are the Weisfeiler–Leman operators $\mathrm{WL}_{k,r}$; the counting logic operators $\mathbb{C}_{k,r}$; and, for any field \mathbb{F} , the invertible map operators $\mathrm{IM}_k^{\mathbb{F}}$. For the former two, we are really interested in the case when r < k. Thus, when $k \le r$, for convenience, we let $\mathrm{WL}_{k,r} \circ \gamma = \mathbb{C}_{k,r} \circ \gamma = \gamma$ for all $\gamma \in \mathcal{P}(V^k)$ and sets V. For r < k define

$$\begin{split} \operatorname{WL}_{k,r} \circ \gamma : V^k &\to \operatorname{Im} \gamma \times \operatorname{Mult}(\Phi^{\gamma,r})^{V^r}, \qquad \vec{v} \mapsto \left(\gamma(\vec{v}), \left\{ \{(\gamma(\vec{v}\langle\vec{\imath},\vec{u}\rangle)_{\vec{\imath} \in [k]^{(r)}} \mid \vec{u} \in V^r\} \right\} \right). \\ \mathbb{C}_{k,r} \circ \gamma : V^k &\to \operatorname{Im} \gamma \times (\operatorname{Mult}(\operatorname{Im} \gamma^{V^r}))^{[k]^{(r)}}, \quad \vec{v} \mapsto \left(\gamma(\vec{v}), \left(\left\{ \{\gamma(\vec{v}\langle\vec{\imath},\vec{u}\rangle) \mid \vec{u} \in V^r\} \right\} \right)_{\vec{\imath} \in [k]^{(r)}} \right). \end{split}$$

Let $\chi_{\vec{\imath},\sigma}^{\gamma,\vec{v}}$ be the adjacency matrix of the binary relation $\{(x,y) \in V^2 \mid \gamma(\vec{v}\langle\vec{\imath},(x,y)\rangle) = \sigma\}$. Similarly to above, set $\mathrm{IM}_1^{\mathbb{F}} \circ \gamma = \mathrm{IM}_2^{\mathbb{F}} \circ \gamma = \gamma$; and, for k > 2 define

$$\operatorname{IM}_{k}^{\mathbb{F}} \circ \gamma : V^{k} \to \operatorname{Im} \gamma \times (\operatorname{Mat}_{V}(\mathbb{F})^{\operatorname{Im} \gamma \times [k]^{(2)}} / \sim), \quad \vec{v} \mapsto (\gamma(\vec{v}), ((\chi_{\vec{\iota}, \sigma}^{\gamma, \vec{v}})_{\sigma \in \operatorname{Im} \gamma})_{\vec{\iota} \in [k]^{(2)}}),$$

where \sim is the equivalence relation on elements of $\operatorname{Mat}_V(\mathbb{F})^{\operatorname{Im}\gamma\times[k]^{(2)}}$ under simultaneous similarity. That is, two tuples are equivalent if they lie in the same orbit of $GL_V(\mathbb{F})$ acting on the tuples by conjugation. Although the reader may find the above definitions rather technical, they are not crucial throughout the paper. Indeed, for k>r and $\vec{u}, \vec{v} \in V^k$, the following facts are sufficient:

- (1) $WL_{k,r} \circ \gamma(\vec{u}) = WL_{k,r} \circ \gamma(\vec{v})$ if and only if $\gamma(\vec{u}) = \gamma(\vec{v})$, and for all $\vec{\phi} \in \Phi^{\gamma,r}$ and $\vec{\iota} \in [k]^{(r)}$
 - $\left| \left\{ \vec{x} \in V^r \mid \gamma(\vec{u}\langle \vec{\imath}, \vec{x} \rangle) = \phi_{\vec{\imath}} \right\} \right| = \left| \left\{ \vec{x} \in V^r \mid \gamma(\vec{v}\langle \vec{\imath}, \vec{x} \rangle) = \phi_{\vec{\imath}} \right\} \right|.$
- (2) $\mathbb{C}_{k,r} \circ \gamma(\vec{u}) = \mathbb{C}_{k,r} \circ \gamma(\vec{v})$ if $\gamma(\vec{u}) = \gamma(\vec{v})$, and for all $\sigma \in \text{Im } \gamma$ and $\vec{\iota} \in [k]^{(r)}$

$$\left| \{ \vec{x} \in V^r \mid \gamma(\vec{u}\langle\vec{\imath}, \vec{x}\rangle) = \sigma \} \right| = \left| \{ \vec{x} \in V^r \mid \gamma(\vec{v}\langle\vec{\imath}, \vec{x}\rangle) = \sigma \} \right|.$$

(3) $\operatorname{IM}_k^{\mathbb{F}} \circ \gamma(\vec{u}) = \operatorname{IM}_k^{\mathbb{F}} \circ \gamma(\vec{v})$ if and only if $\gamma(\vec{u}) = \gamma(\vec{v})$, and for each $\vec{i} \in [k]^{(2)}$ there exists some $S \in \operatorname{GL}_V(\mathbb{F})$ such that for all $\sigma \in \operatorname{Im} \gamma$,

$$S\chi_{\vec{i},\sigma}^{\gamma,\vec{u}}S^{-1}=\chi_{\vec{i},\sigma}^{\gamma,\vec{v}}.$$

From this, one may derive the following *stability conditions*:

Proposition 4.5. For any $\gamma \in \mathcal{P}(V^k)$ and k > r:

- (1) γ is $\operatorname{WL}_{k,r}$ -stable if and only if for all $\vec{v} \in V^k$, $\vec{\iota} \in [k]^{(r)}$, and $\vec{\phi} \in \Phi^{\gamma,r}$, the size of the set $\{\vec{x} \in V^r \mid \gamma(\vec{v}\langle\vec{\iota},\vec{x}\rangle) = \phi_{\vec{\iota}}\}$ is independent of the choice of \vec{v} from the equivalence class $[\vec{v}]_{\gamma}$.
- (2) γ is $\mathbb{C}_{k,r}$ -stable if and only if for all $\vec{v} \in V^k$, $\vec{i} \in [k]^{(r)}$, and $\sigma \in \text{Im } \gamma$, the size of the set $\{\vec{x} \in V^r \mid \gamma(\vec{v} \setminus \vec{i}, \vec{x}) = \sigma\}$ is independent of the choice of \vec{v} from the equivalence class $[\vec{v}]_{\gamma}$.
- (3) γ is $\operatorname{IM}_{k}^{\mathbb{F}}$ -stable if for all $\vec{u}, \vec{v} \in V^{k}$ and $\vec{i} \in [k]^{(2)}$,

$$\gamma(\vec{u}) = \gamma(\vec{v}) \Longrightarrow \exists S \in \mathrm{GL}_V(\mathbb{F}), \, S\chi_{\vec{\imath},\sigma}^{\gamma,\vec{u}}S^{-1} = \chi_{\vec{\imath}\,\sigma}^{\gamma,\vec{v}}, \, \forall \sigma \in \mathrm{Im}\,\gamma.$$

Consider the following families of mappings:

- (1) For all $r \in \mathbb{N}$, $\mathbf{WL}_r = \{WL_{1,r}, WL_{2,r}, \dots\}$.
- (2) For all $r \in \mathbb{N}$, $\mathbf{C}_r = \{\mathbb{C}_{1,r}, \mathbb{C}_{2,r}, \dots\}$.
- (3) For any field \mathbb{F} , $\mathbf{IM}(\mathbb{F}) = \{ \mathrm{IM}_1^{\mathbb{F}}, \mathrm{IM}_2^{\mathbb{F}}, \dots \}$.

Proposition 4.6. The families \mathbf{WL}_r , \mathbf{C}_r , and $\mathbf{IM}(\mathbb{F})$ are refinement procedures for all $r \in \mathbb{N}$ and fields \mathbb{F} .

For the proof of Proposition 4.6 and that of the next auxiliary lemma, we use the following notations and conventions. Fix a graph-like partition $\gamma \in \mathcal{P}(V^k)$. Let $\bar{\gamma} = \operatorname{pr}_{k-1}\gamma$, and for all $\vec{v} \in V^k$, if $\sigma = \gamma(\vec{v})$, let $\bar{\sigma} = \bar{\gamma}(\operatorname{pr}_{k-1}\vec{v})$. For $\pi \in \operatorname{Sym}(k)$, we let $\sigma^{\pi} = \gamma(\vec{v}^{\pi})$ and define an action on $[k]^{(r)}$ by setting $(\pi(\vec{\iota}))_j = \pi(i_j)$ for all $j \in [r]$. For any $\vec{v} \in V^k$, we denote $\vec{v}' = \operatorname{pr}_{k-1}\vec{v} \cdot v_{k-1}$. Note that $\bar{\sigma}$ and σ^{π} are well defined, since γ is graph-like.

Lemma 4.7. *Let* $\vec{i} \in [k]^{(r)}$.

- (1) If k is an entry of \vec{i} , then $\gamma(\vec{v}'\langle\vec{i},\vec{x}\rangle) = \gamma(\vec{v}\langle\vec{i},\vec{x}\rangle)$.
- (2) If k-1 is an entry of \vec{i} , but k is not; then $\gamma(\vec{v}^{\pi}\langle\vec{i},\vec{x}\rangle) = \gamma(\vec{v}'\langle\vec{i},\vec{x}\rangle)$, where $\pi = (k-1, k)$ is a transposition of Sym(k).
- (3) If neither k nor k-1 are entries of $\vec{\imath}$, then $\gamma(\vec{v}\langle\vec{\imath},\vec{x}\rangle) = \sigma \Longrightarrow \gamma(\vec{v}'\langle\vec{\imath},\vec{x}\rangle) = \bar{\sigma}$. In particular,

$$\{\vec{x} \in V^r \mid \gamma(\vec{v}'\langle \vec{i}, \vec{x} \rangle) = \kappa\} = \bigcup_{\substack{\sigma \in \text{Im } \gamma \\ \bar{\sigma} = \kappa}} \{\vec{x} \in V^r \mid \gamma(\vec{v}\langle \vec{i}, \vec{x} \rangle) = \sigma\}. \tag{7}$$

Proof. Statements (1) and (2) are trivial to check.

Let $\vec{w} = \vec{v} \langle \vec{\imath}, \vec{x} \rangle$ with $\gamma(\vec{w}) = \sigma$. Then $\vec{w}' = \vec{v}' \langle \vec{\imath}, \vec{x} \rangle$, so $\gamma(\vec{w}') = \bar{\sigma}$ by definition. From this, and the fact that γ is (k-1)-consistent, the right-hand side of (7) is a subset of the left-hand side. The reverse inclusion follows from the definition of $\bar{\sigma}$ in terms of σ , and statement (3) follows.

Proof of Proposition 4.6. We check that each of the families of mappings satisfy the conditions in Definition 4.4. Note that $\gamma \leq R \circ \gamma$ for any k-refinement operator R. Hence, since γ is graph-like, for any $\vec{u}, \vec{v} \in V^k$, $R \circ \gamma(\vec{u}) = R \circ \gamma(\vec{v})$ implies that $u_i = u_j \Longrightarrow v_i = v_j$ for all $i, j \in [k]$. Thus, to show that $R \circ \gamma$ is graph-like, it suffices to verify that $R \circ \gamma$ is invariant and satisfies (5). That is, for all $\vec{u}, \vec{v} \in V^k$, $R \circ \gamma(\vec{u}) = R \circ \gamma(\vec{v}) \Longrightarrow R \circ \gamma(\vec{v}') = R \circ \gamma(\vec{v}')$.

 WL_r is a refinement procedure. We first show that $WL_{k,r} \circ \gamma$ is invariant and satisfies (5), and is thus a graph-like partition.

Suppose $\operatorname{WL}_{k,r} \circ \gamma(\vec{u}) = \operatorname{WL}_{k,r} \circ \gamma(\vec{v})$. Then $\gamma(\vec{u}) = \gamma(\vec{v})$ by definition, and hence $\gamma(\vec{u}^{\tau}) = \gamma(\vec{v}^{\tau})$ since γ is graph-like and thus invariant. Furthermore, from the invariance of γ , it follows that, for all $\tau \in \operatorname{Sym}(k)$ and $\vec{\phi} \in \Phi^{\gamma,r}$,

$$\{\vec{x} \in V^r \mid \gamma(\vec{v}^\tau \langle \tau(\vec{i}), \vec{x} \rangle) = (\phi_{\vec{i}})^\tau, \ \forall \vec{i} \in [k]^{(r)}\} = \{\vec{x} \in V^r \mid \gamma(\vec{v} \langle \vec{i}, \vec{x} \rangle) = \phi_{\vec{i}}, \ \forall \vec{i} \in [k]^{(r)}\}.$$

Hence, $WL_{k,r} \circ \gamma$ is invariant.

For all $\vec{\phi} \in \Phi^{\gamma,r}$, let $\vec{\phi}^{\dagger} \in \Phi^{\gamma,r}$ be defined as

$$\phi_{\vec{i}}^{\dagger} = \begin{cases} \bar{\phi}_{\vec{i}} & \text{if neither } k \text{ nor } k - 1 \text{ are entries of } \vec{i}, \\ (\phi_{\vec{i}})^{\pi} & \text{if } k - 1 \text{ is an entry of } \vec{i}, \\ \phi_{\vec{i}} & \text{otherwise,} \end{cases}$$

where $\pi = (k-1,k) \in \operatorname{Sym}(k)$. Observe that if $\{\vec{x} \in V^r \mid \gamma(\vec{v}'\langle \vec{\imath}, \vec{x}\rangle) = \psi_{\vec{\imath}}, \ \forall \vec{\imath} \in [k]^{(r)}\}$ is nonempty for some $\vec{v} \in V^k$ and $\vec{\psi} \in \Phi^{\gamma,r}$, then $\vec{\psi} = \vec{\phi}^{\dagger}$ for some $\vec{\phi} \in \Phi^{\gamma,r}$. It follows from the definition of $\vec{\phi}^{\dagger}$ and Lemma 4.7 that

$$\{\vec{x} \in V^r \mid \gamma(\vec{v}'\langle \vec{i}, \vec{x}\rangle) = \psi_{\vec{i}}, \ \forall \vec{i} \in [k]^{(r)}\} = \bigcup_{\substack{\vec{\phi} \in \Phi^{\gamma, r} \\ \vec{\psi} = \vec{\phi}^{\dagger}}} \{\vec{x} \in V^r \mid \gamma(\vec{v}\langle \vec{i}, \vec{x}\rangle) = \phi_{\vec{i}}, \ \forall \vec{i} \in [k]^{(r)}\}.$$
(8)

Since $WL_{k,r} \circ \gamma(\vec{u}) = WL_{k,r} \circ \gamma(\vec{v})$ for all $\vec{\phi} \in \Phi^{\gamma,r}$,

$$\left| \{ \vec{x} \in V^r \mid \gamma(\vec{u}(\vec{i}, \vec{x})) = \phi_{\vec{i}}, \ \forall \vec{i} \in [k]^{(r)} \} \right| = \left| \{ \vec{x} \in V^r \mid \gamma(\vec{v}(\vec{i}, \vec{x})) = \phi_{\vec{i}}, \ \forall \vec{i} \in [k]^{(r)} \} \right|.$$

Thus, because the right-hand side of (8) is a disjoint union, we deduce that

$$\left| \{ \vec{x} \in V^r \mid \gamma(\vec{u}'\langle \vec{\imath}, \vec{x} \rangle) = \phi_{\vec{\imath}}, \ \forall \vec{\imath} \in [k]^{(r)} \} \right| = \left| \{ \vec{x} \in V^r \mid \gamma(\vec{v}'\langle \vec{\imath}, \vec{x} \rangle) = \phi_{\vec{\imath}}, \ \forall \vec{\imath} \in [k]^{(r)} \} \right|,$$

for all $\vec{\phi} \in \Phi^{\gamma,r}$. This implies $\mathrm{WL}_{k,r} \circ \gamma(\vec{u}') = \mathrm{WL}_{k,r} \circ \gamma(\vec{v}')$; and hence, that $\mathrm{WL}_{k,r} \circ \gamma$ is graph-like. Suppose now that γ is $\mathrm{WL}_{k,r}$ -stable. Observe that for all $\vec{\phi} \in \Phi^{\bar{\gamma},r}$,

$$\{\vec{x} \in V^r \mid \bar{\gamma}(\operatorname{pr}_{k-1}\vec{v}\langle\vec{i}, \vec{x}\rangle) = \phi_{\vec{i}}, \ \forall \vec{i} \in [k-1]^{(r)}\} = \bigcup_{\substack{\vec{\psi} \in \Phi^{\gamma, r} \\ \bar{\psi}_{\vec{i}} = \phi_{\vec{i}}, \ \forall \vec{i} \in [k-1]^{(r)}}} \{\vec{x} \in V^r \mid \gamma(\vec{v}\langle\vec{i}, \vec{x}\rangle) = \psi_{\vec{i}}, \ \forall \vec{i} \in [k]^{(r)}\}.$$
(9)

From the $\mathrm{WL}_{k,r}$ -stability of γ and the fact that $\gamma(\vec{u}) = \gamma(\vec{v})$, it follows that for all $\vec{\psi} \in \Phi^{\gamma,r}$,

$$|\{\vec{x} \in V^r \mid \gamma(\vec{u}\langle\vec{\imath}, \vec{x}\rangle) = \psi_{\vec{\imath}}, \ \forall \vec{\imath} \in [k]^{(r)}\}| = |\{\vec{x} \in V^r \mid \gamma(\vec{v}\langle\vec{\imath}, \vec{x}\rangle) = \psi_{\vec{\imath}}, \ \forall \vec{\imath} \in [k]^{(r)}\}|.$$

Since the right-hand side of (9) is a disjoint union, it holds that for all $\vec{\phi} \in \Phi^{\bar{\gamma},r}$,

$$\left|\{\vec{x}\in V^r\mid \overline{\gamma}(\mathrm{pr}_{k-1}\vec{u}\langle\vec{\imath},\vec{x}\rangle)=\phi_{\vec{\imath}},\;\forall\,\vec{\imath}\in [k-1]^{(r)}\}\right|=\left|\{\vec{x}\in V^r\mid \overline{\gamma}(\mathrm{pr}_{k-1}\vec{v}\langle\vec{\imath},\vec{x}\rangle)=\phi_{\vec{\imath}},\;\forall\,\vec{\imath}\in [k-1]^{(r)}\}\right|.$$

Thus, $\bar{\gamma}$ is WL_{k-1,r}-stable.

To see that the stable partition $[\gamma]^{\operatorname{WL}_{k,r}}$ can be computed in time $|V|^{O(k)}$, note that the number of iterations of the refinement operator before the fixed point is reached is at most $|V|^k$. At each step, for each k-tuple \vec{v} , we need to compute the colour $\operatorname{WL}_{k,r} \circ \gamma(\vec{v})$, which involves checking the colour of k|V| distinct tuples. The total time required at each step is therefore $O(k|V|^{k+1})$. Repeating this for $|V|^k$ steps gives us the required bound.

 C_r is a refinement procedure. The proof, in this case, is similar to that of $WL_{k,r}$.

Suppose $\mathbb{C}_{k,r} \circ \gamma(\vec{u}) = \mathbb{C}_{k,r} \circ \gamma(\vec{v})$. The invariance of γ implies that $\gamma(\vec{u}^{\tau}) = \gamma(\vec{v}^{\tau})$ and that

$$\{\vec{x} \in V^r \mid \gamma(\vec{v}^\tau \langle \tau(\vec{i}), \vec{x} \rangle) = \sigma^\tau\} = \{\vec{x} \in V^r \mid \gamma(\vec{v} \langle \vec{i}, \vec{x} \rangle) = \sigma\},\$$

for all $\tau \in \operatorname{Sym}(k)$ and $\vec{i} \in [k]^{(r)}$. Hence,

$$\left| \{ \vec{x} \in V^r \mid \gamma(\vec{u}^\tau \langle \tau(\vec{i}), \vec{x} \rangle) = \sigma^\tau \} \right| = \left| \{ \vec{x} \in V^r \mid \gamma(\vec{v}^\tau \langle \tau(\vec{i}), \vec{x} \rangle) = \sigma^\tau \} \right|.$$

Therefore, $\mathbb{C}_{k,r} \circ \gamma(\vec{u}^{\tau}) = \mathbb{C}_{k,r} \circ \gamma(\vec{v}^{\tau})$; thus showing that $\mathbb{C}_{k,r} \circ \gamma$ is invariant.

As γ is graph-like and $\mathbb{C}_{k,r} \circ \gamma(\vec{u}) = \mathbb{C}_{k,r} \circ \gamma(\vec{v})$, it follows from (1) in Lemma 4.7 that for all $\vec{i} \in [k]^{(r)}$ with some entry equal to k,

$$\left| \{ \vec{x} \in V^r \mid \gamma(\vec{u}'\langle \vec{\imath}, \vec{x} \rangle) = \sigma \} \right| = \left| \{ \vec{x} \in V^r \mid \gamma(\vec{v}'\langle \vec{\imath}, \vec{x} \rangle) = \sigma \} \right|. \tag{10}$$

If $\vec{i} \in [k]^{(r)}$ has an entry equal k-1 but none equal to k, then (2) in Lemma 4.7 implies that

$$\left| \{ \vec{x} \in V^r \mid \gamma(\vec{v}'\langle \vec{i}, \vec{x} \rangle) = \sigma \} \right| = \left| \{ \vec{x} \in V^r \mid \gamma(\vec{v}^\pi \langle \pi(\vec{i}), \vec{x} \rangle) = \sigma \} \right|.$$

From the invariance of $\mathbb{C}_{k,r} \circ \gamma$, we deduce that (10) holds also for such values of \vec{i} . Finally, if no entry of \vec{i} is equal to k or k-1, γ satisfies (7). As the right-hand side of the latter is a disjoint union and for all $\sigma \in \text{Im } \gamma$

$$\left| \{ \vec{x} \in V^r \mid \gamma(\vec{u}\langle\vec{i}, \vec{x}\rangle) = \sigma \} \right| = \left| \{ \vec{x} \in V^r \mid \gamma(\vec{v}\langle\vec{i}, \vec{x}\rangle) = \sigma \} \right|,$$

it follows that (10) holds also when neither k nor k-1 are entries of $\vec{\iota}$. Thus, $C_{k,r} \circ \gamma(\vec{u}') = \mathbb{C}_{k,r} \circ \gamma(\vec{v}')$ and $\mathbb{C}_{k,r} \circ \gamma$ is therefore graph-like.

Suppose γ is $\mathbb{C}_{k,r}$ -stable. Since γ is graph-like, it holds that

$$\{\vec{x} \in V^r \mid \overline{\gamma}(\operatorname{pr}_{k-1}\vec{v}\langle\vec{t}, \vec{x}\rangle) = \kappa\} = \bigcup_{\substack{\sigma \in \operatorname{Im} \gamma \\ \overline{\sigma} = \kappa}} \{\vec{x} \in V^r \mid \gamma(\vec{v}\langle\vec{t}, \vec{x}\rangle) = \sigma\},\tag{11}$$

for all $\vec{i} \in [k-1]^{(r)}$ and $\kappa \in \text{Im } \bar{\gamma}$. Since the right-hand side of the above is a disjoint union and

$$\left| \{ \vec{x} \in V^r \mid \gamma(\vec{u}\langle\vec{i}, \vec{x}\rangle) = \sigma \} \right| = \left| \{ \vec{x} \in V^r \mid \gamma(\vec{v}\langle\vec{i}, \vec{x}\rangle) = \sigma \} \right|,$$

it then follows from (11) that

$$\left| \{ \vec{x} \in V^r \mid \overline{\gamma} (\operatorname{pr}_{k-1} \vec{u} \langle \vec{\imath}, \vec{x} \rangle) = \sigma \} \right| = \left| \{ \vec{x} \in V^r \mid \overline{\gamma} (\operatorname{pr}_{k-1} \vec{v} \langle \vec{\imath}, \vec{x} \rangle) = \sigma \} \right|.$$

Thus, $\bar{\gamma}$ is $\mathbb{C}_{k-1,r}$ -stable.

A very similar argument to the case of $WL_{k,r}$ shows that $[\gamma]^{\mathbb{C}_{k,r}}$ can be computed in time $|V|^{O(k)}$.

 $IM(\mathbb{F})$ is a refinement procedure. We proceed via the same proof strategy as for the above refinement procedures.

Suppose $\operatorname{IM}_k^{\mathbb{F}} \circ \gamma(\vec{u}) = \operatorname{IM}_k^{\mathbb{F}} \circ \gamma(\vec{v})$. Then $\gamma(\vec{u}) = \gamma(\vec{v})$; and hence, $\gamma(\vec{u}^{\tau}) = \gamma(\vec{v}^{\tau})$ for all $\tau \in \operatorname{Sym}(k)$. Note that for any $\sigma \in \operatorname{Im} \gamma$, $\chi_{\vec{t},\sigma}^{\gamma,\vec{v}} = \chi_{\tau(\vec{t})}^{\gamma,\vec{v}^{\tau}}$, σ^{τ} . From this one deduces that $\operatorname{IM}_k^{\mathbb{F}} \circ \gamma(\vec{u}^{\tau}) = \operatorname{IM}_k^{\mathbb{F}} \circ \gamma(\vec{v}^{\tau})$; whence, $\operatorname{IM}_k^{\mathbb{F}} \circ \gamma$ is invariant.

From Lemma 4.7, it follows that $\chi_{\vec{t},\sigma}^{\gamma,\vec{v}'} = \chi_{\vec{t},\sigma}^{\gamma,\vec{v}}$ if k is an entry of \vec{t} ; $\chi_{\vec{t},\sigma}^{\gamma,\vec{v}'} = \chi_{\pi(\vec{t}),\sigma}^{\gamma,\vec{v}^{\pi}}$ if k-1 is an entry of \vec{t} , but k is not; and

$$\chi_{\vec{l},\sigma}^{\gamma,\vec{v}'} = \sum_{\substack{\kappa \in \operatorname{Im} \gamma \\ \vec{\kappa} = \sigma}} \chi_{\vec{l},\kappa}^{\gamma,\vec{v}}$$

if neither k nor k-1 are entries of $\vec{\iota}$. Since $S\chi_{\vec{\iota},\sigma}^{\gamma,\vec{v}}S^{-1}=\chi_{\vec{\iota},\sigma}^{\gamma,\vec{u}}$ for all $\sigma\in \operatorname{Im}\gamma$, then $S\chi_{\vec{\iota},\sigma}^{\gamma,\vec{v}'}S^{-1}=\chi_{\vec{\iota},\sigma}^{\gamma,\vec{u}'}$ for all $\sigma\in \operatorname{Im}\gamma$. From this, $\operatorname{IM}_k^{\mathbb{F}}\circ\gamma(\vec{u}')=\operatorname{IM}_k^{\mathbb{F}}\circ\gamma(\vec{v}')$ follows; and thus, $\operatorname{IM}_k^{\mathbb{F}}\circ\gamma$ is graph-like.

From the fact that γ is graph-like we deduce that for all $\vec{i} \in [k-1]^{(2)}$,

$$\chi_{\vec{\imath},\sigma}^{\vec{\gamma},\operatorname{pr}_{k-1}\vec{v}} = \sum_{\substack{\kappa \in \operatorname{Im} \gamma \\ \vec{v} = \sigma}} \chi_{\vec{\imath},\kappa}^{\gamma,\vec{v}}.$$

Thus, if $S\chi_{\bar{t},\sigma}^{\gamma,\vec{v}}S^{-1}=\chi_{\bar{t},\sigma}^{\gamma,\vec{u}}$ for all $\sigma\in\operatorname{Im}\gamma$, then $S\chi_{\bar{t},\sigma}^{\bar{\gamma},\operatorname{pr}_{k-1}\vec{v}}S^{-1}=\chi_{\bar{t},\sigma}^{\bar{\gamma},\operatorname{pr}_{k-1}\vec{u}}$ for all $\sigma\in\operatorname{Im}\gamma$. In particular, if γ is $\operatorname{IM}_k^{\mathbb{F}}$ -stable, then $\bar{\gamma}$ is $\operatorname{IM}_{k-1}^{\mathbb{F}}$ -stable.

Finally, showing that $[\gamma]^{\text{IM}_k^F}$ can be computed in time $|V|^{O(k)}$ is similar to the previous cases. Again, the number of refinement steps is at most $|V|^k$. However, at each step, and for each pair of tuples, we need to perform a simultaneous similarity test. For this, we rely on the fact that simultaneous similarity is decidable in polynomial time. This follows from the fact that testing simultaneous similarity can be reduced in polynomial time to testing module isomorphism (see [Chistov et al. 1997], for example), and the polynomial-time algorithm for the latter problem over any field is given by Brookbanks and Luks in [2008].

For a refinement procedure $\mathbf{R} = \{R_1, R_2, \dots\}$, let $\mathcal{S}_{\mathbf{R}} = \{\overline{R}_1, \overline{R}_2, \dots\}$ be the family of mappings where for all graphs Γ , $\overline{R}_1(\Gamma) = \Gamma$ and $\overline{R}_k(\Gamma) = \mathrm{pr}_2[\alpha_{k,\Gamma}]^{R_k}$ for $k \geq 2$. Then, from the above discussion, $\mathcal{S}_{\mathbf{R}}$ is a SPAS. We define $\mathcal{S}_{\mathbf{WL},r} = \{\overline{\mathbf{WL}}_{1,r}, \overline{\mathbf{WL}}_{2,r}, \dots\}$, $\mathcal{S}_{\mathbf{C},r} = \{\overline{\mathbb{C}}_{1,r}, \overline{\mathbb{C}}_{2,r}, \dots\}$, and $\mathcal{S}_{\mathbf{IM}(\mathbb{F})} = \{\overline{\mathbf{IM}}_1^{\mathbb{F}}, \overline{\mathbf{IM}}_2^{\mathbb{F}}, \dots\}$ to be the SPASs obtained from the refinement procedures \mathbf{WL}_r , \mathbf{C}_r , and $\mathbf{IM}(\mathbb{F})$, respectively.

5. Proof of Theorem 1.3

Hereafter, when talking of the families of operators \mathbf{WL}_r and \mathbf{C}_r , we write $\mathbf{WL} = \{\mathbf{WL}_1, \mathbf{WL}_2, \ldots\}$ and $\mathbf{C} = \{\mathbb{C}_1, \mathbb{C}_2, \ldots\}$ to denote the families $\mathbf{WL}_1 = \{\mathbf{WL}_{1,1}, \mathbf{WL}_{2,1}, \ldots\}$ and $\mathbf{C}_1 = \{\mathbb{C}_{1,1}, \mathbb{C}_{2,1}, \ldots\}$, respectively. The distinction between the *procedure* \mathbf{WL}_k and the *operator* \mathbf{WL}_k should not cause any confusion, likewise for the procedure \mathbf{C}_k and the operator \mathbb{C}_k .

For showing that for two refinement procedures $\mathbf{X} = \{X_1, X_2, \ldots\}$ and $\mathbf{Y} = \{Y_1, Y_2, \ldots\}$, their respective SPAS satisfy $\mathcal{S}_{\mathbf{X}} \succeq \mathcal{S}_{\mathbf{Y}}$, our general strategy is as follows. For each $k \in \mathbb{N}$, we find some $k' \in \mathbb{N}$ such that for any graph Γ and any Y_k -stable graph-like Γ -partition $\gamma \in \mathcal{P}(V^k)$, there is some $X_{k'}$ -stable graph-like Γ -partition $\gamma' \in \mathcal{P}(V^{k'})$ such that $\operatorname{pr}_2 \gamma' \succeq \operatorname{pr}_2 \gamma$. By taking γ' to be $[\alpha_{k,\Gamma}]^{X_{k'}}$ (which is graph-like by Proposition 4.6), we have that $\overline{X}_k(\Gamma) = \operatorname{pr}_2 \gamma' \succeq \operatorname{pr}_2 \gamma \succeq \overline{Y}_k(\Gamma)$ for all graphs Γ . The last inequality follows from the fact that $\overline{Y}_k(\Gamma)$ is the 2-projection of a minimal Y_k -stable partition refining $\alpha_{k,\Gamma}$.

For the proofs in this section, the operators $WL_{k,r}$ and $\mathbb{C}_{k,r}$ are only considered in the case k > r, since the statements hold trivially when $k \le r$.

Lemma 5.1. For all $k, r \in \mathbb{N}$, any $WL_{k,r}$ -stable partition is $\mathbb{C}_{k,r}$ -stable.

Proof. Let $\gamma \in \mathcal{P}(V^k)$ be $\mathrm{WL}_{k,r}$ -stable. The size of $\{\vec{x} \in V^r \mid \gamma(\vec{v}\langle\vec{i},\vec{x}\rangle) = \phi_{\vec{i}}, \ \forall \vec{i} \in [k]^{(r)}\}$ is then independent of the choice of \vec{v} from the equivalence class $[\vec{v}]_{\gamma}$ for all $\vec{\phi} \in \Phi^{\gamma,r}$, by Proposition 4.5. For each $\vec{j} \in [k]^{(r)}$ and $\sigma \in \mathrm{Im} \ \gamma$,

$$\{\vec{x} \in V^r \mid \gamma(\vec{v}\langle \vec{J}, \vec{x}\rangle) = \sigma\} = \bigcup_{\substack{\vec{\phi} \in \Phi^{\gamma, r} \\ \phi_{\vec{\imath}} = \sigma}} \{\vec{x} \in V^r \mid \gamma(\vec{v}\langle \vec{\imath}, \vec{x}\rangle) = \phi_{\vec{\imath}}, \ \forall \vec{\imath} \in [k]^{(r)}\}.$$

The right-hand side of the above is a disjoint union of sets, whose sizes are independent of the choice of \vec{v} , from the equivalence class $[\vec{v}]_{\gamma}$. Hence, the size of the left-hand side, is also the same for any choice of \vec{v} from the class $[\vec{v}]_{\gamma}$, implying that γ is $\mathbb{C}_{k,r}$ -stable.

Corollary 5.2. For any graph Γ and $k, r \in \mathbb{N}$,

$$\overline{\mathbb{C}}_{k,r}(\Gamma) \leq \overline{\mathrm{WL}}_{k,r}(\Gamma).$$

Lemma 5.3. For all $k, r \in \mathbb{N}$, any \mathbb{C}_k -stable partition is $\mathbb{C}_{k,r}$ -stable.

Proof. Let $\gamma \in \mathcal{P}(V^k)$ be \mathbb{C}_k -stable. Then, by definition, it is $\mathbb{C}_{k,1}$ stable.

Suppose γ is $\mathbb{C}_{k,m}$ -stable for some m < k. Let $\vec{j} \in [k]^{(m)}$. Since m < k, there is some $j' \in [k]$ such that $j' \neq j_i$ for all $i \in [m]$. For all $\sigma, \tau \in \text{Im } \gamma$ and $\vec{v} \in V^k$ such that $\gamma(\vec{v}) = \sigma$, define

$$p_{\sigma\tau} = \left| \{ u \in V \mid \gamma(\vec{v}\langle j', u \rangle) = \tau \} \right|$$

and

$$q_{\sigma\tau} = \left| \{ \vec{w} \in V^m \mid \gamma(\vec{v}\langle \vec{J}, \vec{w} \rangle) = \tau \} \right|.$$

Because γ is \mathbb{C}_k -stable and, by the induction hypothesis, also $\mathbb{C}_{k,m}$ -stable, $p_{\sigma\tau}$ and $q_{\sigma\tau}$ are independent of the choice of \vec{v} from the equivalence class $[\vec{v}]_{\gamma}$.

Observe that

$$\{\vec{w} \cdot u \in V^{m+1} \mid \gamma(\vec{v}\langle \vec{j} \cdot j', \vec{w} \cdot u \rangle) = \tau\} = \{\vec{w} \cdot u \in V^{m+1} \mid \gamma((\vec{v}\langle \vec{j}, \vec{w} \rangle)\langle j', u \rangle) = \tau\}.$$

From this one deduces that

$$\left| \{ \vec{w} \cdot u \in V^{m+1} \mid \gamma(\vec{v} \langle \vec{j} \cdot j', \vec{w} \cdot u \rangle) = \tau \} \right| = \sum_{\alpha \in \operatorname{Im} \gamma} p_{\sigma\alpha} q_{\alpha\tau}.$$

The right-hand side of the above is independent of the choice of \vec{v} from the equivalence class $[\vec{v}]_{\gamma}$. Hence, γ is $\mathbb{C}_{k,m+1}$ -stable. Thus, by induction, it is $\mathbb{C}_{k,r}$ stable for all r.

Corollary 5.4. For any graph Γ and $k, r \in \mathbb{N}$,

$$\overline{\mathbb{C}}_k(\Gamma) \succeq \overline{\mathbb{C}}_{k,r}(\Gamma).$$

Note that Lemmas 5.1 and 5.3 are not restricted to graph-like partitions.

Lemma 5.5. For all $k, r \in \mathbb{N}$, the k-projection of a $\mathbb{C}_{k+r,r}$ -stable graph-like partition is $\mathrm{WL}_{k,r}$ -stable.

Proof. Let \vec{v} , $\vec{v}' \in V^k$, and let $\bar{\gamma} = \operatorname{pr}_k \gamma$ for some $\mathbb{C}_{k+r,r}$ -stable graph-like partition $\gamma \in \mathcal{P}(V^{k+r})$. Because γ is graph-like, by Lemma 3.6 it follows that for all $\vec{i} \in [k+r]^{(r)}$,

$$\gamma(\vec{v}\langle\vec{i}, \operatorname{pr}_{\vec{i}}\vec{v}\rangle) = \gamma(\vec{v}'\langle\vec{i}, \operatorname{pr}_{\vec{i}}\vec{v}'\rangle),$$

where $\vec{j} = (k+1, k+2, \dots, k+r) \in [k+r]^{(r)}$. In particular, for all $\vec{i} \in [k]^{(r)}$,

$$\bar{\gamma}(\operatorname{pr}_{k}\vec{v}\langle\vec{i},\operatorname{pr}_{\vec{i}}\vec{v}\rangle) = \bar{\gamma}(\operatorname{pr}_{k}\vec{v}'\langle\vec{i},\operatorname{pr}_{\vec{i}}\vec{v}'\rangle). \tag{12}$$

For all $\vec{v} \in V^{k+r}$, define $\vec{\Delta} \in (\Phi^{\bar{\gamma},r})^{\operatorname{Im}\gamma}$ to be

$$(\Delta_{\gamma(\vec{v})})_{\vec{i}} = \bar{\gamma}(\operatorname{pr}_{k}\vec{v}\langle\vec{i},\operatorname{pr}_{\vec{i}}\vec{v}\rangle).$$

One deduces from (12) that $\vec{\Delta}$ is well defined.

For any $\vec{\phi} \in \Phi^{\bar{\gamma},r}$, it follows from the definition of $\vec{\Delta}$ that:

$$\{\vec{u} \in V^r \mid \overline{\gamma}(\operatorname{pr}_k \vec{v} \langle \vec{\iota}, \vec{u} \rangle) = \phi_{\vec{\iota}}, \ \forall \vec{\iota} \in [k]^{(r)}\} = \bigcup_{\substack{\sigma \in \operatorname{Im} \gamma \\ \Delta_{\sigma} = \vec{\phi}}} \{\vec{u} \in V^r \mid \gamma(\vec{v} \langle \vec{\jmath}, \vec{u} \rangle) = \sigma\}.$$

 $\mathbb{C}_{k,r}$ -stability of γ implies that the size of the right-hand side of the above is independent of the choice of \vec{v} from the equivalence class $[\vec{v}]_{\gamma}$. Since γ is graph-like, the size of the left-hand side is independent of the choice of $\vec{w} \in V^{k+r}$, where $\operatorname{pr}_k \vec{w} \in [\operatorname{pr}_k \vec{v}]_{\bar{\gamma}}$, and the result follows.

Corollary 5.6. For any graph Γ and $k, r \in \mathbb{N}$,

$$\overline{\mathbb{C}}_{k+r,r}(\Gamma) \succeq \overline{\mathrm{WL}}_{k,r}(\Gamma).$$

In particular,

$$\overline{\mathbb{C}}_{k+1}(\Gamma) \succ \overline{\mathrm{WL}}_k(\Gamma).$$

Lemma 5.7. For all $k, r \in \mathbb{N}$, the (k-r+1)-projection of a $\mathbb{C}_{k,r}$ -stable graph-like partition is \mathbb{C}_{k-r+1} -stable.

Proof. Let $\bar{\gamma} = \operatorname{pr}_{k-r+1} \gamma$, and let $\vec{v} \in V^k$. Fix some $i \in [k-r]$, and let $\vec{j} \in [k]^{(k-r+1)}$ be defined by $\vec{j} = (i, k-r+1, k-r+2, \ldots, k-1, k)$. Fix $u \in \vec{V}$, and let $\vec{w} = (v_{k-r+1}, v_{k-r+1}, \ldots, v_{k-r+1}) \in V^{k-r}$. Since γ is graph-like, for any $u' \in V$,

$$\gamma(\vec{v}\langle\vec{\jmath},u\cdot\vec{w}\rangle) = \gamma(\vec{v}\langle\vec{\jmath},u'\cdot\vec{w}\rangle) \Longleftrightarrow \bar{\gamma}(\mathrm{pr}_k\vec{v}\langle i,u\rangle) = \bar{\gamma}(\mathrm{pr}_k\vec{v}\langle i,u'\rangle);$$

and therefore,

$$\{u' \in V \mid \overline{\gamma}(\operatorname{pr}_k \overrightarrow{v}\langle i, u \rangle) = \overline{\gamma}(\operatorname{pr}_k \overrightarrow{v}\langle i, u' \rangle)\} = \{u' \in V \mid \gamma(\overrightarrow{v}\langle \overrightarrow{\jmath}, u' \cdot \overrightarrow{w} \rangle) = \gamma(\overrightarrow{v}\langle \overrightarrow{\jmath}, u \cdot \overrightarrow{w} \rangle)\}.$$

The size of the right-hand side of the latter is independent of the choice of \vec{v} from the equivalence class $[\vec{v}]_{\gamma}$. Hence, as γ is graph-like, the size of the left-hand side is independent of the choice of $\vec{w} \in V^k$, where $\operatorname{pr}_{k-r+1}\vec{w} \in [\operatorname{pr}_{k-r+1}\vec{v}]_{\overline{\gamma}}$. The result follows.

Corollary 5.8. For any graph Γ and $k, r \in \mathbb{N}$,

$$\overline{\mathbb{C}}_{k,r}(\Gamma) \succeq \overline{\mathbb{C}}_{k-r+1}(\Gamma).$$

The results proved so far involved showing that the projection of a partition satisfying some stability condition satisfied some other stability condition. In order to prove that S_{WL} dominates $S_{\mathbb{C}}$, we *extend*

a WL_k-stable partition to a \mathbb{C}_{k+1} -stable one. Let γ be graph-like and be a WL_k-stable partition of V^k . Define $\hat{\gamma} \in \mathcal{P}(V^{k+1})$ as

$$\hat{\gamma}(\vec{v}, w) = (\gamma(\vec{v}), \gamma(\vec{v}\langle 1, w \rangle), \gamma(\vec{v}\langle 2, w \rangle), \dots, \gamma(\vec{v}\langle k, w \rangle)), \quad \forall \vec{v} \in V^k, w \in V.$$

That is to say, $\hat{\gamma}(\vec{v}, w) = \hat{\gamma}(\vec{v}', w')$ if and only if $\gamma(\vec{v}) = \gamma(\vec{v}')$ and $\gamma(\vec{v}\langle i, w\rangle) = \gamma(\vec{v}'\langle i, w'\rangle)$ for all $i \in [k]$. First, we need to prove that $\hat{\gamma}$ is graph-like.

Lemma 5.9. Let γ and $\hat{\gamma}$ be as above. If γ is invariant, then $\hat{\gamma}$ is invariant.

Proof. Because $\operatorname{Sym}(k+1) = \langle \operatorname{Sym}(k), (k, k+1) \rangle$, it is sufficient to show that $\hat{\gamma}(\vec{u} \cdot u') = \hat{\gamma}(\vec{v} \cdot v') \Longrightarrow \hat{\gamma}((\vec{u} \cdot u')^{\tau}) = \hat{\gamma}((\vec{v} \cdot v')^{\tau})$, where $\tau = (k, k+1)$ is a transposition of $\operatorname{Sym}(k+1)$; or, equivalently,

$$\hat{\gamma}((\operatorname{pr}_{k-1}\vec{u}) \cdot u' \cdot u_k) = \hat{\gamma}((\operatorname{pr}_{k-1}\vec{v}) \cdot v' \cdot v_k). \tag{13}$$

From the definition of $\hat{\gamma}$, it holds that

$$\gamma(\vec{u}\langle k, u'\rangle) = \gamma(\vec{v}\langle k, v'\rangle).$$

Hence,

$$\gamma((\operatorname{pr}_{k-1}\vec{u}) \cdot u') = \gamma((\operatorname{pr}_{k-1}\vec{v}) \cdot v'). \tag{14}$$

Also, $\hat{\gamma}(\vec{u} \cdot u') = \hat{\gamma}(\vec{v} \cdot v')$ implies that $\gamma(\vec{u}) = \gamma(\vec{v})$; and hence, since γ is graph-like,

$$\gamma((\operatorname{pr}_{k-1}\vec{u}) \cdot u_k) = \gamma((\operatorname{pr}_{k-1}\vec{v}) \cdot v_k). \tag{15}$$

Since $\gamma(\vec{u}\langle i, u'\rangle) = \gamma(\vec{v}\langle i, v'\rangle)$ for all $i \in [k]$ and γ is invariant; for any $\tau = (i, k) \in \operatorname{Sym}(k)$, $\gamma(\vec{u}\langle i, u'\rangle^{\tau}) = \gamma(\vec{v}\langle i, v'\rangle^{\tau})$ or, equivalently,

$$\gamma((\operatorname{pr}_{k-1}\vec{u}) \cdot u'\langle i, u_k \rangle) = \gamma((\operatorname{pr}_{k-1}\vec{v}) \cdot v'\langle i, v_k \rangle). \tag{16}$$

Combining (14), (15), and (16), we see that (13) follows.

Consequently, if γ is graph-like, so is $\hat{\gamma}$. Indeed, $\hat{\gamma}$ is k-consistent by construction. Also, since $\hat{\gamma}$ is invariant and γ satisfies (6), then $\hat{\gamma}$ also satisfies (6). Finally, observe that if γ is a Γ -partition for some graph Γ , then so is $\hat{\gamma}$.

Lemma 5.10. Let γ and $\hat{\gamma}$ be as above. If γ is graph-like and WL_k stable, then $\hat{\gamma}$ is \mathbb{C}_{k+1} -stable.

Proof. As observed, if γ is graph-like, then $\hat{\gamma}$ is graph-like by construction. Since γ is WL_k-stable, it follows that for all $\sigma \in \text{Im } \hat{\gamma}$ and $\vec{v} \in V^{k+1}$, the size of $\{x \in V \mid \hat{\gamma}(\vec{v}\langle k+1, x\rangle) = \sigma\}$ is independent of the choice of \vec{v} from the equivalence class $[\vec{v}]_{\hat{\gamma}}$. By invariance of $\hat{\gamma}$, one deduces that for all $i \in [k+1]$, the size of $\{x \in V \mid \hat{\gamma}(\vec{v}\langle i, x\rangle) = \sigma\}$ is independent of the choice of \vec{v} in the equivalence class $[\vec{v}]_{\hat{\gamma}}$. Thus, $\hat{\gamma}$ is \mathbb{C}_{k+1} -stable.

Corollary 5.11. *For any graph* Γ ,

$$\overline{\mathbb{C}}_{k+1}(\Gamma) \succeq \overline{\mathrm{WL}}_k(\Gamma).$$

In particular, combining this with Corollary 5.6, one deduces that for all graphs Γ and $k \in \mathbb{N}$,

$$\overline{\mathbb{C}}_{k+1}(\Gamma) \approx \overline{\mathrm{WL}}_k(\Gamma). \tag{17}$$

This can be seen as a combinatorial reformulation of Theorem 5.2 in [Cai et al. 1992] proven without referring to the Immerman–Lander pebble game.

Proof of Theorem 1.3. Let Γ be any graph. From (17) it easily follows that $S_{WL} \simeq S_C$.

It follows from Corollaries 5.4 and 5.8 that for all $r, k \in \mathbb{N}$,

$$\overline{\mathbb{C}}_k(\Gamma) \succ \overline{\mathbb{C}}_{k,r}(\Gamma) \succ \overline{\mathbb{C}}_{k-r+1}(\Gamma);$$

and hence, for any $r \in \mathbb{N}$,

$$S_{\mathbf{C}} \simeq S_{\mathbf{C},r}$$
.

From Corollaries 5.6 and 5.2, we deduce that for all $r \in \mathbb{N}$,

$$S_{\text{WL},r} \simeq S_{\text{C},r}.$$

6. The Evdokimov-Ponomarenko SPAS

In this section, we apply the language and results of this paper to derive the following statement about the SPAS, $S_{EP} = \{EP_1, EP_2, ...\}$, introduced by Evdokimov and Ponomarenko in [1999].

Theorem 6.1 (Evdokimov, Ponomarenko, 1999). For any graph Γ and $k \in \mathbb{N}$,

$$\overline{\mathrm{WL}}_k(\Gamma) \leq \mathrm{EP}_k(\Gamma) \leq \overline{\mathrm{WL}}_{3k}(\Gamma).$$
 (18)

In particular, $S_{EP} \simeq S_{WL}$.

We first define the mapping EP_k in terms of the refinement operator WL_2 .

For $P \in \mathcal{P}(V^2)$, $k \in \mathbb{N}$, and some symbol Δ , set $P^{(k)}: (V^k)^2 \to \operatorname{Im} P^k \cup \operatorname{Im} P^k \times {\Delta}$ to be the following labelled partition of $(V^k)^2$:

$$P^{(k)}(\vec{u}, \vec{v}) = \begin{cases} (P(u_1, v_1), P(u_2, v_2), \dots, P(u_k, v_k), \Delta) & \text{if } \vec{u} = \vec{v} = (u, u, \dots, u) \text{ for some } u \in V, \\ (P(u_1, v_1), P(u_2, v_2), \dots, P(u_k, v_k)) & \text{otherwise,} \end{cases}$$

for all \vec{u} , $\vec{v} \in V^k$. Note that we have denoted elements of $(V^k)^2$ by (\vec{u}, \vec{v}) as opposed to $(\vec{u} \cdot \vec{v})$ to emphasize the difference between a set of *pairs* of *k*-tuples as opposed to a set of 2k-tuples. Hence, if Γ is a graph, then $[\Gamma^{(k)}]^{\text{WL}_2}$ is a coherent configuration on V^k , which we denote by $\hat{\Gamma}^{(k)}$. Observe that the set

$$I_{\Delta} = \{(u, u, \dots, u) \mid u \in V\} \subset V^k$$

is a union of cells of $\hat{\Gamma}^{(k)}$. Hence, the restriction $\hat{\Gamma}^{(k)}|_{I^2_{\Delta}}$ is a coherent configuration on I_{Δ} . Let $\delta: V^2 \to I^2_{\Delta}$ be the map

$$(u, v) \mapsto (\vec{u}, \vec{v}),$$
 (19)

where $\vec{u} = (u, u, \dots, u) \in V^k$ and $\vec{v} = (v, v, \dots, v) \in V^k$. We define $EP_k(\Gamma) = \hat{\Gamma}^{(k)} \circ \delta$, which is a coherent configuration on V.

Theorem 1.1 in [Evdokimov et al. 1999] shows that the family $S_{EP} = \{EP_1, EP_2, ...\}$ forms a SPAS. Furthermore, the authors also prove the following properties for a binary relation R on V. Set $X_R^k = \{(u, v, ..., v) \in V^k \mid (u, v) \in R\}$.

Proposition 6.2 (Proposition 3.6 in [Evdokimov et al. 1999]). $R = \{(u, v) \in V^2 \mid EP_k(\Gamma)(u, v) = \sigma\}$ for some $\sigma \in Im EP_k(\Gamma)$ if and only if X_R^k is a cell of $\hat{\Gamma}^{(k)}$.

Proposition 6.3 (Proposition 3.6 in [Evdokimov et al. 1999]). The equivalence classes of $[\alpha_{k,\Gamma}]^{WL_k}$ are unions of cells of $\hat{\Gamma}^{(k)}$.

From these, one can deduce the left-most relation of (18).

Lemma 6.4. For any graph Γ and $k \in \mathbb{N}$,

$$\overline{\mathrm{WL}}_k(\Gamma) \leq \mathrm{EP}_k(\Gamma)$$
.

Proof. For k=1, the statement is trivial. Let $k \geq 2$, and set $\Lambda = [\alpha_{k,\Gamma}]^{\operatorname{WL}_k}$. By Proposition 6.3 and because Λ is graph-like, for any $u, v \in V$ the equivalence class $[(v, u, \dots, u)]_{\Lambda}$ is a union of cells of $\hat{\Gamma}^{(k)}$ whose elements are all of the form (v', u', \dots, u') . Since $\overline{\operatorname{WL}}_k(\Gamma) = \operatorname{pr}_2 \Lambda$, it follows from Proposition 6.2 that its equivalence classes are unions of equivalence classes of $\operatorname{EP}_k(\Gamma)$. Hence $\overline{\operatorname{WL}}_k(\Gamma) \leq \operatorname{EP}_k(\Gamma)$. \square

We now apply Lemmas 5.1 and 5.5 to show the right-most relation of (18). For all $k, p \in \mathbb{N}$, let $\psi_{k,p}: V^{pk} \to (V^k)^p$ be the map

$$\vec{v} \mapsto (\vec{w}_1, \vec{w}_2, \dots, \vec{w}_p),$$

where $\vec{w}_i = (v_{1+(i-1)k}, v_{2+(i-1)k}, \dots, v_{k+(i-1)k}).$

Lemma 6.5. For any graph Γ and $k \in \mathbb{N}$,

$$EP_k(\Gamma) \leq \overline{WL}_{3k}(\Gamma)$$
.

Proof. Let $\Lambda = [\alpha_{3k,\Gamma}]^{\mathbb{C}_{3k,k}}$, and set $\hat{\Phi} = \Phi^{\mathrm{pr}_{2k}\Lambda,k}$. By Lemma 5.5, $\mathrm{pr}_{2k}\Lambda$ is a $\mathrm{WL}_{2k,k}$ -stable partition of V^{2k} . Hence, for all $\vec{\phi} \in \hat{\Phi}$, the size of

$$\{\vec{x} \in V^k \mid \operatorname{pr}_{2k} \Lambda(\vec{v}\langle \vec{i}, \vec{x} \rangle) = \phi_{\vec{i}}, \ \forall \vec{i} \in [2k]^{(k)}\}$$

is independent of the choice of \vec{v} from the equivalence class $[\vec{v}]_{\text{pr}_{2k}\Lambda}$. In particular, if $T = \{(1, 2, ..., k), (k+1, k+2, ..., 2k)\}$, then for all $\vec{\xi} \in (\text{Im pr}_{2k}\Lambda)^T$,

$$\{\vec{x} \in V^k \mid \operatorname{pr}_{2k} \Lambda(\vec{v} \langle \vec{\imath}, \vec{x} \rangle) = \xi_i, \ \forall \vec{\imath} \in T\} = \bigcup_{\substack{\vec{\phi} \in \hat{\Phi} \\ \phi_{\vec{\imath}} = \xi_{\vec{\imath}}, \ \forall \vec{\imath} \in T}} \{\vec{x} \in V^k \mid \operatorname{pr}_{2k} \Lambda(\vec{v} \langle \vec{\imath}, \vec{x} \rangle) = \phi_{\vec{\imath}}, \ \forall \vec{\imath} \in [2k]^{(k)}\}.$$

The right-hand side of the above is a disjoint union of sets whose size is independent of the choice of \vec{v} from the equivalence class $[\vec{v}]_{\text{pr}_{2k}\Lambda}$. Hence, the size of the left-hand side is independent of the choice of \vec{v} from the equivalence class $[\vec{v}]_{\text{pr}_{2k}\Lambda}$; and therefore, $(\text{pr}_{2k}\Lambda) \circ \psi_{k,2}^{-1}$ is a WL₂-stable partition of $(V^k)^2$. In particular, it refines $\alpha_{2k,\Gamma} \circ \psi_{k,2}^{-1}$. But $\Gamma^{(k)} \leq \alpha_{2k,\Gamma} \circ \psi_{k,2}^{-1}$; and thus, since $\hat{\gamma}^{(k)}$ is a minimal WL₂-stable partition of V^{2k} refining $\Gamma^{(k)}$, then $\hat{\Gamma}^{(k)} \leq (\text{pr}_{2k}\Lambda) \circ \psi_{k,2}^{-1}$. In particular, $\hat{\Gamma}^{(k)}|_{I_{\Delta}^2} \leq (\text{pr}_{2k}\Lambda) \circ \psi_{k,2}^{-1}|_{I_{\Delta}^2}$; and hence, if δ is as defined in (19),

$$\mathrm{EP}_k(\Gamma) \leq (\mathrm{pr}_{2k}\Lambda) \circ \psi_{k,2}^{-1} \circ \delta.$$

Since Λ is graph-like, Lemma 3.6 implies that for all $u, u', v, v' \in V$,

$$\operatorname{pr}_{2k}\Lambda(u,v,\ldots,v) = \operatorname{pr}_{2k}\Lambda(u',v',\ldots,v') \Longrightarrow \operatorname{pr}_{2k}\Lambda(\vec{u}\cdot\vec{v}) = \operatorname{pr}_{2k}\Lambda(\vec{u}'\cdot\vec{v}), \tag{20}$$

where $\vec{u} = (u, u, \dots, u), \vec{v} = (v, v, \dots, v), \vec{u}' = (u', u', \dots, u'), \vec{v}' = (v', v', \dots, v') \in V^k$. It follows that $\operatorname{pr}_2 \Lambda = (\operatorname{pr}_{2k} \Lambda) \circ \psi_{k,2}^{-1} \circ \delta$. But $\operatorname{pr}_2 \Lambda = \overline{\mathbb{C}}_{3k,k}(\Gamma) \preceq \overline{\operatorname{WL}}_{3k}(\Gamma)$. Thus, by (20),

$$EP_k(\Gamma) \prec \overline{WL}_{3k}(\Gamma).$$

7. Proofs of Theorems 1.4 and 1.5

We proceed using the same strategy as for the proof of Theorem 1.3.

Lemma 7.1. The k-projection of a graph-like $\mathrm{IM}_{k+1}^{\mathbb{F}}$ -stable partition is \mathbb{C}_k -stable for any field \mathbb{F} and $k \in \mathbb{N}$

Proof. Let $\gamma \in \mathcal{P}(V^{k+2})$ be a graph-like $\mathrm{IM}_{k+2}^{\mathbb{F}}$ -stable partition, and set $\overline{\gamma} = \mathrm{pr}_k \gamma$. Suppose $\gamma(\vec{v}) = \gamma(\vec{v}')$. Fix some $j \in [k]$, and let $\vec{\iota} = (j, k+1) \in [k+1]^{(2)}$. As γ is graph-like, there is some $\sigma \in \mathrm{Im}\,\gamma$ such that the matrix $\chi_{\sigma, \vec{\iota}}^{\gamma, \vec{v}}$ is diagonal and nonzero. Since $\gamma(\vec{v}) = \gamma(\vec{v}')$, there is some $S \in \mathrm{GL}_V(\mathbb{F})$ such that $S\chi_{\sigma, \vec{\iota}}^{\gamma, \vec{v}} S^{-1} = \chi_{\sigma, \vec{\iota}}^{\gamma, \vec{v}'}$. So $\chi_{\sigma, \vec{\iota}}^{\gamma, \vec{v}'}$ and $\chi_{\sigma, \vec{\iota}}^{\gamma, \vec{v}'}$ have the same rank and hence, the same number of 1s on the diagonal. Thus, for any $w \in V$, the size of the set $\{u \in V \mid \gamma(\vec{v}\langle\vec{\iota}, (u, u)\rangle) = \sigma\}$ is independent of the choice of \vec{v} from the equivalence class $[\vec{v}]_{\gamma}$. Also,

$$\{u \in V \mid \overline{\gamma}(\operatorname{pr}_{k} \overrightarrow{v}\langle j, u \rangle) = \sigma\} = \{u \in V \mid \gamma(\overrightarrow{v}\langle \overrightarrow{\iota}, (u, u) \rangle) = \sigma\},$$

and hence, as γ is graph-like, the size of the left-hand side is independent of the choice of $\operatorname{pr}_k \vec{v}$ from the equivalence class $[\operatorname{pr}_k \vec{v}]_{\vec{v}}$. The result follows.

Corollary 7.2. For any graph Γ and $k \in \mathbb{N}$,

$$\overline{\mathrm{IM}}_{k+2}^{\mathbb{F}}(\Gamma) \succeq \overline{\mathrm{WL}}_{k}(\Gamma).$$

Proof. This follows from Lemma 7.1 and the fact that $\overline{\mathbb{C}}_{k+1}(\Gamma) \approx \overline{\mathrm{WL}}_k(\Gamma)$.

Lemma 7.3. The k-projection of a graph-like \mathbb{C}_{k+1} -stable partition is $\mathrm{IM}_k^{\mathbb{F}}$ -stable for any \mathbb{F} with $\mathrm{char}(\mathbb{F}) = 0$ or $\mathrm{char}(\mathbb{F}) > |V|$.

Proof. Let $\gamma \in \mathcal{P}(V^{k+1})$ be a graph-like \mathbb{C}_{k+1} -stable partition, and fix some $\vec{\iota} \in [k+1]^{(3)}$ with $i_3 = k+1$. For every $\vec{v} \in V^{k+1}$, define $g_{\vec{v}}$ to be the partition of V^3 given by $g_{\vec{v}}(\vec{x}) = \gamma(\vec{v}\langle\vec{\iota},\vec{x}\rangle)$ for all $\vec{x} \in V^3$. Since γ is \mathbb{C}_{k+1} -stable, it follows that $g_{\vec{v}}$ is \mathbb{C}_3 stable for all $\vec{v} \in V^{k+1}$. As γ is graph-like, $\operatorname{pr}_2 g_{\vec{v}}$ is a rainbow, and it is WL₂-stable by Lemma 5.1 and thus, a coherent configuration on V.

Set $\bar{\gamma} = \operatorname{pr}_k \gamma$, and for all $\vec{y} \in V^2$, let $\bar{g}_{\vec{v}}(\vec{y}) = \bar{\gamma}(\operatorname{pr}_k \vec{v}\langle (i_1, i_2), \vec{y} \rangle)$. Then $\bar{g}_{\vec{v}} = \operatorname{pr}_2 g_{\vec{v}}$. Therefore, all nonempty relations of $\bar{g}_{\vec{v}}$ form a coherent configuration whose \mathbb{F} -adjacency algebra has standard basis

$$\{\chi_{(i_1,i_2),\sigma}^{\bar{\gamma},\operatorname{pr}_k\vec{v}}\mid \exists \vec{y}\in V^2, \, \bar{\gamma}(\operatorname{pr}_k\vec{v}\langle(i_1,i_2),\vec{y}\rangle) = \sigma\}.$$

Thus, if $\bar{\gamma}(\operatorname{pr}_k\vec{u}) = \bar{\gamma}(\operatorname{pr}_k\vec{w})$, then $\bar{g}_{\vec{u}}$ and $\bar{g}_{\vec{w}}$ are algebraically isomorphic coherent configurations. More precisely, one can check that $\operatorname{Im} \bar{g}_{\vec{u}} = \operatorname{Im} \bar{g}_{\vec{v}}$ and that the map

$$\iota: \operatorname{Im} \bar{g}_{\vec{u}} \to \operatorname{Im} \bar{g}_{\vec{v}}, \quad \sigma \mapsto \sigma$$

is an algebraic isomorphism.

For char(\mathbb{F}) = 0 or char(\mathbb{F}) > |V|, it follows from Corollary 2.11 that there is some $S \in GL_V(\mathbb{F})$ such that $S\chi_{(i_1,i_2),\sigma}^{\bar{\gamma},\operatorname{pr}_k\vec{u}}S^{-1} = \chi_{(i_1,i_2),\sigma}^{\bar{\gamma},\operatorname{pr}_k\vec{u}}$ for all $\sigma \in \operatorname{Im}\bar{\gamma}$. The result then follows.

Note that in the above proof, there may be other bijections $\operatorname{Im} \bar{g}_{\vec{u}} \to \operatorname{Im} \bar{g}_{\vec{v}}$ which are algebraic isomorphisms. However, it follows from the definition of the k-refinement operator $\operatorname{IM}_k^{\mathbb{F}}$, that $\vec{u}, \vec{v} \in V^k$ are in the same equivalence class of an $\operatorname{IM}_k^{\mathbb{F}}$ -stable partition, only if ι , as above, is an algebraic isomorphism.

Corollary 7.4. For all $k \in \mathbb{N}$, graph Γ with vertex set V, and field \mathbb{F} such that $\operatorname{char}(\mathbb{F}) = 0$ or $\operatorname{char}(\mathbb{F}) > |V|$,

$$\overline{\mathrm{WL}}_k(\Gamma) \succeq \overline{\mathrm{IM}}_k^{\mathbb{F}}(\Gamma).$$

Proof. This follows from Lemma 7.3 and the fact that $\overline{\mathbb{C}}_{k+1}(\Gamma) \approx \overline{\mathrm{WL}}_k(\Gamma)$.

The statement of Theorem 1.4 comes from taking \mathbb{F} to be of characteristic 0 in Corollaries 7.2 and 7.4. Theorem 1.5 arises instead from combining Corollary 7.4 with a construction due to Holm [2010]. The construction in the proof of Theorem 7.1 in [Holm 2010] gives, for each $k \in \mathbb{N}$ and prime number p, a graph $\Gamma_{k,p}$ for which $\overline{\mathrm{WL}}_k(\Gamma_{k,p})$ is *strictly coarser* than $\mathrm{Sch}(\Gamma_{k,p})$, but $\overline{\mathrm{IM}}_3^{\mathbb{F}}(\Gamma_{k,p}) = \mathrm{Sch}(\Gamma_{k,p})$ for any field \mathbb{F} with $\mathrm{char}(\mathbb{F}) = p$. Furthermore, from the same result, it also follows that $\overline{\mathrm{IM}}_k^{\mathbb{F}}(\Gamma_{k,p})$ is strictly coarser than $\mathrm{Sch}(\Gamma_{k,p})$ whenever $\mathrm{char}(\mathbb{F}) \neq p$. This shows that the SPAS, $\mathcal{S}_{\mathrm{IM}}(\mathbb{F}_1)$ and $\mathcal{S}_{\mathrm{IM}}(\mathbb{F}_2)$, are incomparable whenever $\mathrm{char}(\mathbb{F}_1) \neq \mathrm{char}(\mathbb{F}_2)$.

8. Yet another refinement operator

There is a subtle difference between the definitions of $WL_{k,r}$ and $\mathbb{C}_{k,r}$: the colours of $WL_{k,r} \circ \gamma$ are *multisets of tuples* of colours of γ , whereas the colours of $\mathbb{C}_{k,r} \circ \gamma$ are *tuples of multisets* of colours of γ . We now show that a similar variation in the definition of $IM_k^{\mathbb{F}}$ gives a refinement procedure whose corresponding SPAS is equivalent to $\mathcal{S}_{IM(\mathbb{F})}$.

For every $\gamma \in \mathcal{P}(V^k)$, $\vec{v} \in V^k$, and $\vec{\phi} \in \Phi^{\gamma,2}$, let $\chi_{\vec{\phi}}^{\gamma,\vec{v}}$ be the adjacency matrix of the relation $\{\vec{x} \in V^2 \mid \gamma(\vec{v}\langle\vec{i},(u,v)\rangle) = \phi_{\vec{i}}, \ \forall \vec{i} \in [k]^{(2)}\} \subseteq V^2$. Define the mapping $\mathrm{IMt}_k^{\mathbb{F}}$ by setting $\mathrm{IMt}_1^{\mathbb{F}} \circ \gamma = \mathrm{IMt}_2^{\mathbb{F}} \circ \gamma = \gamma$, and for k > 2,

$$\operatorname{IMt}_{k}^{\mathbb{F}} \circ \gamma : V^{k} \to \operatorname{Im} \gamma \times (\operatorname{Mat}_{V}(\mathbb{F})^{\Phi^{\gamma,2}}/\sim), \quad \vec{v} \mapsto (\gamma(\vec{v}), (\chi_{\vec{\phi}}^{\gamma,\vec{v}})_{\vec{\phi} \in \Phi^{\gamma,2}}),$$

where the equivalence classes of the relation \sim are the orbits of $GL_V(\mathbb{F})$ acting on the tuples by conjugation.

Similarly to Proposition 4.6, one can show that $\mathrm{IMt}_k^{\mathbb{F}}$ is a graph-like k-refinement operator for all $k \in \mathbb{N}$ and that $\mathrm{IMt}(\mathbb{F}) = \{\mathrm{IMt}_1^{\mathbb{F}}, \mathrm{IMt}_2^{\mathbb{F}}, \ldots\}$ is a refinement procedure. Hence, the family $\mathcal{S}_{\mathrm{IMt}(\mathbb{F})} = \{\overline{\mathrm{IMt}}_1^{\mathbb{F}}, \overline{\mathrm{IMt}}_2^{\mathbb{F}}, \ldots\}$ is a SPAS for any field \mathbb{F} . Also, similarly to Proposition 4.5, one may derive the following stability condition:

Proposition 8.1. A partition $\gamma \in \mathcal{P}(V^k)$ is $\mathrm{IMt}_k^{\mathbb{F}}$ -stable if and only if for all $\vec{u}, \vec{v} \in V^k$,

$$\gamma(\vec{u}) = \gamma(\vec{v}) \Longrightarrow \exists S \in \mathrm{GL}_V(\mathbb{F}), \quad \forall \vec{\phi} \in \Phi^{\gamma,2} \, S \chi_{\vec{\phi}}^{\gamma,\vec{u}} S^{-1} = \chi_{\vec{\phi}}^{\gamma,\vec{v}}.$$

In particular, the following result is analogous to Proposition 4.5:

Lemma 8.2. Any $\mathrm{IMt}_k^{\mathbb{F}}$ -stable partition is also $\mathrm{IM}_k^{\mathbb{F}}$ -stable.

Proof. For any $\vec{v} \in V^k$, $\vec{i} \in [k]^{(2)}$, and $\sigma \in \text{Im } \gamma$

$$\chi_{\sigma,\vec{i}}^{\gamma,\vec{v}} = \sum_{\substack{\vec{\phi} \in \Phi^{\gamma,2} \\ \phi_{\vec{i}} = \sigma}} \chi_{\vec{\phi}}^{\gamma,\vec{v}}.$$

Hence, if $\gamma(\vec{u}) = \gamma(\vec{v})$, there is some $S \in GL_V(\mathbb{F})$ such that $S\chi_{\vec{\phi}}^{\gamma,\vec{v}}S^{-1} = \chi_{\vec{\phi}}^{\gamma,\vec{u}}$ for all $\vec{\phi} \in \Phi^{\gamma,2}$. Thus,

$$S\chi_{\sigma,\vec{i}}^{\gamma,\vec{v}}S^{-1} = \sum_{\substack{\vec{\phi} \in \Phi^{\gamma,2} \\ \phi_{\vec{i}} = \sigma}} S\chi_{\vec{\phi}}^{\gamma,\vec{v}}S^{-1} = \sum_{\substack{\vec{\phi} \in \Phi^{\gamma,2} \\ \phi_{\vec{i}} = \sigma}} \chi_{\vec{\phi}}^{\gamma,\vec{u}} = \chi_{\sigma,\vec{i}}^{\gamma,\vec{u}};$$

from which follows that γ is $IM_k^{\mathbb{F}}$ -stable.

Corollary 8.3. *For any graph* Γ , *field* \mathbb{F} , *and* $k \in \mathbb{N}$,

$$\overline{\mathrm{IM}}_{k}^{\mathbb{F}}(\Gamma) \leq \overline{\mathrm{IM}} \mathfrak{t}_{k}^{\mathbb{F}}(\Gamma).$$

In the following result, the argument, is analogous to that of Lemma 5.5:

Lemma 8.4. The k-projection of an $\mathrm{IM}_{k+2}^{\mathbb{F}}$ -stable partition is $\mathrm{IMt}_{k}^{\mathbb{F}}$ -stable.

Proof. Let \vec{u} , $\vec{v} \in V^k$ be such that $\gamma(\vec{u}) = \gamma(\vec{v})$, and set $\overline{\gamma} = \operatorname{pr}_k \gamma$. Because γ is graph-like, it follows from Lemma 3.6 that for all $\vec{\iota} \in [k+2]^{(2)}$, $\gamma(\vec{v}\langle\vec{\iota},\operatorname{pr}_{\vec{\jmath}}\vec{v}\rangle) = \gamma(\vec{u}\langle\vec{\iota},\operatorname{pr}_{\vec{\jmath}}\vec{u}\rangle)$, where $\vec{\jmath} = (k+1,k+2) \in [k+2]^{(2)}$. In particular, for all $\vec{\iota} \in [k]^{(2)}$,

$$\bar{\gamma}(\operatorname{pr}_{k}\vec{v}\langle\vec{i},\operatorname{pr}_{7}\vec{v}\rangle) = \bar{\gamma}(\operatorname{pr}_{k}\vec{u}\langle\vec{i},\operatorname{pr}_{7}\vec{u}\rangle). \tag{21}$$

Define $\vec{\Delta} \in (\Phi^{\bar{\gamma},2})^{\operatorname{Im}\gamma}$ such that

$$(\Delta_{\gamma(\vec{v})})_{\vec{i}} = \overline{\gamma}(\operatorname{pr}_{k}\vec{v}\langle\vec{i},\operatorname{pr}_{\vec{i}}\vec{v}\rangle).$$

One deduces from (21) that $\vec{\Delta}$ is well defined.

For any $\vec{\phi} \in \Phi$,

$$\{\vec{u} \in V^2 \mid \overline{\gamma}(\operatorname{pr}_k \vec{v} \langle \vec{i}, \vec{u} \rangle) = \phi_{\vec{i}}, \ \forall \vec{i} \in [k]^{(2)}\} = \bigcup_{\substack{\sigma \in \operatorname{Im} \gamma \\ \Delta_{\sigma} = \vec{\phi}}} \{\vec{u} \in V^2 \mid \gamma(\vec{v} \langle \vec{j}, \vec{u} \rangle) = \sigma\}.$$

Hence, for all $\vec{\phi} \in \Phi$ it holds that

$$\chi_{\vec{\phi}}^{\gamma,\vec{v}} = \sum_{\substack{\sigma \in \operatorname{Im} \gamma \\ \Delta_{\sigma} = \vec{\phi}}} \chi_{\sigma,\vec{J}}^{\gamma,\vec{v}}.$$

Because γ is $\mathrm{IM}_{k+2}^{\mathbb{F}}$ -stable and $\gamma(\vec{u}) = \gamma(\vec{v})$, for each $\vec{i} \in [k]^{(2)}$ there is some $S \in \mathrm{GL}_V(\mathbb{F})$ such that $S\chi_{\sigma,\vec{i}}^{\gamma,\vec{v}}S^{-1} = \chi_{\sigma,\vec{i}}^{\gamma,\vec{u}}$ for all $\sigma \in \mathrm{Im}\,\gamma$. Thus, for all $\vec{\phi} \in (\mathrm{Im}\,\bar{\gamma})^{[k]^{(2)}}$,

$$S\chi_{\vec{\phi}}^{\gamma,\vec{v}}S^{-1} = \sum_{\substack{\sigma \in \operatorname{Im} \gamma \\ \Delta_{\sigma} = \vec{\phi}}} S\chi_{\sigma,\vec{j}}^{\gamma,\vec{v}}S^{-1} = \sum_{\substack{\sigma \in \operatorname{Im} \gamma \\ \Delta_{\sigma} = \vec{\phi}}} \chi_{\sigma,\vec{j}}^{\vec{u}} = \chi_{\vec{\phi}}^{\gamma,\vec{u}};$$

whence, $\bar{\gamma}(\vec{u}) = \bar{\gamma}(\vec{v})$. Therefore, $\bar{\gamma}$ is $\mathrm{IMt}_k^{\mathbb{F}}$ -stable.

Corollary 8.5. For any graph Γ , field \mathbb{F} , and $k \in \mathbb{N}$,

$$\overline{\mathrm{IM}} \mathfrak{t}_{k}^{\mathbb{F}}(\Gamma) \prec \overline{\mathrm{IM}}_{k+2}^{\mathbb{F}}(\Gamma).$$

Theorem 1.6 follows from Corollaries 8.5 and 8.3.

9. Conclusions

The Weisfeiler–Leman algorithm is much studied in the context of graph isomorphism. It is really a family of algorithms, graded by a dimension parameter. A large number of other families of algorithms have been shown to give essentially the same graded approximations of isomorphism. The Schurian polynomial approximation schemes of Evdokimov et al. provide a general framework for comparing these families of algorithms. The invertible map operators of Dawar and Holm provide another such family of algorithms (or, more formally in the language of this paper, *refinement procedure*), but one that has greater distinguishing power than the Weisfeiler–Leman family. In the same way that \mathbf{WL}_r and \mathbf{C}_r were obtained from \mathbf{WL} and \mathbf{C} , one can generalize $\mathbf{IM}(\mathbb{F})$ as follows: for every $k, r \in \mathbb{N}$, define the k-refinement operator $\mathbf{IM}_{k,r}^{\mathbb{F}}$ by setting $\mathbf{IM}_{k,r}^{\mathbb{F}} \circ \gamma = \gamma$ when $k \leq 2r$. When k > 2r define:

$$\operatorname{IM}_{k,r}^{\mathbb{F}} \circ \gamma : V^{k} \times \to \operatorname{Im} \gamma \times (\operatorname{Mat}_{V^{r}}(\mathbb{F})^{\operatorname{Im} \gamma \times [k]^{(2r)}} / \sim), \quad \vec{v} \mapsto (\gamma(\vec{v}), ((\chi_{\vec{t}, \sigma}^{\gamma, \vec{v}})_{\sigma \in \operatorname{Im} \gamma})_{\vec{t} \in [k]^{(2r)}}),$$

where $\chi_{\vec{\imath},\sigma}^{\gamma,\vec{\imath}}$ is the adjacency matrix of the relation $\{(\vec{x},\vec{y}) \mid \gamma(\vec{\imath}(\vec{\imath},\vec{x}\cdot\vec{y})) = \sigma\} \subseteq (V^r)^2$ and \sim is the relation whose equivalence classes are the orbits of $GL_{V^r}(\mathbb{F})$ acting on the tuples by conjugation. One can show that $\mathbf{IM}_r(\mathbb{F}) = \{\mathbf{IM}_{1,r}^{\mathbb{F}}, \mathbf{IM}_{2,r}^{\mathbb{F}}, \dots\}$ is a refinement procedure for all $r \in \mathbb{N}$. One can thus derive from it a SPAS, $S_{\mathbf{IM}(\mathbb{F}),r}$, in the same manner as described in Section 4. While we were able to show that the refinement procedures \mathbf{WL}_r and \mathbf{C}_r do not yield SPASs more powerful than that yielded by \mathbf{WL} , the exact relation between $S_{\mathbf{IM}(\mathbb{F}),r}$ and $S_{\mathbf{IM}(\mathbb{F}),r}$ is still unclear and an interesting open question.

References

[Babai 2016] L. Babai, "Graph isomorphism in quasipolynomial time", pp. 684–697 in STOC'16: Proceedings of the 48th Annual ACM SIGACT Symposium on Theory of Computing (Cambridge, 2016), edited by D. Wichs and Y. Mansour, ACM, New York, 2016. MR Zbl

[Brooksbank and Luks 2008] P. A. Brooksbank and E. M. Luks, "Testing isomorphism of modules", *J. Algebra* **320**:11 (2008), 4020–4029. MR Zbl

[Cai et al. 1992] J.-Y. Cai, M. Fürer, and N. Immerman, "An optimal lower bound on the number of variables for graph identification", *Combinatorica* 12:4 (1992), 389–410. MR Zbl

[Cameron 1999] P. J. Cameron, *Permutation groups*, Lond. Math. Soc. Student Texts **45**, Cambridge Univ. Press, 1999. MR Zbl

[Chen and Ponomarenko 2019] G. Chen and I. Ponomarenko, *Coherent configurations*, Central China Normal Univ. Press, Wuhan, 2019.

[Chistov et al. 1997] A. Chistov, G. Ivanyos, and M. Karpinski, "Polynomial time algorithms for modules over finite dimensional algebras", pp. 68–74 in *Proceedings of the 1997 International Symposium on Symbolic and Algebraic Computation* (Kihei, HI, 1997), edited by W. W. Küchlin, ACM, New York, 1997. MR Zbl

[Dawar and Holm 2017] A. Dawar and B. Holm, "Pebble games with algebraic rules", Fund. Inform. 150:3-4 (2017), 281–316. MR Zbl

[Dawar et al. 2019] A. Dawar, E. Grädel, and W. Pakusa, "Approximations of isomorphism and logics with linear-algebraic operators", art. id. 112 in 46th International Colloquium on Automata, Languages, and Programming (Patras, Greece, 2019), edited by C. Baier et al., LIPIcs. Leibniz Int. Proc. Inform. 132, Schloss Dagstuhl. Leibniz-Zent. Inform., Wadern, Germany, 2019. MR

[Evdokimov and Ponomarenko 1999] S. Evdokimov and I. Ponomarenko, "On highly closed cellular algebras and highly closed isomorphisms", *Electron. J. Combin.* **6** (1999), art. id. 18. MR Zbl

[Evdokimov et al. 1999] S. Evdokimov, M. Karpinski, and I. Ponomarenko, "On a new high-dimensional Weisfeiler–Lehman algorithm", *J. Algebraic Combin.* **10**:1 (1999), 29–45. MR Zbl

[Higman 1987] D. G. Higman, "Coherent algebras", Linear Algebra Appl. 93 (1987), 209-239. MR Zbl

[Holm 2010] B. Holm, *Descriptive complexity of linear algebra*, Ph.D. thesis, University of Cambridge, 2010, Available at https://tinyurl.com/holmbphd.

[Mathon 1979] R. Mathon, "A note on the graph isomorphism counting problem", *Inform. Process. Lett.* **8**:3 (1979), 131–136. MR Zbl

[Zieschang 1996] P.-H. Zieschang, An algebraic approach to association schemes, Lecture Notes in Math. 1628, Springer, 1996. MR Zbl

Received 3 Aug 2019. Revised 14 May 2020.

ANUJ DAWAR:

anuj.dawar@cl.cam.ac.uk

Department of Computer Science and Technology, University of Cambridge, Cambridge, United Kingdom

DANNY VAGNOZZI:

danny.vagnozzi@gmail.com

Department of Computer Science and Technology, University of Cambridge, Cambridge, United Kingdom

Moscow Journal of Combinatorics and Number Theory

msp.org/moscow

EDITORS-IN-CHIEF

Yann Bugeaud Université de Strasbourg (France)

bugeaud@math.unistra.fr

Nikolay Moshchevitin Lomonosov Moscow State University (Russia)

moshchevitin@gmail.com

Andrei Raigorodskii Moscow Institute of Physics and Technology (Russia)

mraigor@yandex.ru

Ilya D. Shkredov Steklov Mathematical Institute (Russia)

ilya.shkredov@gmail.com

EDITORIAL BOARD

Iskander Aliev Cardiff University (United Kingdom)

Vladimir Dolnikov Moscow Institute of Physics and Technology (Russia)

Nikolay Dolbilin Steklov Mathematical Institute (Russia)

Oleg German Moscow Lomonosov State University (Russia)

Michael Hoffman United States Naval Academy

Grigory Kabatiansky Russian Academy of Sciences (Russia)

Roman Karasev Moscow Institute of Physics and Technology (Russia)

Gyula O. H. Katona Hungarian Academy of Sciences (Hungary)

Alex V. Kontorovich Rutgers University (United States)

Maxim Korolev Steklov Mathematical Institute (Russia)

Christian Krattenthaler Universität Wien (Austria)

Antanas Laurinčikas Vilnius University (Lithuania)

Vsevolod Lev University of Haifa at Oranim (Israel)

János Pach EPFL Lausanne(Switzerland) and Rényi Institute (Hungary)

Rom Pinchasi Israel Institute of Technology – Technion (Israel)

Alexander Razborov Institut de Mathématiques de Luminy (France)

Joël Rivat Université d'Aix-Marseille (France) Tanguy Rivoal Institut Fourier, CNRS (France)

Damien Roy University of Ottawa (Canada)

Vladislav Salikhov Bryansk State Technical University (Russia)
Tom Sanders University of Oxford (United Kingdom)

Alexander A. Sapozhenko Lomonosov Moscow State University (Russia)

József Solymosi University of British Columbia (Canada)

Andreas Strömbergsson Uppsala University (Sweden)

Benjamin Sudakov University of California, Los Angeles (United States)

Jörg Thuswaldner University of Leoben (Austria) Kai-Man Tsang Hong Kong University (China)

Maryna Viazovska EPFL Lausanne (Switzerland)

Barak Weiss Tel Aviv University (Israel)

PRODUCTION

Silvio Levy (Scientific Editor)

production@msp.org

Cover design: Blake Knoll, Alex Scorpan and Silvio Levy

See inside back cover or msp.org/moscow for submission instructions.

The subscription price for 2020 is US \$310/year for the electronic version, and \$365/year (+\$20, if shipping outside the US) for print and electronic. Subscriptions, requests for back issues and changes of subscriber address should be sent to MSP.

Moscow Journal of Combinatorics and Number Theory (ISSN 2640-7361 electronic, 2220-5438 printed) at Mathematical Sciences Publishers, 798 Evans Hall #3840, c/o University of California, Berkeley, CA 94720-3840 is published continuously online. Periodical rate postage paid at Berkeley, CA 94704, and additional mailing offices.

MJCNT peer review and production are managed by EditFlow® from MSP.

PUBLISHED BY

mathematical sciences publishers

nonprofit scientific publishing

http://msp.org/

© 2020 Mathematical Sciences Publishers

Thematic Issue on Complexity

Provenance analysis for logic and games	203
Erich Grädel and Val Tannen	
Generalizations of k-dimensional Weisfeiler–Leman stabilization	229
Anuj Dawar and Danny Vagnozzi	
On defining linear orders by automata	253
Bruno Courcelle, Irène Durand and Michael Raskin	
First-order definitions of subgraph isomorphism through the adjacency and order relations	293
OLEG GRIGORYAN, MIKHAIL MAKAROV and MAKSIM ZHUKOVSKII	
On limit points of spectra of first-order sentences with quantifier depth 4 YURY YAROVIKOV	303
Weakly distinguishing graph polynomials on addable properties JOHANN A. MAKOWSKY and VSEVOLOD RAKITA	333